Logic & Probabilistic Circuits

Representation Reasoning Theory

Guy Van den Broeck University of California, Los Angeles guyvdb@cs.ucla.edu **YooJung Choi** Arizona State University yj.choi@asu.edu

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Objective

Circuits are an assembly language for tractable logic and probabilistic reasoning

Even though <u>logic</u> is central to this Simons program, we will couch this tutorial in <u>probability</u>...

- Most AI and DB interest in tractable logic circuits for the past 15 years has been as a means of doing probabilistic inference
- Much richer query languages <3
- We live in the age of probabilistic generative Al... :-)

We will spare you most of the machine learning details, and instead focus on representations, query languages, reasoning algorithms, and connections to theory.

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1. What are probabilistic queries? Are current models tractable? (Guy)

- 2. What are probabilistic circuits and why are they tractable? (Guy)
- 3. What is the connection to logical circuit languages? (YooJung)
- 4. How do I compile my favorite model into a circuit? (YooJung)
- 5. How are circuit size and tractability related? (YooJung)
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Acknowledgements

This tutorial is based on our (joint) tutorials and slides from



- Robert Peharz
- Nicola Di Mauro







The Alphabet Soup of probabilistic models



Intractable and tractable models



tractability is a spectrum



Expressive models without compromises



a unifying framework for tractable models

Why tractable inference?

or the inherent trade-off of tractability vs. expressiveness

q₁: What is the probability that today is a Monday and there is a traffic jam on Westwood Blvd.?



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q₁: What is the probability that today is a Monday and there is a traffic jam on Westwood Blvd.?

$$\begin{split} \mathbf{X} &= \{\mathsf{Day},\mathsf{Time},\mathsf{Jam}_{\mathsf{Str1}},\mathsf{Jam}_{\mathsf{Str2}},\ldots,\mathsf{Jam}_{\mathsf{StrN}}\}\\ \mathbf{q}_1(\mathbf{m}) &= p_{\mathbf{m}}(\mathsf{Day}=\mathsf{Mon},\mathsf{Jam}_{\mathsf{Wwood}}=1) \end{split}$$



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q₁: What is the probability that today is a Monday and there is a traffic jam on Westwood Blvd.?

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 \Rightarrow marginals



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q₂: Which day is most likely to have a traffic jam on my route to campus?

 $\mathbf{X} = \{\mathsf{Day}, \mathsf{Time}, \mathsf{Jam}_{\mathsf{Str1}}, \mathsf{Jam}_{\mathsf{Str2}}, \dots, \mathsf{Jam}_{\mathsf{StrN}}\}$

$$\mathbf{q}_2(\mathbf{m}) = \operatorname{argmax}_{\mathsf{d}} p_{\mathbf{m}}(\mathsf{Day} = \mathsf{d} \land igvee_{i \in \mathsf{route}} \operatorname{\mathsf{Jam}}_{\mathsf{Str}i})$$



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 \Rightarrow marginals + MAP + logical events



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Tractable Probabilistic Inference

A class of queries Q is tractable on a family of probabilistic models \mathcal{M} iff for any query $\mathbf{q} \in Q$ and model $\mathbf{m} \in \mathcal{M}$ **exactly** computing $\mathbf{q}(\mathbf{m})$ runs in time $O(\operatorname{poly}(|\mathbf{m}|))$.

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 \Rightarrow often poly will in fact be **linear**!



tractable bands

Complete evidence (EVI)

q₃: What is the probability that today is a Monday at 12.00 and there is a traffic jam only on Westwood Blvd.?



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q₃: What is the probability that today is a Monday at 12.00 and there is a traffic jam only on Westwood Blvd.?

$$\begin{split} \mathbf{X} &= \{\mathsf{Day},\mathsf{Time},\mathsf{Jam}_{\mathsf{Wwood}},\mathsf{Jam}_{\mathsf{Str2}},\ldots,\mathsf{Jam}_{\mathsf{StrN}}\}\\ \mathbf{q}_3(\mathbf{m}) &= p_{\mathbf{m}}(\mathbf{X} = \{\mathsf{Mon},12.00,1,0,\ldots,0\}) \end{split}$$



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Complete evidence (EVI)

q₃: What is the probability that today is a Monday at 12.00 and there is a traffic jam only on Westwood Blvd.?

$$f X = \{Day, Time, Jam_{Wwood}, Jam_{Str2}, \dots, Jam_{StrN}\}$$

 $f q_3(f m) = p_{f m}(f X = \{Mon, 12.00, 1, 0, \dots, 0\})$

...fundamental in *maximum likelihood learning*

$$\theta_{\mathbf{m}}^{\mathsf{MLE}} = \operatorname{argmax}_{\theta} \prod_{\mathbf{x} \in \mathcal{D}} p_{\mathbf{m}}(\mathbf{x}; \theta)$$



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Generative Adversarial Networks

$$\min_{\theta} \max_{\phi} \mathbb{E}_{\mathbf{x} \sim p_{\mathsf{data}}(\mathbf{x})} \left[\log D_{\phi}(\mathbf{x}) \right] + \mathbb{E}_{\mathbf{z} \sim p(\mathbf{z})} \left[\log(1 - D_{\phi}(G_{\theta}(\mathbf{z}))) \right]$$



Goodfellow et al., "Generative adversarial nets", 2014

Generative Adversarial Networks

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Goodfellow et al., "Generative adversarial nets", 2014



tractable bands

Variational Autoencoders

 $p_{\theta}(\mathbf{x}) = \int p_{\theta}(\mathbf{x} \mid \mathbf{z}) p(\mathbf{z}) d\mathbf{z}$

an explicit likelihood model!



Rezende et al., "Stochastic backprop. and approximate inference in deep generative models", 2014 Kingma and Welling, "Auto-Encoding Variational Bayes", 2014

Variational Autoencoders

$$\log p_{\theta}(\mathbf{x}) \geq \mathbb{E}_{\mathbf{z} \sim q_{\phi}(\mathbf{z} \mid \mathbf{x})} \left[\log p_{\theta}(\mathbf{x} \mid \mathbf{z}) \right] - \mathbb{KL}(q_{\phi}(\mathbf{z} \mid \mathbf{x}) || p(\mathbf{z}))$$

an explicit likelihood model!

... but computing $\log p_{ heta}(\mathbf{x})$ is intractable

 \Rightarrow an infinite and uncountable mixture \Rightarrow no tractable EVI

we need to optimize the ELBO...

 \Rightarrow which is "tricky"





tractable bands
Normalizing flows

$$p_{\mathbf{X}}(\mathbf{x}) = p_{\mathbf{Z}}(f^{-1}(\mathbf{x})) \left| \det \left(\frac{\delta f^{-1}}{\delta \mathbf{x}} \right) \right|$$

an explicit likelihood!

tractable EVI queries!

many neural variants

RealNVP (Dinh et al. 2016), MAF (Papamakarios et al. 2017 MADE (Germain et al. 2015),

 \Rightarrow

PixelRNN (Oord et al. 2016)



Normalizing flows

$$p_{\mathbf{X}}(\mathbf{x}) = p_{\mathbf{Z}}(f^{-1}(\mathbf{x})) \left| \det \left(\frac{\delta f^{-1}}{\delta \mathbf{x}} \right) \right|$$



 \Rightarrow tractable EVI queries!



RealNVP (Dinh et al. 2016),

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MADE (Germain et al. 2015),

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q₁: What is the probability that today is a Monday at 12.00 and there is a traffic jam only on Westwood Blvd.?



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$$\mathbf{q}_1(\mathbf{m}) = p_{\mathbf{m}}(\mathsf{Day} = \mathsf{Mon}, \mathsf{Jam}_{\mathsf{Wwood}} = 1)$$



q₁: What is the probability that today is a Monday ot 12.00 and there is a traffic jam only on Westwood Blvd.?

$$\mathbf{q}_1(\mathbf{m}) = p_{\mathbf{m}}(\mathsf{Day} = \mathsf{Mon}, \mathsf{Jam}_{\mathsf{Wwood}} = 1)$$

General: $p_{\mathbf{m}}(\mathbf{e}) = \int p_{\mathbf{m}}(\mathbf{e}, \mathbf{H}) d\mathbf{H}$

where $\mathbf{E} \subset \mathbf{X}, \ \mathbf{H} = \mathbf{X} \setminus \mathbf{E}$



q₁: What is the probability that today is a Monday ot 12.00 and there is a traffic jam only on Westwood Blvd.?

$$\mathbf{q}_1(\mathbf{m}) = p_{\mathbf{m}}(\mathsf{Day} = \mathsf{Mon}, \mathsf{Jam}_{\mathsf{Wwood}} = 1)$$

tractable MAR \implies tractable **conditional queries** (CON):

$$p_{\mathbf{m}}(\mathbf{q} \mid \mathbf{e}) = \frac{p_{\mathbf{m}}(\mathbf{q}, \mathbf{e})}{p_{\mathbf{m}}(\mathbf{e})}$$



Tractable MAR : scene understanding





Fast and exact marginalization over unseen or "do not care" parts in the scene

Stelzner et al., "Faster Attend-Infer-Repeat with Tractable Probabilistic Models", 2019Kossen et al., "Structured Object-Aware Physics Prediction for Video Modeling and Planning", 201923/266

Normalizing flows

$$p_{\mathbf{X}}(\mathbf{x}) = p_{\mathbf{Z}}(f^{-1}(\mathbf{x})) \left| \det \left(\frac{\delta f^{-1}}{\delta \mathbf{x}} \right) \right|$$

an explicit likelihood!

 \implies tractable EVI queries!



Normalizing flows

$$p_{\mathbf{X}}(\mathbf{x}) = p_{\mathbf{Z}}(f^{-1}(\mathbf{x})) \left| \det \left(\frac{\delta f^{-1}}{\delta \mathbf{x}} \right) \right|$$

an explicit likelihood!

 \implies tractable EVI queries!

MAR is generally intractable: we cannot easily integrate over high-dimensional f





tractable bands

Probabilistic Graphical Models (PGMs)

Declarative semantics: a clean separation of modeling assumptions from inference

- Nodes: random variables
- Edges: dependencies



Inference:

conditioning (Darwiche 2001; Sang et al. 2005)
elimination (Zhang et al. 1994; Dechter 1998)
message passing (Yedidia et al. 2001; Dechter et al. 2002; Choi et al. 2010; Sontag et al. 2011)

Complexity of MAR on PGMs

Exact complexity: Computing MAR and CON is *#P-hard*

Approximation complexity: Computing MAR and CON approximately within a relative error of $2^{n^{1-\epsilon}}$ for any fixed ϵ is *NP-hard*

⇒ (Dagum et al. 1993; Roth 1996)

⇒ (Cooper 1990; Roth 1996)



Treewidth:

Informally, how tree-like is the graphical model **m**?

Fixed-parameter tractable: MAR and CON on a graphical model \mathbf{m} with treewidth w take time $O(|\mathbf{X}| \cdot 2^w)$ (Dechter 1998; Koller et al. 2009).

 \Rightarrow what about bounding the treewidth by design?

Low-treewidth PGMs



If treewidth is bounded (e.g. $\simeq 20$), exact MAR and CON inference is possible in practice

Tree distributions

A *tree-structured BN* (*Meilă et al. 2000*) where each $X_i \in \mathbf{X}$ has *at most* one parent Pa_{X_i} .



$$p(\mathbf{X}) = \prod_{i=1}^{n} p(x_i | \operatorname{Pa}_{x_i})$$

Exact querying: EVI, MAR, CON tasks *linear* for trees: $O(|\mathbf{X}|)$

Exact learning from d examples takes $O(|\mathbf{X}|^2 \cdot d)$ with the classical Chow-Liu algorithm¹

¹Chow et al., "Approximating discrete probability distributions with dependence trees", 1968 **31**,266



tractable bands



Expressiveness: Ability to represent rich and complex classes of distributions



Bounded-treewidth PGMs lose the ability to represent all possible distributions ...

