# Logic \& Probabilistic Circuits 

# Representation 

## Reasoning

Theory

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## Objective

- Circuits are an assembly language for tractable logic and probabilistic reasoning

Even though logic is central to this Simons program,
we will couch this tutorial in probability...

- Most AI and DB interest in tractable logic circuits for the past 15 years has been as a means of doing probabilistic inference
Much richer query languages <3
$\square$ We live in the age of probabilistic generative A/... :-)
We will spare you most of the machine learning details, and instead focus on representations, query languages, reasoning algorithms, and connections to theory.


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## Questions to be answered today

1. What are probabilistic queries? Are current models tractable? (Guy)
2. What are probabilistic circuits and why are they tractable? (Guy)
3. What is the connection to logical circuit languages? (YooJung)
4. How do I comnile my favorite model into a circuit? (Yoolung)
5. How are circuit size and tractability related? (YooJung)
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7. Are all tractable distributions probabilistic circuits? (Guy)
8. How to learn probabilistic circuits from data? (Guy)

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## Acknowledgements

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Antonio Vergari
Robert Peharz
$\square$ Nicola Di Mauro
■ Honghua Zhang

- Benjie Wang



## The Alphabet Soup of probabilistic models



## Intractable and tractable models



## tractability is a spectrum



## Expressive models without compromises


a unifying framework for tractable models

## Why tractable inference?

or the inherent trade-off of tractability vs. expressiveness

## Why probabilistic inference?

$\mathrm{q}_{1}$ : What is the probability that today is a Monday and there is a traffic jam on Westwood Blvd.?

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## Why probabilistic inference?

$\mathrm{q}_{1}$ : What is the probability that today is a Monday and there is a traffic jam on Westwood Blvd.?

$$
\begin{aligned}
& \mathbf{X}=\left\{\text { Day }, \text { Time, } \text { Jam }_{\text {Str } 1}, \text { Jam }_{\text {Str2 } 2}, \ldots, \text { Jam }_{\text {StrN }}\right\} \\
& \mathrm{q}_{1}(\mathbf{m})=p_{\mathrm{m}}\left(\text { Day }=\text { Mon }, \text { Jam }_{\mathrm{W}_{\text {wood }}}=1\right)
\end{aligned}
$$


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## Why probabilistic inference?

$\mathrm{q}_{1}$ : What is the probability that today is a Monday and there is a traffic jam on Westwood Blvd.?
$\mathbf{X}=\left\{\right.$ Day, Time, Jam $_{\text {Str1 }}$, Jam $_{\text {Str2 }}, \ldots$, Jam $\left._{\text {StrN }}\right\}$
$\mathrm{q}_{1}(\mathbf{m})=p_{\mathrm{m}}\left(\right.$ Day $=$ Mon, $\left.\operatorname{Jam}_{\mathrm{W}_{\mathrm{wood}}}=1\right)$

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## Why probabilistic inference?

$\mathrm{q}_{2}$ : Which day is most likely to have a traffic jam on my route to campus?
$\mathbf{X}=\left\{\right.$ Day, Time, Jamstr1, Jam $_{\text {Str2 }}, \ldots$, Jam $\left._{\text {StrN }}\right\}$
$\mathrm{q}_{2}(\mathbf{m})=\operatorname{argmax}_{\mathrm{d}} p_{\mathrm{m}}\left(\right.$ Day $\left.=\mathrm{d} \wedge \bigvee_{i \in \text { route }} \operatorname{Jam}_{\text {Stri }}\right)$

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$$
\Rightarrow \text { marginals + MAP + logical events }
$$

## Tractable Probabilistic Inference

A class of queries $\mathcal{Q}$ is tractable on a family of probabilistic models $\mathcal{M}$ iff for any query $\mathrm{q} \in \mathcal{Q}$ and model $\mathrm{m} \in \mathcal{M}$ exactly computing $q(\mathbf{m})$ runs in time $O($ poly $(|\mathbf{m}|))$.

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often poly will in fact be linear!

tractable bands

## Complete evidence (EVI)

$\mathrm{q}_{3}$ : What is the probability that today is a Monday at 12.00 and there is a traffic jam only on Westwood Blvd.?