Cohen et al., "On the expressive power of deep learning: A tensor analysis", 2016 Martens and Medabalimi, "On the Expressive Efficiency of Sum Product Networks", 2014



Mixtures as a convex combination of k (simpler) probabilistic models



$$p(X) = w_1 \cdot p_1(X) + w_2 \cdot p_2(X)$$

(77)

(77)

--->

EVI, MAR, CON queries scale linearly in \boldsymbol{k}



Mixtures as a convex combination of k (simpler) probabilistic models



$$p(X) = p(Z = 1) \cdot p_1(X|Z = 1)$$
$$+ p(Z = 2) \cdot p_2(X|Z = 2)$$

Mixtures are marginalizing a *categorical latent variable* Z with k values

 \Rightarrow increased expressiveness

Expressiveness and efficiency

Expressiveness: Ability to represent rich and effective classes of functions

 \Rightarrow mixture of Gaussians can approximate any probability density!

Cohen et al., "On the expressive power of deep learning: A tensor analysis", 2016 Martens and Medabalimi, "On the Expressive Efficiency of Sum Product Networks", 2014

Expressiveness and efficiency

Expressiveness: Ability to represent rich and effective classes of functions

⇒ mixture of Gaussians can approximate any probability density!

Expressive efficiency (aka Succinctness):

Ability to represent rich and effective classes of functions compactly

 \Rightarrow but how many components does a Gaussian mixture need?

Cohen et al., "On the expressive power of deep learning: A tensor analysis", 2016 Martens and Medabalimi, "On the Expressive Efficiency of Sum Product Networks", 2014



















solution: deep mixtures as in deep generative models



tractable bands

aka Most Probable Explanation (MPE)

q₅: Which combination of roads is most likely to be jammed on Monday at 9am?



aka Most Probable Explanation (MPE)

q₅: Which combination of roads is most likely to be jammed on Monday at 9am?

$$\mathbf{q}_5(\mathbf{m}) = \operatorname{argmax}_{\mathbf{j}} p_{\mathbf{m}}(\mathbf{j}_1, \mathbf{j}_2, \dots \mid \mathsf{Day} = \mathsf{M}, \mathsf{Time} = \mathbf{9})$$



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General: $\operatorname{argmax}_{\mathbf{q}} \, p_{\mathbf{m}}(\mathbf{q} \mid \mathbf{e})$

where $\mathbf{Q} \cup \mathbf{E} = \mathbf{X}$



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aka Most Probable Explanation (MPE)

q₅: Which combination of roads is most likely to be jammed on Monday at 9am?

...*intractable* for latent variable models!

$$\max_{\mathbf{q}} p_{\mathbf{m}}(\mathbf{q} \mid \mathbf{e}) = \max_{\mathbf{q}} \sum_{\mathbf{z}} p_{\mathbf{m}}(\mathbf{q}, \mathbf{z} \mid \mathbf{e})$$
$$\neq \sum_{\mathbf{z}} \max_{\mathbf{q}} p_{\mathbf{m}}(\mathbf{q}, \mathbf{z} \mid \mathbf{e})$$



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MAP inference : image inpainting



Predicting *arbitrary patches* given a *single* model without the need of retraining.

Poon and Domingos, "Sum-Product Networks: a New Deep Architecture", 2011 Sguerra and Cozman, "Image classification using sum-product networks for autonomous flight of micro aerial vehicles", 2016



tractable bands
aka Bayesian Network MAP

q₆: Which combination of roads is most likely to be jammed on Monday at 9am?



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General: $\operatorname{argmax}_{\mathbf{q}} p_{\mathbf{m}}(\mathbf{q} \mid \mathbf{e})$ = $\operatorname{argmax}_{\mathbf{q}} \sum_{\mathbf{h}} p_{\mathbf{m}}(\mathbf{q}, \mathbf{h} \mid \mathbf{e})$ where $\mathbf{Q} \cup \mathbf{H} \cup \mathbf{E} = \mathbf{X}$



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$$\mathbf{q}_6(\mathbf{m}) = \operatorname{argmax}_{\mathbf{j}} p_{\mathbf{m}}(\mathbf{j}_1, \mathbf{j}_2, \dots \mid \mathsf{Time} = \mathbf{9})$$

→ NP^{PP}-complete (Park et al. 2006)
⇒ NP-hard for trees (de Campos 2011)
⇒ NP-hard even for Naive Bayes (ibid.)



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tractable bands

q₂: Which day is most likely to have a traffic jam on my route to campus?



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q₂: Which day is most likely to have a traffic jam on my route to campus?

 $\mathbf{q}_{2}(\mathbf{m}) = \operatorname{argmax}_{\mathsf{d}} p_{\mathbf{m}}(\mathsf{Day} = \mathsf{d} \land \bigvee_{i \in \mathsf{route}} \mathsf{Jam}_{\mathsf{Str}\,i})$ $\implies marginals + MAP + logical events$



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q₂: Which day is most likely to have a traffic jam on my route to campus?

q₇: What is the probability of seeing more traffic jams in Westwood than Hollywood?



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Bekker et al., "Tractable Learning for Complex Probability Queries", 2015

q₂: Which day is most likely to have a traffic jam on my route to campus?

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 \Rightarrow counts + group comparison



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q₂: Which day is most likely to have a traffic jam on my route to campus?

q₇: What is the probability of seeing more traffic jams in Westwood than Hollywood?

 \mathbf{q}_8 : Is traffic more uncertain on weekdays?



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q₂: Which day is most likely to have a traffic jam on my route to campus?

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 \mathbf{q}_8 : Is traffic more uncertain on weekdays?

 \Rightarrow information-theoretic queries



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- **q**₂: Which day is most likely to have a traffic jam on my route to campus?
- **q**₇: What is the probability of seeing more traffic jams in Westwood than Hollywood?
- \mathbf{q}_8 : Is traffic more uncertain on weekdays?
- \mathbf{q}_9 : What is the causal effect of doing road works?



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- **q**₂: Which day is most likely to have a traffic jam on my route to campus?
- **q**₇: What is the probability of seeing more traffic jams in Westwood than Hollywood?
- **q**₈: Is traffic more uncertain on weekdays?
- \mathbf{q}_9 : What is the causal effect of doing road works?

 \Rightarrow causal backdoor estimation



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tractable bands



tractable bands



A completely disconnected graph. Example: Product of Bernoullis (PoBs)



Complete evidence, marginals and MAP, MMAP inference is *linear*!

 \Rightarrow but definitely not expressive...



tractable bands





Expressive models are not very tractable...



and tractable ones are not very expressive...



probabilistic circuits are at the "sweet spot"

Questions answered today

1. What are probabilistic queries? Are current models tractable? (Guy)

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Probabilistic Circuits



Given a reasoning task can we design a class of expressive models that is tractable for it?



Given a reasoning task can we design a class of deep computational graphs that is tractable for it?





Expressive models are not very tractable...



Tractable models are not that expressive...



Circuits can be both expressive and tractable!





then make it more expressive!



impose structure!

Input distributions

as computational nodes



Base case: a single node encoding a distribution

 \Rightarrow e.g., Gaussian PDF continuous random variable

Input distributions

as computational nodes



Base case: a single node encoding a distribution

 \Rightarrow e.g., indicators for X or $\neg X$ for Boolean random variable

Input distributions

as computational nodes



Simple distributions are tractable "black boxes" for:

- EVI: output $p(\mathbf{x})$ (density or mass)
- MAR: output 1 (normalized) or Z (unnormalized)
- MAP: output the mode
as computational graphs



$$p(X) = w_1 \cdot p_1(X_1) + w_2 \cdot p_2(X_1)$$

 \Rightarrow

translating inference to data structures...

as computational graphs



$$p(X_1) = 0.2 \cdot p_1(X_1) + 0.8 \cdot p_2(X_1)$$

 \Rightarrow ...e.g., as a weighted sum unit over Gaussian input distributions

as computational graphs



$$p(X = 5) = 0.2 \cdot p_1(X_1 = 5) + 0.8 \cdot p_2(X_1 = 5)$$

_

as computational graphs



A simplified notation:



scopes attached to inputs

edge directions omitted

as computational graphs

$$p(X_1, X_2, X_3) = p(X_1) \cdot p(X_2) \cdot p(X_3)$$



 \Rightarrow e.g. modeling a multivariate Gaussian with diagonal covariance matrix...

 \Rightarrow

as computational graphs

$$p(X_1, X_2, X_3) = p(X_1) \cdot p(X_2) \cdot p(X_3)$$



...with a product node over some univariate Gaussian distribution

as computational graphs

$$p(x_1, x_2, x_3) = p(x_1) \cdot p(x_2) \cdot p(x_3)$$





 \Rightarrow feedforward evaluation

as computational graphs

$$p(x_1, x_2, x_3) = p(x_1) \cdot p(x_2) \cdot p(x_3)$$





 \Rightarrow feedforward evaluation











Building PCs in Python with SPFlow



import spn.structure.leaves.parametric.Parametric as param
from param import Categorical, Gaussian

Molina et al., "SPFlow: An easy and extensible library for deep probabilistic learning using sum-product networks", 2019

EVI queries = **feedforward** evaluation

$$p(X_1 = -1.85, X_2 = 0.5, X_3 = -1.3, X_4 = 0.2)$$



EVI queries = **feedforward** evaluation

$$p(X_1 = -1.85, X_2 = 0.5, X_3 = -1.3, X_4 = 0.2)$$



EVI queries = **feedforward** evaluation

$$p(X_1 = -1.85, X_2 = 0.5, X_3 = -1.3, X_4 = 0.2) = 0.75$$



Just sum, products and distributions?



just arbitrarily compose them like a neural network!

Just sum, products and distributions?



just arbitrarily compose them like a neural-network!

 \Rightarrow structural properties needed for tractability

Which structural constraints ensure tractability?



A product node is decomposable if its children depend on disjoint sets of variables

 \implies just like in factorization!



decomposable circuit



non-decomposable circuit

Darwiche and Marquis, "A knowledge compilation map", 2002



aka completeness

A sum node is smooth if its children depend of the same variable sets

 \Rightarrow otherwise not accounting for some variables



Darwiche and Marquis, "A knowledge compilation map", 2002

Computing arbitrary integrations (or summations)

 \Rightarrow linear in circuit size!

E.g., suppose we want to compute Z (the distribution's normalizing constant):

$$\int oldsymbol{p}(\mathbf{x}) d\mathbf{x}$$

If $oldsymbol{p}(\mathbf{x}) = \sum_i w_i oldsymbol{p}_i(\mathbf{x})$, (smoothness):

$$\int \mathbf{p}(\mathbf{x}) d\mathbf{x} = \int \sum_{i} w_{i} \mathbf{p}_{i}(\mathbf{x}) d\mathbf{x} =$$
$$= \sum_{i} w_{i} \int \mathbf{p}_{i}(\mathbf{x}) d\mathbf{x}$$

 \Rightarrow

integrals are "pushed down" to children



If $m{p}(\mathbf{x},\mathbf{y},\mathbf{z})=m{p}(\mathbf{x})m{p}(\mathbf{y})m{p}(\mathbf{z})$, (decomposability):

$$\int \int \int \mathbf{p}(\mathbf{x}, \mathbf{y}, \mathbf{z}) d\mathbf{x} d\mathbf{y} d\mathbf{z} =$$
$$= \int \int \int \int \mathbf{p}(\mathbf{x}) \mathbf{p}(\mathbf{y}) \mathbf{p}(\mathbf{z}) d\mathbf{x} d\mathbf{y} d\mathbf{z} =$$
$$= \int \mathbf{p}(\mathbf{x}) d\mathbf{x} \int \mathbf{p}(\mathbf{y}) d\mathbf{y} \int \mathbf{p}(\mathbf{z}) d\mathbf{z}$$

 \Rightarrow integrals decompose into easier ones



Forward pass evaluation for MAR

 \Rightarrow linear in circuit size!

E.g. to compute $p(x_2, x_4)$:

leafs over X_1 and X_3 output $\boldsymbol{Z}_i = \int p(x_i) dx_i$

for normalized leaf distributions: 1.0

leafs over X_2 and X_4 output EVI

feedforward evaluation (bottom-up)



Forward pass evaluation for MAR

 \Rightarrow linear in circuit size!

E.g. to compute $p(x_2, x_4)$: leafs over X_1 and X_3 output $\mathbf{Z}_i = \int p(x_i) dx_i$ for normalized leaf distributions: 1.0 leafs over X_2 and X_4 output *EVI* feedforward evaluation (bottom-up)



Forward pass evaluation for MAR

 \Rightarrow linear in circuit size!

E.g. to compute $p(x_2, x_4)$: leafs over X_1 and X_3 output $\mathbf{Z}_i = \int p(x_i) dx_i$ \Rightarrow for normalized leaf distributions: 1.0 leafs over X_2 and X_4 output *EVI* feedforward evaluation (bottom-up)



Tractable MAR



Peharz et al., "Einsum Networks: Fast and Scalable Learning of Tractable Probabilistic Circuits", 2020

Analogously, for arbitrary conditional queries:

$$p(\mathbf{q} \mid \mathbf{e}) = \frac{p(\mathbf{q}, \mathbf{e})}{p(\mathbf{e})}$$

1. evaluate $p(\mathbf{q}, \mathbf{e}) \implies$ one feedforward pass2. evaluate $p(\mathbf{e}) \implies$ another feedforward pass \implies ...still linear in circuit size!



Tractable CON



Peharz et al., "Einsum Networks: Fast and Scalable Learning of Tractable Probabilistic Circuits", 2020

Generative models are still hard to control

more reasoning!

Generate image



Stable Diffusion

Generate a sentence using "frisbee", "caught" and "dog", following the given order.



Generate a sentence using "frisbee", "caught" and "dog", following the given order.



After a perfect throw, the <u>frisbee</u> glided through the air, and the <u>dog</u>, with incredible agility, <u>caught</u> it mid-flight.

ChatGPT



Generate a sentence using "frisbee", "caught" and "dog", following the given order.



After a perfect throw, the <u>frisbee</u> glided through the air, and the <u>dog</u>, with incredible agility, <u>caught</u> it mid-flight.

ChatGPT



That's not correct. Generate a sentence using "frisbee", "caught" and "dog". The keywords should appear in the order as specified.



Here's the correct sentence: The <u>dog caught</u> the <u>frisbee</u> in mid-air, showing off its amazing catching skills.

ChatGPT



Generate a sentence using "frisbee", "caught" and "dog", following the given order.



After a perfect throw, the frisbee glided through the air, and the **dog**, with incredible agility, **caught** it mid-flight.

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Here's the correct sentence: The dog caught the frisbee in mid-air, showing off its amazing catching skills.

ChatGPT



A <u>frisbee</u> is <u>caught</u> by a <u>dog</u>. A pair of <u>frisbee</u> players are <u>caught</u> in a <u>dog</u> fight.

Gel aTo
What do we have?

Prefix: "The weather is"

Constraint a: text contains "winter"

Model only does p(next-token|prefix) =

cold	0.05
warm	0.10

What do we need?

Prefix: "The weather is"

Constraint a: text contains "winter"

Generate from $p(\text{next-token}|\text{prefix}, \alpha) = \frac{\text{cold}}{\text{warm}} \frac{0.50}{0.01}$

$$\propto \sum_{ ext{text}} p(ext{next-token, text, prefix}, lpha)$$

Marginalization!

Computing $p(\alpha | x_{1:t+1})$

For α in conjunctive normal form (CNF):

$$(\mathsf{w}_{1,1} \lor \ldots \lor \mathsf{w}_{1,d1}) \land \ldots \land (\mathsf{w}_{m,1} \lor \ldots \lor \mathsf{w}_{m,dm})$$

where each w_{ij} is a keyword (i.e. a string of tokens), representing the constraint that w_{ij} appears in the generated text.

e.g., α = ("swims" V "like swimming") \land ("lake" V "pool")

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e.g., α = ("swims" V "like swimming") \land ("lake" V "pool")

Efficient algorithm:

For m clauses and sequence length n, time-complexity for generation is $O(2^{|m|}n)$ when p is a *hidden Markov model* (see general probabilistic circuit case later).

Trick: dynamic programming with clever preprocessing and local belief updates

CommonGen: a Challenging Benchmark

Given 3-5 concepts (keywords), our goal is to generate a sentence using all keywords, which can appear in any order and any form of inflections. e.g.,

Input: snow drive car

Reference 1: A car drives down a snow covered road.

Reference 2: Two cars drove through the snow.

$$(\mathsf{w}_{1,1} \lor \ldots \lor \mathsf{w}_{1,d1}) \land \ldots \land (\mathsf{w}_{m,1} \lor \ldots \lor \mathsf{w}_{m,dm})$$

Each clause represents the inflections for one keyword.

GeLaTo Overview

Lexical Constraint *a*: sentence contains keyword "winter" **Constrained Generation**: $Pr(x_{t+1} | \alpha, x_{1:t} =$ "the weather is") X intractable efficient Pre-trained Tractable Language Model Probabilistic Model Minimize KL-divergence $\Pr_{LM}(x_{t+1} | x_{1:t})$ $\Pr_{TPM}(\alpha | x_{t+1}, x_{1:t})$ x_{t+1} x_{t+1} 0.05 0.50 cold cold 0.10 0.01 warm warm

Honghua Zhang, Meihua Dang, Nanyun Peng and Guy Van den Broeck. Tractable Control for Autoregressive Language Generation, 2023.

GeLaTo Overview

Lexical Constraint *a*: sentence contains keyword "winter" **Constrained Generation**: $Pr(x_{t+1} | \alpha, x_{1:t} =$ "the weather is") X intractable efficient Pre-trained Tractable Language Model Probabilistic Model Minimize KL-divergence $\Pr_{LM}(x_{t+1} | x_{1:t})$ $\Pr_{TPM}(\alpha | x_{t+1}, x_{1:t})$ x_{t+1} x_{t+1} 0.05 0.50 cold cold 0.10 0.01 warm warm $p(x_{t+1} | \alpha, x_{1:t})$ x_{t+1} 0.025 cold 0.001 warm

Honghua Zhang, Meihua Dang, Nanyun Peng and Guy Van den Broeck. Tractable Control for Autoregressive Language Generation, 2023.

Step 2: Control p_{gpt} via p_{hmm}

Unsupervised

Language model is not fine-tuned/prompted to satisfy constraints

By Bayes rule: $p_{gpt}(x_{t+1} | x_{1:t}, \alpha) \propto p_{gpt}(\alpha | x_{1:t+1}) \cdot p_{gpt}(x_{t+1} | x_{1:t})$

Assume $p_{hmm}(\alpha | x_{1:t+1}) \approx p_{gpl}(\alpha | x_{1:t+1})$, we generate from:

 $p(x_{t+1} | x_{1:t}, \alpha) \propto p_{hmm}(\alpha | x_{1:t+1}) \cdot p_{gpt}(x_{t+1} | x_{1:t})$

Mathad		Generation Quality								Constraint Satisfaction			
Method	ROUGE-L		BLEU-4		CIDEr		SPICE		Coverage		Success Rate		
Unsupervised	dev	test	dev	test	dev	test	dev	test	dev	test	dev	test	
InsNet (Lu et al., 2022a)	-	-	18.7	-	1.0	100	-	-	100.0	Ξ.	100.0	-	
NeuroLogic (Lu et al., 2021)	-	41.9	-	24.7	-	14.4	-	27.5	-	96.7	-	-	
A*esque (Lu et al., 2022b)	-	44.3	-	28.6	-	15.6	-	29.6	-	97.1	-	-	
NADO (Meng et al., 2022)	-	-	26.2	-	-	-	-	-	96.1	-	-	-	
GeLaTo	44.6	44.1	29.9	29.4	16.0	15.8	31.3	31.0	100.0	100.0	100.0	100.0	

Honghua Zhang, Meihua Dang, Nanyun Peng and Guy Van den Broeck. Tractable Control for Autoregressive Language Generation, 2023.

Step 2: Control
$$p_{gpt}$$
 via p_{hmm}

<u>Supervised</u>

Language model is fine-tuned to perform constrained generation (e.g. seq2seq)

Empirically $p_{HMM}(\alpha | x_{1:t+1}) \approx p_{gpt}(\alpha | x_{1:t+1})$ does not hold well enough; we view $p_{HMM}(x_{t+1} | x_{1:t}, \alpha)$ and $p_{gpt}(x_{t+1} | x_{1:t})$ as classifiers trained for the same task with different biases; thus we generate from their <u>weighted</u> geometric mean:

 $p(x_{t+1} | x_{1:t}, \alpha) \propto p_{hmm}(x_{t+1} | x_{1:t}, \alpha)^{w} \cdot p_{gpt}(x_{t+1} | x_{1:t})^{1-w}$

Mathad		Generation Quality								Constraint Satisfaction			
Method	ROU	GE-L	L BLEU-4		CIDEr		SPICE		Coverage		Success Rate		
Supervised	dev	test	dev	test	dev	test	dev	test	dev	test	dev	test	
NeuroLogic (Lu et al., 2021)	-	42.8	-	26.7	12	14.7	-	30.5	-	97.7	-	93.9†	
A*esque (Lu et al., 2022b)	-	43.6	-	28.2	-	15.2	-	30.8	-	97.8	-	97.9 [†]	
NADO (Meng et al., 2022)	44.4	-	30.8	-	16.1^{\dagger}	-	32.0 [†]	-	97.1	-	88.8 [†]	-	
GeLaTo	46.0	45.6	34.1	32.9	16.7	16.8	31.3	31.9	100.0	100.0	100.0	100.0	

Advantages of GeLaTo:

- 1. Constraint α is <u>guaranteed to be satisfied</u>: for any next-token x_{t+1} that would make α unsatisfiable, $p(x_{t+1} | x_{1:t}, \alpha) = 0$ for both settings.
- 2. Training p_{hmm} does not depend on α , which is only imposed at inference (generation) time. Once p_{hmm} is trained, we can impose whatever α .
- 3. We can impose <u>additional tractable constraints</u>:
 - The keywords are generated following a particular order.
 - (Some) keywords must appear at a particular position.
 - (Some) keywords must not appear in the generated sentence.

Conclusion: you can control an intractable generative model using a tractable probabilistic circuit.



We can also decompose bottom-up a MAP query:

 $\max_{\mathbf{q}} p(\mathbf{q} \mid \mathbf{e})$



We *cannot* decompose bottom-up a MAP query:

 $\max_{\mathbf{q}} p(\mathbf{q} \mid \mathbf{e})$

since for a sum node we are marginalizing out a latent variable

$$\max_{\mathbf{q}} \sum_{i} w_{i} p_{i}(\mathbf{q}, \mathbf{e}) = \max_{\mathbf{q}} \sum_{\mathbf{z}} p(\mathbf{q}, \mathbf{z}, \mathbf{e}) \neq \sum_{\mathbf{z}} \max_{\mathbf{q}} p(\mathbf{q}, \mathbf{z}, \mathbf{e})$$

⇒ MAP for latent variable models is intractable (Conaty et al. 2017)



aka selectivity

A sum node is deterministic if the output of only one children is non zero for any input

 \Rightarrow e.g. if their distributions have disjoint support



deterministic circuit



non-deterministic circuit

Computing maximization with arbitrary evidence e

 \Rightarrow linear in circuit size!

E.g., suppose we want to compute:

$$\max_{\mathbf{q}} p(\mathbf{q} \mid \mathbf{e})$$



If
$$p(\mathbf{q}, \mathbf{e}) = \sum_{i} w_i p_i(\mathbf{q}, \mathbf{e}) = \max_i w_i p_i(\mathbf{q}, \mathbf{e})$$
,
(*deterministic* sum node):

$$\max_{\mathbf{q}} \mathbf{p}(\mathbf{q}, \mathbf{e}) = \max_{\mathbf{q}} \sum_{i} w_{i} \mathbf{p}_{i}(\mathbf{q}, \mathbf{e})$$
$$= \max_{\mathbf{q}} \max_{i} w_{i} \mathbf{p}_{i}(\mathbf{q}, \mathbf{e})$$
$$= \max_{i} \max_{\mathbf{q}} w_{i} \mathbf{p}_{i}(\mathbf{q}, \mathbf{e})$$



one non-zero child term, thus sum is max



If
$$p(q, e) = p(q_x, e_x, q_y, e_y) = p(q_x, e_x)p(q_y, e_y)$$

(*decomposable* product node):

$$\max_{\mathbf{q}} p(\mathbf{q} \mid \mathbf{e}) = \max_{\mathbf{q}} p(\mathbf{q}, \mathbf{e})$$
$$= \max_{\mathbf{q}_{\mathbf{x}}, \mathbf{q}_{\mathbf{y}}} p(\mathbf{q}_{\mathbf{x}}, \mathbf{e}_{\mathbf{x}}, \mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}})$$
$$= \max_{\mathbf{q}_{\mathbf{x}}} p(\mathbf{q}_{\mathbf{x}}, \mathbf{e}_{\mathbf{x}}) \cdot \max_{\mathbf{q}_{\mathbf{y}}} p(\mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}})$$
$$\implies \text{ solving optimization independently}$$



Evaluating the circuit twice: bottom-up and top-down

 \implies still linear in circuit size!