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## Complete evidence (EVI)

$\mathrm{q}_{3}$ : What is the probability that today is a Monday at 12.00 and there is a traffic jam only on Westwood Blvd.?
$\mathbf{X}=\left\{\right.$ Day, Time, Jam $_{\text {Wwood }}$, Jam $_{\text {Str2 }}, \ldots$, Jam $\left._{\text {StrN }}\right\}$
$\mathrm{q}_{3}(\mathbf{m})=p_{\mathrm{m}}(\mathbf{X}=\{$ Mon, $12.00,1,0, \ldots, 0\})$

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## Complete evidence (EVI)

$\mathrm{q}_{3}$ : What is the probability that today is a Monday at 12.00 and there is a traffic jam only on Westwood Blvd.?
$\mathbf{X}=\left\{\right.$ Day, Time, Jam $_{\text {Wwood }}$, Jam $_{\text {Str2 }}, \ldots$, Jam $\left._{\text {StrN }}\right\}$
$\mathrm{q}_{3}(\mathbf{m})=p_{\mathbf{m}}(\mathbf{X}=\{$ Mon, $12.00,1,0, \ldots, 0\})$
...fundamental in maximum likelihood learning

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$$
\theta_{\mathrm{m}}^{\mathrm{MLE}}=\operatorname{argmax}_{\theta} \prod_{\mathbf{x} \in \mathcal{D}} p_{\mathrm{m}}(\mathbf{x} ; \theta)
$$

## Generative Adversarial Networks

$\min _{\theta} \max _{\phi} \mathbb{E}_{\mathbf{x} \sim p_{\text {data }}(\mathbf{x})}\left[\log D_{\phi}(\mathbf{x})\right]+\mathbb{E}_{\mathbf{z} \sim p(\mathbf{z})}\left[\log \left(1-D_{\phi}\left(G_{\theta}(\mathbf{z})\right)\right)\right]$


## 

$\min _{\theta} \max _{\phi} \mathbb{E}_{\mathbf{x} \sim p_{\text {data }}(\mathbf{x})}\left[\log D_{\phi}(\mathbf{x})\right]+\mathbb{E}_{\mathbf{z} \sim p(\mathbf{z})}\left[\log \left(1-D_{\phi}\left(G_{\theta}(\mathbf{z})\right)\right)\right]$

- no explicit likelihood! $\Rightarrow$ adversarial training instead of MLE
$\Rightarrow$ no tractable EVI
- good sample quality
$\Rightarrow$ but lots of samples needed for MC
- unstable training
$\Rightarrow$ mode collapse




## Variational Autoencoders

$$
p_{\theta}(\mathbf{x})=\int p_{\theta}(\mathbf{x} \mid \mathbf{z}) p(\mathbf{z}) d \mathbf{z}
$$

$\square$ an explicit likelihood model!


[^0]

$\log p_{\theta}(\mathbf{x}) \geq \mathbb{E}_{\mathbf{z} \sim q_{\phi}(\mathbf{z} \mid \mathbf{x})}\left[\log p_{\theta}(\mathbf{x} \mid \mathbf{z})\right]-\mathbb{K} \mathbb{L}\left(q_{\phi}(\mathbf{z} \mid \mathbf{x}) \| p(\mathbf{z})\right)$
$\square$ an explicit likelihood model!

- ... but computing $\log p_{\theta}(\mathbf{x})$ is intractable
$\Rightarrow$ an infinite and uncountable mixture $\Rightarrow$ no tractable EVIwe need to optimize the ELBO...

$\Rightarrow$ which is "tricky"



## Normalizing flows

$$
p_{\mathbf{X}}(\mathbf{x})=p_{\mathbf{Z}}\left(f^{-1}(\mathbf{x})\right)\left|\operatorname{det}\left(\frac{\delta f^{-1}}{\delta \mathbf{x}}\right)\right|
$$

$\square$ an explicit likelihood!
$\Rightarrow$ tractable EVI queries!

- many neural variants
$\square$ RealNVP (Dinh et al. 2016),
MAF (Papamakarios et al. 201.



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$\square$ many neural variants
RealNVP (Dinh et al. 2016),
MAF (Papamakarios et al. 2017)
MADE (Germain et al. 2015),
PixelRNN (Oord et al. 2016)

## Marginal queries (MAR)

$\mathrm{q}_{1}$ : What is the probability that today is a Monday 1200 and there is a traffic jam on Westwood Blvd.?