Evaluating the circuit twice: bottom-up and top-down

still linear in circuit size!

- 1. turn sum into max nodes and distributions into max distributions



Evaluating the circuit twice: bottom-up and top-down

still linear in circuit size!

- 1. turn sum into max nodes and distributions into max distributions
- 2. evaluate $p(x_2, x_4)$ bottom-up



Evaluating the circuit twice: bottom-up and top-down

still linear in circuit size!

- 1. turn sum into max nodes and distributions into max distributions
- 2. evaluate $p(x_2, x_4)$ bottom-up
- 3. retrieve max activations top-down





Evaluating the circuit twice: bottom-up and top-down

still linear in circuit size!

- 1. turn sum into max nodes and distributions into max distributions
- 2. evaluate $p(x_2, x_4)$ bottom-up
- 3. retrieve max activations top-down
- 4. compute **MAP states** for X_1 and X_3 at leaves



MAP inference : image segmentation



Semantic segmentation is MAP over joint pixel and label space

Even approximate MAP for non-deterministic circuits (SPNs) delivers good performances.

Rathke et al., "Locally adaptive probabilistic models for global segmentation of pathological oct scans", 2017 Yuan et al., "Modeling spatial layout for scene image understanding via a novel multiscale sum-product network", 2016 Friesen and Domingos, "Submodular Sum-product Networks for Scene Understanding", 2016

How expressive?

Dataset	Sparse PC (ours)	HCLT	RatSPN	IDF	BitSwap	BB-ANS	McBits
MNIST	1.14	1.20	1.67	1.90	1.27	1.39	1.98
EMNIST(MNIST)	1.52	1.77	2.56	2.07	1.88	2.04	2.19
EMNIST(Letters)	1.58	1.80	2.73	1.95	1.84	2.26	3.12
EMNIST(Balanced)	1.60	1.82	2.78	2.15	1.96	2.23	2.88
EMNIST(ByClass)	1.54	1.85	2.72	1.98	1.87	2.23	3.14
FashionMNIST	3.27	3.34	4.29	3.47	3.28	3.66	3.72

competitive with Flows and VAEs!

Dang et al., "Sparse Probabilistic Circuits via Pruning and Growing", 2022

How scalable?

Dataset		TPM	DGMs				
	LVD (ours)	HCLT	EiNet	RAT-SPN	Glow	RealNVP	BIVA
ImageNet32	4.39±0.01	4.82	5.63	6.90	4.09	4.28	3.96
ImageNet64	4.12±0.00	4.67	5.69	6.82	3.81	3.98	-
CIFAR	4.38±0.02	4.61	5.81	6.95	3.35	3.49	3.08



up to billions of parameters

Liu et al., "Scaling Up Probabilistic Circuits by Latent Variable Distillation", 2022

Questions answered today

- 1. What are probabilistic queries? Are current models tractable? (Guy)
- 2. What are probabilistic circuits and why are they tractable? (Guy)
- What is the connection to logical circuit languages? (YooJung)
- 4. How do I compile my favorite model into a circuit? (YooJung)
- 5. How are circuit size and tractability related? (YooJung)
- 6. What's the most impressive query we can efficiently compute? (YooJung)
- 7. Are all tractable distributions probabilistic circuits? (Guy)
- 8. *How to learn probabilistic circuits from data?* (Guy)

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Logical Circuits

Tractability to other semi-rings

Tractable probabilistic inference exploits *efficient summation for decomposable functions* in the probability commutative semiring:

 $(\mathbb{R}, +, \times, 0, 1)$

analogously efficient computations can be done in other semi-rings:

 $(\mathbb{S},\oplus,\otimes,0_\oplus,1_\otimes)$



Algebraic model counting (Kimmig et al. 2017), Semi-ring

programming (Belle et al. 2016)

Historically, very well studied for boolean functions:

$$(\mathbb{B} = \{0, 1\}, \lor, \land, 0, 1) \implies \text{logical circuits!}$$

Logical circuits







s/d-D/NNFs (Darwiche et al. 2002a)

O/BDDs (Bryant 1986)

SDDs (Darwiche 2011a)

Logical circuits are compact representations for boolean functions...



structural properties

...and like probabilitistic circuits, one can define **structural properties**: (structured) decomposability, smoothness, determinism allowing for tractable computations



Darwiche and Marquis, "A knowledge compilation map", 2002



a knowledge compilation map

...inducing *a hierarchy of tractable logical circuit families*



Darwiche and Marquis, "A knowledge compilation map", 2002

Knowledge Compilation





Decomposability (DNNF)



Determinism (d-DNNF)





Model counting problem: given a Boolean formula Δ , compute the number of satisfying assignments.

Weighted model counting (WMC):

$$WMC(\Delta, w) = \sum_{\mathbf{x} \models \Delta} \prod_{l \in \mathbf{x}} w(l)$$

$$\Rightarrow$$
 linear in circuit size!
Decomposability + determinism = tractable (W)MC

To compute $\mathrm{WMC}(\Delta, w)$:

Turn OR gates to sum nodes and AND gates to product nodes

Replace each literal l with its weight w(l)

bottom-up evaluation



Decomposability + determinism = tractable (W)MC

To compute $WMC(\Delta, w)$:

Turn OR gates to sum nodes and AND gates to product nodes

Replace each literal l with its weight w(l)

bottom-up evaluation



Probabilistic inference by WMC

connection to probabilistic circuits through WMC

- 1. Encode probabilistic model as WMC formula (Δ,w)
- 2. Compile Δ into a logical circuit (e.g. d-DNNF, OBDD, SDD, etc.)
- 3. Tractable MAR/CON by tractable WMC on circuit
- 4. Answer complex queries tractably by enforcing more structural properties!

Probabilistic inference by WMC

connection to probabilistic circuits through WMC

Resulting compiled WMC circuit equivalent to probabilistic circuit

 \Rightarrow parameter variables o edge parameters



Compiled circuit of WMC encoding

Equivalent probabilistic circuit

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via compilation

Bottom-up *compilation*: starting from leaves...



via compilation

...compile a leaf CPT



p(A|C=0)



via compilation

...compile a leaf CPT







via compilation

...compile a leaf CPT...for all leaves...





via compilation

...and recurse over parents...





via compilation

...while reusing previously compiled nodes!...



via compilation





Hidden Markov Models

as computational graphs





Compilation : probabilistic programming



Chavira et al., "Compiling relational Bayesian networks for exact inference", 2006 Holtzen et al., "Symbolic Exact Inference for Discrete Probabilistic Programs", 2019 De Raedt et al.; Riguzzi; Fierens et al.; Vlasselaer et al., "ProbLog: A Probabilistic Prolog and Its Application in Link Discovery."; "A top down interpreter for LPAD and CP-logic"; "Inference and Learning in Probabilistic Logic Programs using Weighted Boolean Formulas"; "Anytime Inference in Probabilistic Logic Programs with Tp-compilation", 2007; 2007; 2015; 2015 Olteanu et al.; Van den Broeck et al., "Using OBDDs for efficient query evaluation on probabilistic databases"; <u>Query Processing on Probabilistic Data: A Survey</u>, 2008; 2017 Vlasselaer et al., "Exploiting Local and Repeated Structure in Dynamic Bayesian Networks", 2016

Decision Diagrams

FBDDs (Free binary decision diagrams; *read-once*)

OBDDs (Ordered BDDs)

SDDs (Sentential decision diagrams)



 \Rightarrow BDD as circuit

Darwiche and Marquis, "A knowledge compilation map", 2002







Partitioned Determinism (SDDs)

Darwiche, IJCAI 2011



Partitioned Determinism (SDDs)

Darwiche, IJCAI 2011



Decision Diagrams



Darwiche, "SDD: A new canonical representation of propositional knowledge bases", 2011

Probability of logical events

q₈: What is the probability of having a traffic jam on my route to campus?



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Probability of logical events

q₈: What is the probability of having a traffic jam on my route to campus?

 $\mathbf{q}_8(\mathbf{m}) = p_{\mathbf{m}}(\bigvee_{i \in \mathsf{route}} \operatorname{\mathsf{Jam}}_{\mathsf{Str}\,i})$

⇒ marginals + logical events



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Computing $\boldsymbol{p}(\alpha)$: the probability of arbitrary logical formula

Multilinear in circuit sizes if the logical circuit:

is smooth, structured decomposable, deterministic

shares the same vtree





If
$$p(\mathbf{x}) = \sum_{i} w_{i} p_{i}(\mathbf{x}), \boldsymbol{\alpha} = \bigvee_{j} \boldsymbol{\alpha}_{j},$$

(smooth p)
(smooth + deterministic $\boldsymbol{\alpha}$):

$$p(\boldsymbol{\alpha}) = \sum_{i} w_{i} p_{i} \left(\bigvee_{j} \boldsymbol{\alpha}_{j}\right) = \sum_{i} w_{i} \sum_{j} p_{i} \left(\boldsymbol{\alpha}_{j}\right) \overset{(\mathbf{\alpha})}{\underset{X_{2}}{\overset{(\mathbf{\alpha})}{\underset{X_{1}}{\overset{(\mathbf{\alpha})}{\overset{(\mathbf{\alpha})}{\underset{X_{2}}{\overset{(\mathbf{\alpha})}{\underset{X_{1}}{\overset{(\mathbf{\alpha})}{\underset{X_{2}}{\overset{(\mathbf{\alpha})}{\underset{X_{1}}{\overset{(\mathbf{\alpha})}{\underset{X_{2}}{\overset{(\mathbf{\alpha})}{\underset{X_{1}}{\overset{(\mathbf{\alpha})}{\underset{X_{2}}{\overset{(\mathbf{\alpha})}{\underset{X_{1}}{\overset{(\mathbf{\alpha})}{\underset{X_{2}}{\overset{(\mathbf{\alpha})}{\underset{X_{2}}{\overset{(\mathbf{\alpha})}{\underset{X_{2}}{\underset{X_{1}}{\overset{(\mathbf{\alpha})}{\underset{X_{2}}{\underset{X_{2}}{\underset{X_{1}}{\overset{(\mathbf{\alpha})}{\underset{X_{2}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X_{2}}{\underset{X_{2}}{\underset{X_{1}}{\underset{X_{2}}{\underset{X$$

probabilities are "pushed down" to children

If $p(\mathbf{x},\mathbf{y}) = p(\mathbf{x})p(\mathbf{y})$, $\boldsymbol{\alpha} = \boldsymbol{\beta} \wedge \gamma$, (structured decomposability):

$$p(\alpha) = p(\beta \wedge \gamma) \cdot p(\beta \wedge \gamma) = p(\beta) \cdot p(\gamma)$$



probabilities decompose into simpler ones





To compute $p(\alpha)$:

compute the probability for each **pair** of probabilistic and logical circuit nodes for the **same vtree node**

cache the values!

eedforward evaluation (bottom-up)





To compute $p(\alpha)$:

compute the probability for each **pair** of probabilistic and logical circuit nodes for the **same vtree node**

cache the values!

feedforward evaluation (bottom-up)





structured decomposability = tractable...

Symmetric and **group queries** (exactly-k, odd-number, etc.) (Bekker et al. 2015)

For the "right" vtree

- Marginal MAP (Oztok et al. 2016)
- Probability of logical circuit event in probabilistic circuit (Choi et al. 2015b)
- Multiply two probabilistic circuits (Shen et al. 2016)
- KL Divergence between probabilistic circuits (Liang et al. 2017)
- Same-decision probability (Oztok et al. 2016)
- Expected same-decision probability (Choi et al. 2017)
- Expected classifier agreement (Choi et al. 2018)
- Expected predictions (Khosravi et al. 2019b)

Questions answered today

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Succinctness of circuits

Expressive efficiency

Tractability is defined with respect to the size of the model.

How do structural constraints affect **expressive efficiency** (**succinctness**) of probabilistic/logical circuits?



Succinctness of circuits

Expressive efficiency

A family of circuits \mathcal{M}_1 is **at least as succinct as** \mathcal{M}_2 iff for every $\mathbf{m}_2 \in \mathcal{M}_2$, there exists $\mathbf{m}_1 \in \mathcal{M}_1$ that represents the same function and $|m_1| \leq |\mathsf{poly}(m_2)|$.

Succinctness of circuits

Expressive efficiency

Strict succinctness ordering: DNNF < d-DNNF < FBDD < OBDD



Darwiche and Marquis, "A knowledge compilation map", 2002
Succinctness of circuits

Expressive efficiency

Strict succinctness ordering: **DNNF < d-DNNF** < FBDD < OBDD



The Sauerhoff function has DNNF of size $O(n^2)$ but d-DNNF of size $2^{\Omega(n)}$ (Bova et al. 2016).

Succinctness of circuits

Expressive efficiency

Strict succinctness ordering: **DNNF < d-DNNF** < FBDD < OBDD

d-DNNF ≰ DNNF unless the polynomial hierarchy collapses (*Darwiche et al. 2002a*).

The Sauerhoff function has DNNF of size $O(n^2)$ but d-DNNF of size $2^{\Omega(n)}$ (Bova et al. 2016).



Unconditional exponential separation for d-DNNF $\not\leq$ DNNF

 Using a connection between circuits and communication complexity

Succinctness of circuits

Expressive efficiency

SDD < **OBDD**: *SDDs are strictly more succinct than OBDDs*

SDD \leq OBDD: OBDDs are SDDs with right-linear vtrees

SDD \geq OBDD: The *hidden weighted bit function* has SDD of size $O(n^3)$ but OBDD of size $2^{\Omega}(n)$.



Bova, "SDDs are exponentially more succinct than OBDDs", 2016





How precise is the characterization of tractable circuits by structural properties?

Smoothness + decomposability = tractable MAR

Recall: Smoothness and decomposability allow marginal inference by feedforward (sum-product) evaluation.







Smoothness + decomposability = tractable MAR

Recall: Smoothness and decomposability allow marginal inference by feedforward (sum-product) evaluation.

Are these properties necessary?



Smoothness + decomposability = tractable MAR

Recall: Smoothness and decomposability allow marginal inference by feedforward (sum-product) evaluation.

 \Rightarrow Yes! Otherwise, integrals do not decompose.

Are these properties necessary?







Determinism + decomposability = tractable MAP

Recall: Determinism and decomposability allow MAP inference by feedforward (max-product) evaluation.







Determinism + decomposability = tractable MAP

Recall: Determinism and decomposability allow MAP inference by feedforward (max-product) evaluation.

However, decomposability is not necessary!



Determinism + decomposability = tractable MAP

Recall: Determinism and decomposability allow MAP inference by feedforward (max-product) evaluation.













A product node is consistent if any variable shared between its children appears in a single leaf node

 \Rightarrow decomposability implies consistency



consistent circuit



inconsistent circuit

Determinism + consistency = tractable MAP

Determinism + **consistency** = **tractable MAP**

If
$$\max_{\mathbf{q}_{\mathsf{shared}}} \frac{p(\mathbf{q}, \mathbf{e})}{p(\mathbf{q}_{\mathbf{x}}, \mathbf{e}_{\mathbf{x}})} \cdot \max_{\mathbf{q}_{\mathsf{shared}}} \frac{p(\mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}})}{p(\mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}})}$$
 (consistent):

$$\begin{aligned} \max_{\mathbf{q}} p(\mathbf{q}, \mathbf{e}) &= \max_{\mathbf{q}_{\mathbf{x}}, \mathbf{q}_{\mathbf{y}}} p(\mathbf{q}_{\mathbf{x}}, \mathbf{e}_{\mathbf{x}}, \mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}}) \\ &= \max_{\mathbf{q}_{\mathbf{x}}} p(\mathbf{q}_{\mathbf{x}}, \mathbf{e}_{\mathbf{x}}) \cdot \max_{\mathbf{q}_{\mathbf{y}}} p(\mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}}) \end{aligned}$$

 \Rightarrow solving optimization independently





Are smooth & decomposable circuits as succinct as deterministic & consistent ones, or vice versa?



Smooth & decomposable circuits strictly more succinct than deterministic & decomposable ones

Smooth & consistent circuits are equally succinct as smooth & decomposable ones



Smooth & decomposable circuits strictly more succinct than deterministic & decomposable ones Smooth & consistent circuits are equally succinct as smooth & decomposable ones



Smooth & decomposable circuits strictly more succinct than deterministic & decomposable ones Smooth & consistent circuits are equally succinct as smooth & decomposable ones



Consider following circuit over Boolean variables: $\prod_{i}^{r} (Y_i \cdot Z_{i1} + (\neg Y_i) \cdot Z_{i2}), \quad Z_{ij} \in \mathbf{X}$

Size linear in the number of variables

Deterministic and consistent

Marginal (with no evidence) is the solution to #P-hard SAT' problem (*Valiant 1979*) \Rightarrow **no tractable circuit for marginals!**



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Consider the marginal distribution $p(\mathbf{X})$ from a naive Bayes distribution $p(\mathbf{X}, C)$:

Linear-size smooth and decomposable circuit

MAP of $p(\mathbf{X})$ solves marginal MAP of $p(\mathbf{X}, C)$ which is NP-hard (*de Campos 2011*) \Rightarrow no tractable circuit for MAP!



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MAP of $p(\mathbf{X})$ solves marginal MAP of $p(\mathbf{X}, C)$ which is NP-hard (*de Campos 2011*) \Rightarrow no tractable circuit for MAP!



Neither smooth & decomposable nor deterministic & consistent circuits are more succinct than the other!

> Choose tractable circuit family based on your query

More theoretical questions remaining *"Complete the m*

 \Rightarrow



Neither smooth & decomposable nor deterministic & consistent circuits are more succinct than the other!

⇒ Choose tractable circuit family based on your query

More theoretical questions remaining

"Complete the map"



Succinctness map for *monotone* circuits

(s)mooth, (d)eterministic, (D)ecomposable, (w)eak (D)ecomposable (i.e. consistent)

Colnet and Mengel, "A Compilation of Succinctness Results for Arithmetic Circuits", 2021



Succinctness map for monotone circuits



Succinctness map for *positive* circuits (non-negative output, but weights may be negative)



(s)mooth, (d)eterministic, (D)ecomposable, (w)eak (D)ecomposable (i.e. consistent)

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Given a class of queries can we systematically find a class of probabilistic circuits that is tractable for it?



Integral expressions that can be formed by composing these operators

+ , \times , pow , log , exp and /



 \Rightarrow many divergences and information-theoretic queries

A language for queries

Integral expressions that can be formed by composing these operators

+ , imes , pow , log , exp and /

 \Rightarrow many divergences and information-theoretic queries

Represented as *higher-order computational graphs*—pipelines—operating over circuits! *re-using intermediate transformations across queries*

$\mathbb{KLD}(p \mid\mid q) = \int_{\mathsf{val}(\mathbf{X})} p(\mathbf{x}) \times \log\left(p(\mathbf{x})/q(\mathbf{x})\right) \, d\mathbf{X}$



$$\mathbb{KLD}(p \mid\mid q) = \int_{\mathsf{val}(\mathbf{X})} p(\mathbf{x}) \times \log \left(p(\mathbf{x}) / q(\mathbf{x}) \right) \, d\mathbf{X}$$



$\mathbb{KLD}(p \mid\mid q) = \int_{\mathsf{val}(\mathbf{X})} p(\mathbf{x}) \times \log (p(\mathbf{x})/q(\mathbf{x})) \ d\mathbf{X}$



$\mathbb{KLD}(p \mid\mid q) = \int_{\mathsf{val}(\mathbf{X})} p(\mathbf{x}) \times \log\left(p(\mathbf{x})/q(\mathbf{x})\right) \, d\mathbf{X}$



$\mathbb{XENT}(p \mid\mid q) = \int p(\mathbf{x}) \times \log q(\mathbf{x}) \, d\mathbf{X}$


$$\mathbb{E}_{\mathbf{x}^m \sim p(\mathbf{x}^m | \mathbf{x}^o)} \left[q^{lpha}(\mathbf{x}^m, \mathbf{x}^o) \right]$$





Two circuits are *compatible* if they have the same *hierarchical scope partitioning*

 \Rightarrow generalizes "structured decomposability with same vtree"





non-compatible circuits

Tractable operators





smooth, decomposable compatible

Tractable operators



smooth, decomposable deterministic

smooth, decomposable



Building an atlas of composable tractable atomic operations



To perform tractable integration we need *s* to be *smooth and decomposable*...



hence we need p and r to be smooth, decomposable and $\emph{compatible}$...



therefore *q* must be smooth, decomposable and *deterministic*...



we can compute \mathbb{XENT} tractably if p and q are smooth, decomposable, compatible and q is deterministic...

	Query	Tract. Conditions	Hardness
CROSS ENTROPY	$-\int p(oldsymbol{x})\log q(oldsymbol{x})\mathrm{d}\mathbf{X}$	Cmp, q Det	#P-hard w/o Det
SHANNON ENTROPY	$-\sum p(oldsymbol{x})\log p(oldsymbol{x})$	Sm, Dec, Det	coNP-hard w/o Det
Rényi Entropy	$(1-\alpha)^{-1}\log \int p^{\alpha}(\boldsymbol{x}) d\mathbf{X}, \alpha \in \mathbb{N}$	SD	#P-hard w/o SD
	$(1-\alpha)^{-1}\log\int p^{\alpha}(\boldsymbol{x}) d\mathbf{X}, \alpha \in \mathbb{R}_+$	Sm, Dec, Det	#P-hard w/o Det
MUTUAL INFORMATION	$\int p(oldsymbol{x},oldsymbol{y}) \log(p(oldsymbol{x},oldsymbol{y})/(p(oldsymbol{x})p(oldsymbol{y})))$	Sm, SD, Det*	coNP-hard w/o SD
KULLBACK-LEIBLER DIV.	$\int p(oldsymbol{x}) \log(p(oldsymbol{x})/q(oldsymbol{x})) doldsymbol{X}$	Cmp, Det	#P-hard w/o Det
PÉNVI'S ALDHA DIV	$(1-\alpha)^{-1}\log \int p^{\alpha}(\boldsymbol{x})q^{1-\alpha}(\boldsymbol{x}) d\mathbf{X}, \alpha \in \mathbb{N}$	Cmp, q Det	#P-hard w/o Det
KENTI S ALPHA DIV.	$(1-\alpha)^{-1}\log \int p^{\alpha}(\boldsymbol{x})q^{1-\alpha}(\boldsymbol{x}) d\mathbf{X}, \alpha \in \mathbb{R}$	Cmp, Det	#P-hard w/o Det
ITAKURA-SAITO DIV.	$\int [p(oldsymbol{x})/q(oldsymbol{x}) - \log(p(oldsymbol{x})/q(oldsymbol{x})) - 1] d \mathbf{X}$	Cmp, Det	#P-hard w/o Det
CAUCHY-SCHWARZ DIV.	$-\lograc{\int p(oldsymbol{x})q(oldsymbol{x})doldsymbol{X}}{\sqrt{\int p^2(oldsymbol{x})doldsymbol{X}\int q^2(oldsymbol{x})doldsymbol{X}}}$	Cmp	#P-hard w/o Cmp
SQUARED LOSS	$\int (p(oldsymbol{x}) - q(oldsymbol{x}))^2 d \mathbf{X}$	Cmp	#P-hard w/o Cmp

compositionally derive the tractability of many more queries

Vergari et al., "A Compositional Atlas of Tractable Circuit Operations: From Simple Transformations to Complex Information-Theoretic Queries", 2021