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## Marginal queries (MAR)

$\mathrm{q}_{1}$ : What is the probability that today is a Monday 2 and there is a traffic jam on Westwood Blvd.?
$\mathrm{q}_{1}(\mathbf{m})=p_{\mathrm{m}}\left(\right.$ Day $=$ Mon, $\left.\operatorname{Jam}_{\mathrm{W}_{\text {wood }}}=1\right)$

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## Marginal queries (MAR)

$\mathrm{q}_{1}$ : What is the probability that today is a Monday 2 and there is a traffic jam on Westwood Blvd.?
$\mathrm{q}_{1}(\mathbf{m})=p_{\mathrm{m}}\left(\right.$ Day $=$ Mon, $\left.\operatorname{Jam}_{\mathrm{W}_{\text {wood }}}=1\right)$

General: $p_{\mathrm{m}}(\mathbf{e})=\int p_{\mathrm{m}}(\mathbf{e}, \mathbf{H}) d \mathbf{H}$

© fineartamerica.com where $\mathbf{E} \subset \mathbf{X}, \quad \mathbf{H}=\mathbf{X} \backslash \mathbf{E}$

## Marginal queries (MAR)

$\mathrm{q}_{1}$ : What is the probability that today is a Monday 2 and there is a traffic jam on Westwood Blvd.?
$\mathrm{q}_{1}(\mathbf{m})=p_{\mathrm{m}}\left(\right.$ Day $=$ Mon, $\left.\operatorname{Jam}_{\mathrm{W}_{\text {wood }}}=1\right)$
tractable MAR $\Rightarrow$ tractable conditional queries (CON):

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$$
p_{\mathrm{m}}(\mathbf{q} \mid \mathbf{e})=\frac{p_{\mathrm{m}}(\mathbf{q}, \mathbf{e})}{p_{\mathrm{m}}(\mathbf{e})}
$$

## Tractable MAR: scene understanding




Fast and exact marginalization over unseen or "do not care" parts in the scene
Stelzner et al., "Faster Attend-Infer-Repeat with Tractable Probabilistic Models", 2019
Kossen et al., "Structured Object-Aware Physics Prediction for Video Modeling and Planning", 2019

## Normalizing flows

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p_{\mathbf{X}}(\mathbf{x})=p_{\mathbf{Z}}\left(f^{-1}(\mathbf{x})\right)\left|\operatorname{det}\left(\frac{\delta f^{-1}}{\delta \mathbf{x}}\right)\right|
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$\square$ an explicit likelihood!
$\Rightarrow$ tractable EVI queries!


##  

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p_{\mathbf{X}}(\mathbf{x})=p_{\mathbf{Z}}\left(f^{-1}(\mathbf{x})\right)\left|\operatorname{det}\left(\frac{\delta f^{-1}}{\delta \mathbf{x}}\right)\right|
$$

- an explicit likelihood!
$\Rightarrow$ tractable EVI queries!


## MAR is generally intractable:

 we cannot easily integrate over high-dimensional $f$


## Probabilistic Graphical Models (PGMs)

Declarative semantics: a clean separation of modeling assumptions from inference

Nodes: random variables
Edges: dependencies


Inference: $\quad$ conditioning (Darwiche 2001; Sang et al. 2005)

- elimination (Zhang et al. 1994; Dechter 1998)
$\square$ message passing (Yedidia et al. 2001; Dechter
et al. 2002; Choi et al. 2010; Sontag et al. 2011)


## Complexity of MAR on PGMs

Exact complexity: Computing MAR and CON is \#P-hard $\Rightarrow \quad$ (Cooper 1990; Roth 1996)

Approximation complexity: Computing MAR and CON approximately within a relative error of $2^{n^{1-\epsilon}}$ for any fixed $\epsilon$ is NP-hard

## Treewidth!

## Treewidth:

Informally, how tree-like is the graphical model m?

Fixed-parameter tractable: MAR and CON on a graphical model m with treewidth $w$ take time $O\left(|\mathbf{X}| \cdot 2^{w}\right)$ (Dechter 1998; Kolle et al. 2009). $\Rightarrow \quad$ what about bounding the treewidth by design?

## Low-treewidth PGMs


Trees
(Meilă et al. 2000)

Polytrees
(Dasgupta 1999)

Thin Junction trees
(Bach et al. 2001)

If treewidth is bounded (e.g. $\cong 20$ ), exact MAR and CON inference is possible in practice

## Tree distributions

A tree-structured BN (Meilă et al. 2000) where each $X_{i} \in \mathbf{X}$ has at most one parent $\mathrm{Pa}_{X_{i}}$.


$$
p(\mathbf{X})=\prod_{i=1}^{n} p\left(x_{i} \mid \mathrm{Pa}_{x_{i}}\right)
$$

Exact querying: EVI, MAR, CON tasks linear for trees: $O(|\mathbf{X}|)$
Exact learning from $d$ examples takes $O\left(|\mathbf{X}|^{2} \cdot d\right)$ with the classical Chow-Liu algorithm ${ }^{1}$


## What do we lose?

Expressiveness: Ability to represent rich and complex classes of distributions


Bounded-treewidth PGMs lose the ability to represent all possible distributions ...