	Query	Tract. Conditions	Hardness
CROSS ENTROPY	$-\int p(oldsymbol{x})\log q(oldsymbol{x})\mathrm{d}\mathbf{X}$	Cmp, q Det	#P-hard w/o Det
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Rényi Entropy	$(1-\alpha)^{-1}\log \int p^{\alpha}(\boldsymbol{x}) d\mathbf{X}, \alpha \in \mathbb{N}$	SD	#P-hard w/o SD
	$(1-\alpha)^{-1}\log\int p^{\alpha}(\boldsymbol{x}) d\mathbf{X}, \alpha \in \mathbb{R}_+$	Sm, Dec, Det	#P-hard w/o Det
MUTUAL INFORMATION	$\int p(oldsymbol{x},oldsymbol{y}) \log(p(oldsymbol{x},oldsymbol{y})/(p(oldsymbol{x})p(oldsymbol{y})))$	Sm, SD, Det*	coNP-hard w/o SD
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KENTI S ALPHA DIV.	$(1-\alpha)^{-1}\log \int p^{\alpha}(\boldsymbol{x})q^{1-\alpha}(\boldsymbol{x}) d\mathbf{X}, \alpha \in \mathbb{R}$	Cmp, Det	#P-hard w/o Det
ITAKURA-SAITO DIV.	$\int [p(oldsymbol{x})/q(oldsymbol{x}) - \log(p(oldsymbol{x})/q(oldsymbol{x})) - 1] d \mathbf{X}$	Cmp, Det	#P-hard w/o Det
CAUCHY-SCHWARZ DIV.	$-\lograc{\int p(oldsymbol{x})q(oldsymbol{x})doldsymbol{X}}{\sqrt{\int p^2(oldsymbol{x})doldsymbol{X}\int q^2(oldsymbol{x})doldsymbol{X}}}$	Cmp	#P-hard w/o Cmp
SQUARED LOSS	$\int (p(oldsymbol{x}) - q(oldsymbol{x}))^2 d \mathbf{X}$	Cmp	#P-hard w/o Cmp

and **prove hardness** when some input properties are not satisfied

Vergari et al., "A Compositional Atlas of Tractable Circuit Operations: From Simple Transformations to Complex Information-Theoretic Queries", 2021

Composable tractable sub-routines

```
function kld(p, q)
                           function xent(p, q)
                               r = log(q)
    r = quotient(p, q)
                               s = product(p, r)
    s = log(r)
                               return -integrate(s)
    t = product(p, s)
    return integrate(t)
                           end
end
                           function alphadiv(p, q, alpha=1.5)
function ent(p)
                               r = real_pow(p, alpha)
    q = log(p)
                               s = real_pow(q, 1.0-alpha)
    r = product(p, q)
                               t = product(r, s)
    return -integrate(s)
                               return log(integrate(t)) / (1.0-alpha)
end
                           end
```

Efficient inference algorithms in a couple lines of Julia code!²

²https://github.com/UCLA-StarAI/circuit-ops-atlas



- 1. Learning and reasoning with symbolic constraints
- 2. Expected predictions: handling missing values, fairness
- 3. Exact inference of causal effects

 \Rightarrow using tractable operators



Symbolic constraints

"How can neural nets reason and learn with symbolic constraints reliably and efficiently?"





Ground Truth

e.g. predict shortest path in a map





given x // e.g. a tile map

Ground Truth

Vlastelica et al., "Differentiation of blackbox combinatorial solvers", 2020





given \mathbf{x} // e.g. a tile map find $\mathbf{y}^* = \operatorname{argmax}_{\mathbf{y}} p_{\theta}(\mathbf{y} \mid \mathbf{x})$ // e.g. a configurations of edges in a grid

Ground Truth

Vlastelica et al., "Differentiation of blackbox combinatorial solvers", 2020





given $\mathbf{x} \quad // e.g.$ a tile map find $\mathbf{y}^* = \operatorname{argmax}_{\mathbf{y}} p_{\theta}(\mathbf{y} \mid \mathbf{x}) \quad // e.g.$ a configurations of edges in a grid s.t. $\mathbf{y} \models \mathsf{K} \quad // e.g.$, that form a valid path

Ground Truth

Vlastelica et al., "Differentiation of blackbox combinatorial solvers", 2020

When?



given $\mathbf{x} \quad // e.g.$ a tile map find $\mathbf{y}^* = \operatorname{argmax}_{\mathbf{y}} p_{\theta}(\mathbf{y} \mid \mathbf{x}) \quad // e.g.$ a configurations of edges in a grid s.t. $\mathbf{y} \models \mathsf{K} \quad // e.g.$, that form a valid path

// for a 12×12 grid, 2^{144} states but only 10^{10} valid ones!

Ground Truth

Vlastelica et al., "Differentiation of blackbox combinatorial solvers", 2020

When?



given \mathbf{x} // e.g. a feature map find $\mathbf{y}^* = \operatorname{argmax}_{\mathbf{y}} p_{\theta}(\mathbf{y} \mid \mathbf{x})$ // e.g. labels of classes s.t. $\mathbf{y} \models \mathsf{K}$ // e.g., constraints over superclasses

$$\mathsf{K}: (Y_{\mathsf{cat}} \implies Y_{\mathsf{animal}}) \land (Y_{\mathsf{dog}} \implies Y_{\mathsf{animal}})$$

hierarchical multi-label classification





Ground Truth

ResNet-18

neural nets struggle to satisfy domain constraints!





take an unreliable neural network architecture...





.....and replace the last layer with a semantic probabilistic layer





 $m{q}_{m{\Theta}}(\mathbf{y} \mid g(\mathbf{z}))$ is an expressive distribution over labels

 $c_{\mathsf{K}}(\mathbf{x},\mathbf{y})$ encodes the constraint $\mathbbm{1}\{\mathbf{x},\mathbf{y}\models\mathsf{K}\}$

Ahmed et al., "Semantic Probabilistic Layers for Neuro-Symbolic Learning", 2022









a conditional circuit $q(\mathbf{y}; \boldsymbol{\Theta} = g(\mathbf{z}))$





and a logical circuit $\boldsymbol{c}(\mathbf{y},\mathbf{x})$ encoding K

Tractable products





smooth, decomposable compatible

exactly compute \boldsymbol{Z} in time $O(|\boldsymbol{q}||\boldsymbol{c}|)$



$$\mathsf{K}: \, (Y_1 = 1 \implies Y_3 = 1) \\ \land \quad (Y_2 = 1 \implies Y_3 = 1)$$

1) Take any logical constraint

SPL recipe

$$\begin{split} \mathsf{K}:\, (Y_1=1 \implies Y_3=1) \\ \wedge \quad (Y_2=1 \implies Y_3=1) \end{split}$$

$$\begin{array}{c} 1 \left(Y_{5}=1\right) \bigodot \ref{eq: starting starting$$

1) Take any logical constraint

2) Compile it into a constraint circuit

SPL recipe

$$\mathsf{K} : (Y_1 = 1 \implies Y_3 = 1)$$

$$\land \quad (Y_2 = 1 \implies Y_3 = 1)$$





1) Take any logical constraint

2) Compile it into a constraint circuit

3) Multiply it by a circuit distribution

SPL recipe

$$\mathsf{K} : (Y_1 = 1 \implies Y_3 = 1)$$

$$\land \quad (Y_2 = 1 \implies Y_3 = 1)$$





1) Take any logical constraint

2) Compile it into a constraint circuit

3) Multiply it by a circuit distribution

4) train end-to-end by sgd!

Guaranteeing consistency

Ground Truth



cost: 39.31



cost: ∞



cost: ∞



SPL

cost: 45.09



cost: 57.31



cost: ∞



cost: ∞



cost: 58.09

Expected predictions

Reasoning about the output of a classifier or regressor $m{f}$ given a distribution $m{p}$ over the input features

$$\mathbb{E}_{p}[f] = \int_{\mathsf{val}(\mathbf{X})} p(\mathbf{x}) \times f(\mathbf{x}) \, d\mathbf{X}$$

$$p \longrightarrow f$$

$$r$$

$$f$$

Handling missing values at test time



Given a partial observation \mathbf{x}^{o} , what is the expected output from f?

$$\mathop{\mathbb{E}}_{\mathbf{x}^m \sim p(\mathbf{x}^m | \mathbf{x}^o)} \left[f(\mathbf{x}^m, \mathbf{x}^o) \right]$$

Khosravi et al., "On Tractable Computation of Expected Predictions", 2019

Fairness analysis



using ProbabilisticCircuits
pc = load_prob_circuit(zoo_psdd_file("insurance.psdd"));
rc = load_logistic_circuit(zoo_lc_file("insurance.circuit"), 1);

```
q: Is the predictive model biased by gender?
```

```
groups = make_observations([["male"], ["female"]])
exps, _ = Expectation(pc, rc, groups);
println("Female : \$ $(exps[2])");
println("Male : \$ $(exps[1])");
println("Diff : \$ $(exps[2] - exps[1])");
Female : $ 14170.125469335406
Male : $ 13196.548926381849
Diff : $ 973.5765429535568
```
Given subsets $\boldsymbol{A}, \boldsymbol{Y} \subseteq \boldsymbol{X}$, interested in *causal effect* $p(\boldsymbol{Y}|do(\boldsymbol{A}))$.

Given subsets $A, Y \subseteq X$, interested in *causal effect* p(Y|do(A)). In general, $p(Y|do(A)) \neq p(Y|A)$ (correlation is not causation).

Given subsets $\mathbf{A}, \mathbf{Y} \subseteq \mathbf{X}$, interested in *causal effect* $p(\mathbf{Y}|do(\mathbf{A}))$. In general, $p(\mathbf{Y}|do(\mathbf{A})) \neq p(\mathbf{Y}|\mathbf{A})$ (correlation is not causation).

Specify (qualitative) assumptions on the system using a causal diagram G (here A, Y, Z, K ⊆ X)):



Given subsets $\mathbf{A}, \mathbf{Y} \subseteq \mathbf{X}$, interested in *causal effect* $p(\mathbf{Y}|do(\mathbf{A}))$. In general, $p(\mathbf{Y}|do(\mathbf{A})) \neq p(\mathbf{Y}|\mathbf{A})$ (correlation is not causation).

Specify (qualitative) assumptions on the system using a causal diagram G (here A, Y, Z, K ⊆ X)):



► Given causal diagram G, can derive expressions for causal effect p(Y|A) using do-calculus (Pearl 1995).

$\sum_{\boldsymbol{Z}} p(\boldsymbol{Z}) p(\boldsymbol{Y} \boldsymbol{A}, \boldsymbol{Z})$	$\frac{\sum_{\mathbf{K}} p(\mathbf{A}, \mathbf{Y} \mathbf{K}, \mathbf{Z}) p(\mathbf{K})}{\sum_{\mathbf{K}} p(\mathbf{A} \mathbf{K}, \mathbf{Z}) p(\mathbf{K})}$		
(a) Backdoor	(b) Napkin		

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Tractability of Exact Causal Inference

Consider the backdoor query, for **fixed** values of the treatment \boldsymbol{a} and outcome \boldsymbol{y} :

$$p(\mathbf{y}|do(\mathbf{a})) := \sum_{\mathbf{Z}} p(\mathbf{Z}) imes p(\mathbf{y}|\mathbf{a}, \mathbf{Z})$$

.

Tractability of Exact Causal Inference

Consider the backdoor query, for **fixed** values of the treatment a and outcome y:

$$p(\mathbf{y}|do(\mathbf{a})) := \sum_{\mathbf{Z}} p(\mathbf{Z}) \times p(\mathbf{y}|\mathbf{a}, \mathbf{Z})$$

Theorem (Wang & Kwiatkowska 2023)

If p is given as a structured decomposable and deterministic circuit, then the backdoor query is #P-hard to compute.

Applying the Atlas of Tractable Operations

Break down do-calculus query into compositions of basic operations, such as marginalization, products, and powers:

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Break down do-calculus query into compositions of basic operations, such as marginalization, products, and powers:



Problem: Cannot guarantee that input to POW is deterministic, even if p(X) is deterministic.

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Marginal Determinism

Definition (Marginal Determinism, Choi et al. 2020)

Given a subset of variables $Q \subseteq X$, a PC is Q-deterministic if the children of a sum node T correspond to different values of Q (for sum nodes with $sc(T) \cap Q \neq \emptyset$).

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Motivation: If a circuit is marginally deterministic w.r.t Q, then we can marginalize out $X \setminus Q$ and obtain a deterministic circuit!

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If (the circuit encoding) p(X) is $(A \cup Z)$ -deterministic, then the input to POW is guaranteed to be deterministic.

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 \implies all operations are tractable according to Atlas

If (the circuit encoding) p(X) is $(A \cup Z)$ -deterministic, then the input to POW is guaranteed to be deterministic.



(a) Pipeline for $COND(\cdot, \boldsymbol{A} \cup \boldsymbol{Z})$

(b) Pipeline for entire backdoor query

⇒ all operations are tractable according to Atlas ⇒ can compute causal effect in $O(|p|^3)$ time (can improve to $O(|p|^2)$)

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Open Questions

- Are all causal queries derived by the do-calculus tractable in PTIME (for some non-trivial marginal determinism condition)?
- What is the optimal complexity for these queries?

Questions answered today

- 1. What are probabilistic queries? Are current models tractable? (Guy)
- 2. What are probabilistic circuits and why are they tractable? (Guy)
- 3. What is the connection to logical circuit languages? (YooJung)
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Questions answered today

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tractability vs expressive efficiency

SmoothdecomposabledeterministicstructureddecomposablePCs?

	smooth	dec.	det.	str.dec.
Arithmetic Circuits (ACs) (Darwiche 2003)	~	~	~	×
Sum-Product Networks (SPNs) (Poon et al. 2011)	~	~	×	X
Cutset Networks (CNets) (Rahman et al. 2014)	~	~	~	X
Probabilistic Decision Graphs (Jaeger 2004)	~	~	~	~
(Affine) ADDs (Hoey et al. 1999; Sanner et al. 2005)	~	~	~	~
AndOrGraphs (Dechter et al. 2007)	~	~	~	~
PSDDs (Kisa et al. 2014)	~	~	~	~

Low-treewidh PGMs

Tree, polytrees and Thin Junction trees can be turned into



circuits





Arithmetic Circuits (ACs)





 \Rightarrow parameters are attached to the leaves \Rightarrow ...but can be moved to the sum node edges (Rooshenas et al. 2014)

Lowd and Rooshenas, "Learning Markov Networks With Arithmetic Circuits", 2013

Sum-Product Networks (SPNs)









deterministic SPNs are also called selective (Peharz et al. 2014)

Cutset Networks (CNets)

CNets

(Rahman et al. 2014) are



smooth

deterministic





Rahman et al., "Cutset Networks: A Simple, Tractable, and Scalable Approach for Improving the Accuracy of Chow-Liu Trees", 2014 Di Mauro et al., "Learning Accurate Cutset Networks by Exploiting Decomposability", 2015

Probabilistic Sentential Decision Diagrams







Kisa et al., "Probabilistic sentential decision diagrams", 2014 Choi et al., "Tractable learning for structured probability spaces: A case study in learning preference distributions", 2015 Shen et al., "Conditional PSDDs: Modeling and learning with modular knowledge", 2018

Probabilistic Decision Graphs







Jaeger, "Probabilistic decision graphs—combining verification and AI techniques for probabilistic inference", 2004 Jaeger et al., "Learning probabilistic decision graphs", 2006

AndOrGraphs







Dechter and Mateescu, "AND/OR search spaces for graphical models", 2007 Marinescu and Dechter, "Best-first AND/OR search for 0/1 integer programming", 2007

Probabilistic circuits seem awfully general.

Are all tractable probabilistic models probabilistic circuits?



Enter: Determinantal Point Processes (DPPs)

DPPs are models where probabilities are specified by (sub)determinants

$$L = \begin{bmatrix} 1 & 0.9 & 0.8 & 0\\ 0.9 & 0.97 & 0.96 & 0\\ 0.8 & 0.96 & 1 & 0\\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$\Pr_L(X_1 = 1, X_2 = 0, X_3 = 1, X_4 = 0) = \frac{1}{\det(L+I)} \det(L_{\{1,2\}})$$

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Tractable likelihoods and marginals

Global Negative Dependence

Diversity in recommendation systems

$$\Pr_L(X_1 = 1, X_2 = 0, X_3 = 1, X_4 = 0) = \frac{1}{\det(L+I)} \det(L_{\{1,2\}})$$

Are all tractable probabilistic models probabilistic circuits?



Relationship between PCs and DPPs



We cannot tractably represent DPPs with subclasses of PCs



We cannot tractably represent DPPs with subclasses of PCs



We cannot tractably represent DPPs with subclasses of PCs


We cannot tractably represent DPPs with subclasses of PCs



Honghua Zhang, Steven Holtzen and Guy Van den Broeck. On the Relationship Between Probabilistic Circuits and Determinantal Point Processes, UAI, 2020.

Theorem (Martens and Medabalimi, 2014). Let P_n be the uniform distribution over spanning trees on K_n . For $n \ge 20$, the size of any smooth and decomposable PC that represents P_n is at least $2^{n/30240}$.

Based on arithmetic circuit lower bounds by Ran Raz and Amir Yehudayoff

Decomposable PCs are Syntactically Multilinear Arithmetic Circuits:

Definition 7 (Multilinear Arithmetic Circuit) If every node of an arithmetic circuit Φ over y computes a multilinear polynomial in y, Φ is said to be a *(semantically) multilinear arithmetic circuit*. And if for every product node in Φ , the scopes of its child nodes are pair-wise disjoint, Φ is said to be a *syntactically multilinear arithmetic circuit*.

DPPs have No Compact Decomposable PCs

Theorem (Snell, 1995). The uniform distribution over spanning trees on the complete graph K_n is a DPP over $\binom{n}{2}$ edges.

Theorem (Martens and Medabalimi, 2014). Let P_n be the uniform distribution over spanning trees on K_n . For $n \ge 20$, the size of any smooth and decomposable PC that represents P_n is at least $2^{n/30240}$.



Probabilistic Generating Circuits



A Tractable Unifying Framework for PCs and DPPs



Probability Generating Functions

X_1	X_2	X_3	\Pr_{eta}
0	0	0	0.02
0	0	1	0.08
0	1	0	0.12
0	1	1	0.48
1	0	0	0.02
1	0	1	0.08
1	1	0	0.04
1	1	1	0.16

$$\rightarrow$$

$$g_{\beta} = \underbrace{0.16z_1z_2z_3}_{+ 0.48z_2z_3} + 0.04z_1z_2 + 0.08z_1z_3 + 0.02z_1 + 0.48z_2z_3 + 0.12z_2 + 0.08z_3 + 0.02.$$

Probability Generating Functions

X_1	X_2	X_3	Pr_{eta}
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$$g_{\beta} = \underbrace{0.16z_1z_2z_3}_{+ 0.4z_1z_2 + 0.04z_1z_2 + 0.08z_1z_3 + 0.02z_1}_{+ 0.48z_2z_3 + 0.12z_2 + 0.08z_3 + 0.02.}$$

Probabilistic Generating Circuits (PGCs)



- 1. Sum nodes \bigoplus with weighted edges to children.
- 2. Product nodes 🚫 with unweighted edges to children.
- 3. Leaf nodes: z_i or constant.

PCs as PGCs

(Smooth & Decomposable) PCs represents probability mass functions:
$$\begin{split} m_{\beta} &= 0.16X_{1}X_{2}X_{3} + 0.04X_{1}X_{2}\overline{X_{3}} + 0.08X_{1}\overline{X_{2}}X_{3} + 0.02X_{1}\overline{X_{2}}\overline{X_{3}} \\ &+ 0.48\overline{X_{1}}X_{2}X_{3} + 0.12\overline{X_{1}}X_{2}\overline{X_{3}} + 0.08\overline{X_{1}}\overline{X_{2}}X_{3} + 0.02\overline{X_{1}}\overline{X_{2}}\overline{X_{3}} \end{split}$$

PGCs represent probability generating functions:

$$g_{\beta} = 0.16z_1z_2z_3 + 0.04z_1z_2 + 0.08z_1z_3 + 0.02z_1 + 0.48z_2z_3 + 0.12z_2 + 0.08z_1z_3 + 0.02$$

Given a smooth & decomposable PC, by setting $\overline{X_i}$ to 1, and X_i to z_i , we obtain a PGC that represents the PC.

Tractable Likelihood (EVID) or Marginals (MAR)?



PGCs Support Tractable Likelihoods/Marginals



PGCs Support Tractable Likelihoods/Marginals



PGCs Support Tractable Likelihoods/Marginals



Example



Example



Example



Inference Time Complexity

Given a PGC of size *m* (#edges) over *n* random variables.

Algorithm 1 (*Zhang et al., ICML 2021*):

Bottom-up pass w/ z_i = t, 0 or 1 Product/sum of degree-n polynomials at each node

 $O(mn^2)$

or O(mn log n log log n)

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Bottom-up pass w/ z_i = t, 0 or 1



Product/sum of degree-n polynomials at each node

 $O(mn^2)$

O(mn)

or O(mn log n log log n)

Algorithm 2 (*Harviainen et al., UAI 2023*):

observation: the output of a PGC is a degree-n polynomial w/ respect to t

- + PGCs are tractable when semantically multilinear
 - No need for PC decomposability/syntactic multilinearity or other properties...
- Checking Validity of PGCs is Hard

Theorem (*Harviainen et al.*). It is NP-hard to check if a PGC encodes a valid probability generating polynomial

DPPs as PGCs

The generating polynomial for a DPP with kernel *L* is given by:

$$g_L = rac{1}{\det(L+I)} \det(I + L \operatorname{diag}(z_1, \dots, z_n)).$$

We need it as a sum of products to obtain a Probabilistic Generating Circuit

DPPs as PGCs

The generating polynomial for a DPP with kernel *L* is given by:



We need it as a sum of products to obtain a Probabilistic Generating Circuit

DPPs as PGCs



Experiment Results: Amazon Baby Registries

	DPP	Strudel	EiNet	MT	SimplePGC
apparel	-9.88	-9.51	-9.24	-9.31	$-9.10^{*\dagger \circ}$
bath	-8.55	-8.38	-8.49	-8.53	$-8.29^{*\dagger \circ}$
bedding	-8.65	-8.50	-8.55	-8.59	$-8.41^{*\dagger\circ}$
carseats	-4.74	-4.79	-4.72	-4.76	$-4.64^{*\dagger\circ}$
diaper	-10.61	-9.90	-9.86	-9.93	$-9.72^{*\dagger \circ}$
feeding	-11.86	-11.42	-11.27	-11.30	$-11.17^{*\dagger \circ}$
furniture	-4.38	-4.39	-4.38	-4.43	$-4.34^{*\dagger\circ}$
gear	-9.14	-9.15	-9.18	-9.23	$-9.04^{*\dagger \circ}$
gifts	-3.51	-3.39	-3.42	-3.48	-3.47°
health	-7.40	-7.37	-7.47	-7.49	$-7.24^{*\dagger\circ}$
media	-8.36	-7.62	-7.82	-7.93	$-7.69^{+\circ}$
moms	-3.55	-3.52	-3.48	-3.54	-3.53°
safety	-4.28	-4.43	-4.39	-4.36	$-4.28^{*\dagger\circ}$
strollers	-5.30	-5.07	-5.07	-5.14	$-5.00^{*\dagger \circ}$
toys	-8.05	-7.61	-7.84	-7.88	$-7.62^{\dagger\circ}$

SimplePGC achieves SOTA result on 11/15 datasets

Beyond DPPs: Strongly Rayleigh Distributions

DPPs are strongly Rayleigh distributions

Definition. A probability distribution over binary random variables X_1, \ldots, X_n (or equivalently, subsets of $[n] := \{1, 2, \ldots, n\}$) is strongly Rayleigh if its probability generating polynomial g is real-stable; that is, for $z_i \in \mathbb{C}$, if $Im(z_i) > 0$ for all z_i , then $g(z_1, \cdots, z_n) \neq 0$.

We can efficiently sample from strongly Rayleigh distributions by MCMC (with polynomial bound on mixing time)

Efficient Sampling from SR Distributions

Theorem (Li et al., 2016). Let π be a strongly Rayleigh distribution over [n], we can efficiently sample from π by sampling from its symmetric homogenization π_{sh} ; for $S \subset [2n]$, define

$$\pi_{sh}(S) := \begin{cases} \pi(S \cap [n]) \binom{n}{S \cap [n]}^{-1}, & \text{if } |S| = n \\ 0, & otherwise \end{cases}$$

in particular, π_{sh} is also strongly Rayleigh and the mixing time of a Gibbsexchange sampler with initial set S_0 is bounded as

$$\tau(\epsilon) \le 2n^2 \left(\log \binom{n}{|S_0|} + \log \pi(S_0)^{-1} + \log \epsilon^{-1} \right)$$

Relationship between PGCs and SR Distributions



Relationship between PGCs and SR Distributions



Not All SR Distributions have Compact PGCs (Bläser 2023)

Let $K_{m,n} = (U \cup V, E)$ be a complete bipartite graph, the signed double function generating polynomial is defined as

$$DF_{m,n}(e) = \sum_{F,H} (-1)^{|F|+|H|} \prod_{(i,j)\in F} e_{i,j} \prod_{(i',j')\in H} e_{i',j'}$$

where the sum is taken over all partial functions $U \to V$ and $V \to U$, respectively. Each pair of (F, H) is a double function of $K_{m,n}$.