[^1]
## Mixtures

Mixtures as a convex combination of $k$ (simpler) probabilistic models


$$
p(X)=w_{1} \cdot p_{1}(X)+w_{2} \cdot p_{2}(X)
$$

EVI, MAR, CON queries scale linearly in $k$

## Mixtures

Mixtures as a convex combination of $k$ (simpler) probabilistic models


$$
\begin{aligned}
p(X)= & p(Z=1) \cdot p_{1}(X \mid Z=1) \\
& +p(Z=\mathbf{2}) \cdot p_{2}(X \mid Z=\mathbf{2})
\end{aligned}
$$

Mixtures are marginalizing a categorical latent variable $Z$ with $k$ values
$\Rightarrow$ increased expressiveness

## Expressiveness and efficiency

Expressiveness: Ability to represent rich and effective classes of functions
$\Rightarrow$ mixture of Gaussians can approximate any probability density!

[^2]
## Expressiveness and efficiency

Expressiveness: Ability to represent rich and effective classes of functions
$\Rightarrow$ mixture of Gaussians can approximate any probability density!

## Expressive efficiency (aka Succinctness):

Ability to represent rich and effective classes of functions compactly $\Rightarrow$ but how many components does a Gaussian mixture need?

Cohen et al., "On the expressive power of deep learning: A tensor analysis", 2016

## How expressive efficient are mixtures?



## How expressive efficient are mixtures?



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## How expressive efficient are mixtures?



## How expressive efficient are mixtures?


$\Rightarrow$ solution: deep mixtures as in deep generative models


## Maximum A Posteriori (MAP)

aka Most Probable Explanation (MPE)
$\mathrm{q}_{5}$ : Which combination of roads is most likely to be jammed on Monday at 9am?

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## Maximum A Posteriori (MAP)

aka Most Probable Explanation (MPE)
$q_{5}$ : Which combination of roads is most likely to be jammed on Monday at 9am?

$$
\mathrm{q}_{5}(\mathbf{m})=\operatorname{argmax}_{\mathbf{j}} p_{\mathbf{m}}\left(\mathbf{j}_{1}, \mathbf{j}_{2}, \ldots \mid \text { Day }=\mathrm{M}, \text { Time }=9\right)
$$


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## Maximum A Posteriori (MAP)

aka Most Probable Explanation (MPE)
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General: $\operatorname{argmax}_{\mathbf{q}} p_{\mathrm{m}}(\mathbf{q} \mid \mathbf{e})$

$$
\text { where } \mathbf{Q} \cup \mathbf{E}=\mathbf{X}
$$


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## Maximum A Posteriori (MAP)

aka Most Probable Explanation (MPE)
$\mathrm{q}_{5}$ : Which combination of roads is most likely to be jammed on Monday at 9am?
...intractable for latent variable models!

$$
\begin{aligned}
\max _{\mathbf{q}} p_{\mathbf{m}}(\mathbf{q} \mid \mathbf{e}) & =\max _{\mathbf{q}} \sum_{\mathbf{z}} p_{\mathbf{m}}(\mathbf{q}, \mathbf{z} \mid \mathbf{e}) \\
& \neq \sum_{\mathbf{z}} \max _{\mathbf{q}} p_{\mathbf{m}}(\mathbf{q}, \mathbf{z} \mid \mathbf{e})
\end{aligned}
$$


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## MAP inference: image inpainting



# Predicting arbitrary patches 

given a single model
without the need of retraining.

[^3]

## Marginal MAP (MMAP)

aka Bayesian Network MAP
$\mathrm{q}_{6}$ : Which combination of roads is most likely to be jammed at 9am?

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=\operatorname{argmax}_{\mathbf{q}} \sum_{\mathbf{h}} p_{\mathrm{m}}(\mathbf{q}, \mathbf{h} \mid \mathbf{e})
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$$

$\Rightarrow N P^{P P}$-complete (Park et al. 2006)
$\Rightarrow$ NP-hard for trees (de Campos 2011)
$\Rightarrow$ NP-hard even for Naive Bayes (ibid.)

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## Advanced queries

$\mathrm{q}_{2}$ : Which day is most likely to have a traffic jam on my route to campus?