Figure 4. The thick edges are a matching of size two.

Figure 5. The thick edges form a total function $U \rightarrow V$, which is not injective.

Figure 6. The thick edges form a partial function from V to U.



 v_1

 v_2

 v_3

Not All SR Distributions have Compact PGCs (Bläser 2023)

$$DF_{m,n}(e) = \sum_{F,H} (-1)^{|F|+|H|} \prod_{(i,j)\in F} e_{i,j} \prod_{(i',j')\in H} e_{i',j'}$$

$$Generalize to bipartite multigraph K_{m,n}^{(d)}$$

$$d: each edge from U to V has d copies$$

$$DF_{m,n}^{(d)}(e^{(d)}) = \sum_{F,H} (-1)^{|F|+|H|} \prod_{(i,j)\in F\setminus H} \sum_{\delta=1}^{d} e_{i,j}^{(\delta)} \prod_{(i',j')\in H\cap F} \sum_{1\leq\delta'<\gamma\leq d} e_{i',j'}^{(\delta')} e_{i',j'}^{(\gamma)} \prod_{(i'',j'')\in H\setminus F} \sum_{\delta''=1}^{d} e_{i'',j''}^{(\delta'')}$$

$$DF_{n,n}^{(n+2)} \text{ is real-stable and its evaluation is \#P-hard.}$$

 $DF_{n,n}^{(n+2)}$ does not define an SR distribution as it has negative coefficients

Not All SR Distributions have Compact PGCs (Bläser 2023)

Definition. For a polynomial $f(z_1, \ldots, z_n)$ with z_i of degree k_i , the inversion of f is defined as $\prod_i z_i^{k_i} f(-1/z_1, \ldots, -1/z_i, \ldots, -1/z_n)$.

The inversion of a real stable polynomial is also real stable

Let P_n be the inversion of $DF_{n,n}^{(n+2)}$, then P_n is a mutilinear and real stable polynomial with all coefficients non-negative.

Theorem (Bläser, 2023). Assuming $P^{\#P} \nsubseteq P/Poly$. Let \hat{P}_n be the normalized P_n , then \hat{P}_n cannot be represented as polynomial-size PGCs.

Relationship between PGCs and SR Distributions



Probabilistic **generating** circuits seem awfully general.

Are all tractable probabilistic models probabilistic **generating** circuits?



Questions answered today

- 1. What are probabilistic queries? Are current models tractable? (Guy)
- 2. What are probabilistic circuits and why are they tractable? (Guy)
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Building Probabilistic Circuits



Origins: Compilation
Compiling probabilistic graphical models

Arithmetic circuits

(Darwiche 2002, 2003, 2009)

- Compile a given Bayesian network into an **arithmetic circuit**—a smooth, decomposable and deterministic PCs
- Either via logic encoding of Bayesian network + knowledge compilation
- Or record "execution trace" (sum and product operations) of traditional inference algorithms (junction tree, variable elimination)





Selected references

Logic circuits, interplay between structural properties and tractable reasoning

(Darwiche et al. 2002a)

Converting probabilistic graphical models via knowledge compilation

(Darwiche 2002)

Logic circuit compilers

(Darwiche 2004; Muise et al. 2012; Bova et al. 2015; Lagniez et al. 2017; Oztok et al. 2018)

Neuro-symbolic models using logic circuits

(Ahmed et al. 2022a,b)

Parameter Learning

Gradient descent (of course)

PCs are computational graphs

Hence we can just learn them as any other neural net using SGD

Use re-parameterization if parameters should satisfy constraints:

soft-max for sum-weights (non-negative, sum-to-one)

soft-plus for variances

low-rank plus diagonal for covariance matrices

Allows for conditional distributions

Conditional PCs

(Shao et al. 2019)



Maximum likelihood (frequentist)

PCs can be interpreted as *hierarchical latent variable models*, where each sum node corresponds to a discrete latent variable *(Peharz et al. 2016)*. This allows to perform **classical maximum-likelihood** estimation.





Closed-form maximum likelihood

When the circuit is **deterministic**, there is even an *closed-form ML solution*, which is incredible fast:

```
julia> using ProbabilisticCircuits;
julia> data, structure = load(...);
julia> num_examples(data)
17412
julia> num_edges(structure)
270448
julia> @btime estimate_parameters(structure, data);
63.585 ms (1182350 allocations: 65.97 MiB)
```



<u>Custom SIMD and CUDA kernels</u> to parallelize over layers and training examples. https://github.com/Juice-jl/

Expectation-Maximization

When the PC is not deterministic, we can still apply *expectation-maximization* (*Peharz et al. 2016*). EM can piggy-back on autodfiff:

```
train_x, valid_x, test_x = get_mnist_images([7])
```

```
PC = EinsumNetwork.EinsumNetwork(graph, args)
PC.initialize()
PC.to('cuda')
```

Expectation-Maximization

```
for epoch_count in range(10):
    train_ll, valid_ll, test_ll = compute_loglikelihood()
    start_t = time.time()

    for idx in get_batches(train_x, 100):
        outputs = PC.forward(train_x[idx, :])
        log_likelihood = EinsumNetwork.log_likelihoods(outputs).sum()
        log_likelihood.backward()
        PC.em_process_batch()
```

print_performance(epoch_count, train_ll, valid_ll, test_ll, time.time() - start_t)

https://github.com/cambridge-mlg/EinsumNetworks

Expectation-Maximization

train sample: 5175
parameters: 1573486

[epoch	0]	train LL	-140936.80	valid LL	-140955.72	test LL ·	-141033.80	elapsed	time	3.621	sec
[epoch	1]	train LL	-15916.14	valid LL	-15693.25	test LL	-15976.43	elapsed	time	3.438	sec
[epoch	2]	train LL	-10865.67	valid LL	-10616.72	test LL	-10943.56	elapsed	time	3.436	sec
[epoch	3]	train LL	-10388.53	valid LL	-10158.84	test LL	-10475.49	elapsed	time	3.473	sec
[epoch	4]	train LL	-10264.11	valid LL	-10041.66	test LL	-10352.59	elapsed	time	3.497	sec
[epoch	5]	train LL	-10212.66	valid LL	-10001.09	test LL	-10319.35	elapsed	time	3.584	sec
[epoch	6]	train LL	-10192.21	valid LL	-9965.98	test LL	-10314.84	elapsed	time	3.508	sec
[epoch	7]	train LL	-10153.97	valid LL	-9920.09	test LL	-10261.41	elapsed	time	3.446	sec
[epoch	8]	train LL	-10112.95	valid LL	-9882.48	test LL	-10236.34	elapsed	time	3.579	sec
[epoch	9]	train LL	-10093.31	valid LL	-9862.15	test LL	-10200.94	elapsed	time	3.483	sec

Peharz et al., "Einsum Networks: Fast and Scalable Learning of Tractable Probabilistic Circuits", 2020

Structure Learning

Region graphs

Laying out the PC structure on a high level

Region graphs (RGs) describe decompositional structure

RGs are bipartite, directed graphs containing regions (\mathcal{R}) and partitions (\mathcal{P})

Input and output nodes of the RG are regions

Regions have a *scope* (receptive field), denoted as $sc(\mathcal{R}) \subseteq \mathbf{X}$ For every partition \mathcal{P} it holds that

$$sc(\mathcal{R}_{out}) = \bigcup_{\mathcal{R}_{in} \in inputs(\mathcal{P})} sc(\mathcal{R}_{in})$$
$$sc(\mathcal{R}') \cap sc(\mathcal{R}'') = \emptyset, \qquad \mathcal{R}' \neq \mathcal{R}'' \in inputs(\mathcal{P})$$

Example region graph



(Here, every partition has 2 input regions. This is often assumed, but not necessary.)



Equip each input region with non-linear units













- Equip each *input region* (leaf) \mathcal{R} with K units ϕ_1, \ldots, ϕ_K , which are non-linear functions over $sc(\mathcal{R})$. Usually, ϕ_1, \ldots, ϕ_K are probability densities. K can be different for each input region.
- Equip *each other region* with K sum units. K can be different for each internal region. Often, for the root region K = 1, when PC is used as density estimator.
- Equip each partition \mathcal{P} with as many products as there are combinations of units in the input regions to \mathcal{P} , selecting one unit from each region. Formally, if \mathcal{P} has input regions $\mathcal{R}_1, \mathcal{R}_2, \ldots, \mathcal{R}_I$, insert one product $\prod_{i=1}^{I} u_i$ for each $(u_1, u_2, \ldots, u_I) \in \mathcal{R}_1 \times \mathcal{R}_2 \times \cdots \times \mathcal{R}_I$.

Connect each $\prod_{i=1}^{I} u_i$ in \mathcal{P} to all sum units in the output regions of \mathcal{P} .

- Resulting PC has alternating sum and product units (not a strong constraint)
- We can easily scale the PC (overparameterize, increase expressivity) by equipping regions with more units
- RGs can be seen as a *vectorized version of PCs* each region and partition can be seen as as a module
- Resulting PC will be *smooth and decomposable*, i.e., we can integrate, marginalize, and take conditionals
- After the PC has been constructed, we might discard the RG

Scaling up image models

Latent Variable Distillation

Detect	TPMs					DGMs			
Dataset	LVD (ours)	HCLT	EiNet	RAT-SPN	Glow	RealNVP	BIVA		
ImageNet32	4.38	4.82	5.63	6.90	4.09	4.28	3.96		
ImageNet64	4.12	4.67	5.69	6.82	3.81	3.98	-		
CIFAR	4.37	4.61	5.81	6.95	3.35	3.49	3.08		



Liu et al., "Scaling Up Probabilistic Circuits by Latent Variable Distillation", 2022

How to construct and learn RGs?



The "no-learning" option

(Peharz et al. 2019)

Generating a random region graph, by recursively splitting ${f X}$ into two random parts:



Image-tailored circuit structure

"Recursive image slicing"

(Poon et al. 2011)

Images yield a natural region graph by using axis-aligned splits:

- Start with the full image (=output region)
- Define partitions by applying *horizontal* and *vertical* splits
 - Recurse on the newly generated sub-images (internal regions)
 - Structure somewhat reminiscent to convolutions
 - Generates RGs which are "true DAGs," i.e. regions get re-used









"Recursive data slicing"

(Gens et al. 2013)

Expand regions with **clustering**



"Recursive data slicing"

(Gens et al. 2013)

Number of clusters = number of partitions





"Recursive data slicing"

(Gens et al. 2013)

Try to find independent groups of variables (e.g. independence tests)





"Recursive data slicing"

(Gens et al. 2013)

Success \rightarrow *partition* into new regions



"Recursive data slicing"

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Try to find independent groups of variables (e.g. independence tests)



"Recursive data slicing"

(Gens et al. 2013)

Success \rightarrow *partition* into new regions


"Recursive data slicing"

(Gens et al. 2013)





"Recursive data slicing"

(Gens et al. 2013)

Single variable \rightarrow *input region*



"Recursive data slicing"

(Gens et al. 2013)

Expand regions with **clustering**



"Recursive data slicing"

(Gens et al. 2013)

Number of clusters = number of partitions

And so on...



"Recursive data slicing"

(Gens et al. 2013)

- Stopping conditions: minimal number of features, samples, depth, ...
- Clustering ratios also deliver (initial) parameters
- Smooth & Decomposable Circuits
- **Tractable integration**





Selected references

- **ID-SPN** (Rooshenas et al. 2014)
- LearnSPN-b/T/B (Vergari et al. 2015)
- For heterogeneous data (Molina et al. 2018)
- Using **k-means** (Butz et al. 2018) Or **SVD** splits (Adel et al. 2015)
 - Learning DAGs (Dennis et al. 2015; Jaini et al. 2018)
- Approximating independence tests (Di Mauro et al. 2018)

Cutset networks

Besides clustering, **decision tree learning** can be used as PC learner. **Cutset networks**, decision trees over simple probabilistic models (Chow-Liu trees) (*Rahman et al. 2014*):



Cutset networks can easily be converted into **smooth**, **decomposable and deterministic PCs**.

Decision trees as PCs

Also vanilla decision tree learners can be used to learn PCs, by augmenting the leaves with distributions over inputs (*Correia et al. 2020*). Allows to treat **missing features** and **outlier detection**.





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