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$\mathrm{q}_{2}$ : Which day is most likely to have a traffic jam on my route to campus?

$$
\mathrm{q}_{2}(\mathbf{m})=\operatorname{argmax}_{\mathrm{d}} p_{\mathrm{m}}\left(\text { Day }=\mathrm{d} \wedge \bigvee_{i \in \text { route }} \operatorname{Jam}_{\mathrm{Str} i}\right)
$$

$$
\Rightarrow \text { marginals + MAP + logical events }
$$


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## Advanced queries

$\mathrm{q}_{2}$ : Which day is most likely to have a traffic jam on my route to campus?
$\mathrm{q}_{7}$ : What is the probability of seeing more traffic jams in Westwood than Hollywood?

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## Advanced queries

$\mathrm{q}_{2}$ : Which day is most likely to have a traffic jam on my route to campus?
$\mathrm{q}_{7}$ : What is the probability of seeing more traffic jams in Westwood than Hollywood?

$$
\Rightarrow \text { counts + group comparison }
$$


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## Advanced queries

$\mathrm{q}_{2}$ : Which day is most likely to have a traffic jam on my route to campus?
$\mathrm{q}_{7}$ : What is the probability of seeing more traffic jams in Westwood than Hollywood?
$\mathrm{q}_{8}$ : Is traffic more uncertain on weekdays?

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## Advanced queries

$\mathrm{q}_{2}$ : Which day is most likely to have a traffic jam on my route to campus?
$\mathrm{q}_{7}$ : What is the probability of seeing more traffic jams in Westwood than Hollywood?
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## $\Rightarrow$ information-theoretic queries


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## Advanced queries

$\mathrm{q}_{2}$ : Which day is most likely to have a traffic jam on my route to campus?
$\mathrm{q}_{7}$ : What is the probability of seeing more traffic jams in Westwood than Hollywood?
$\mathrm{q}_{8}$ : Is traffic more uncertain on weekdays?
$\mathrm{q}_{9}$ : What is the causal effect of doing road works?

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## Advanced queries

$\mathrm{q}_{2}$ : Which day is most likely to have a traffic jam on my route to campus?
$\mathrm{q}_{7}$ : What is the probability of seeing more traffic jams in Westwood than Hollywood?
$\mathrm{q}_{8}$ : Is traffic more uncertain on weekdays?
$\mathrm{q}_{9}$ : What is the causal effect of doing road works?
$\Rightarrow$ causal backdoor estimation

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## Fully factorized models

A completely disconnected graph. Example: Product of Bernoullis (PoBs)


$$
p(\mathbf{x})=\prod_{i=1}^{n} p\left(x_{i}\right)
$$



$x_{5}$

Complete evidence, marginals and MAP, MMAP inference is linear!
$\Rightarrow$ but definitely not expressive...

larger tractable bands

larger tractable bands


## Expressive models are not very tractable...



## and tractable ones are not very expressive...



## probabilistic circuits are at the "sweet spot"

## Questions answered today

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## Probabilistic Circuits

## Goal

Given a reasoning task can we design a class of expressive models that is tractable for it?

## Goal

Given a reasoning task can we design
a class of deep computational graphs
that is tractable for it?
more tractable

more tractable


## Expressive models are not very tractable...

more tractable


## Tractable models are not that expressive...

more tractable


## Circuits can be both expressive and tractable!



## Start simple...




## Input distributions

as computational nodes


Base case: a single node encoding a distribution
$\Rightarrow$ e.g., Gaussian PDF continuous random variable

## Input distributions

as computational nodes


Base case: a single node encoding a distribution
$\Rightarrow$ e.g., indicators for $X$ or $\neg X$ for Boolean random variable

## Input distributions

as computational nodes


Simple distributions are tractable "black boxes" for:
$\square$ EVI: output $p(\mathbf{x})$ (density or mass)
MAR: output 1 (normalized) or $Z$ (unnormalized)
MAP: output the mode

## Mixture models

as computational graphs


$$
p(X)=w_{1} \cdot p_{1}\left(X_{1}\right)+w_{2} \cdot p_{2}\left(X_{1}\right)
$$

$\Rightarrow$ translating inference to data structures...

## Mixture models

as computational graphs


$$
p\left(X_{1}\right)=0.2 \cdot p_{1}\left(X_{1}\right)+0.8 \cdot p_{2}\left(X_{1}\right)
$$

$\Rightarrow$
...e.g., as a weighted sum unit over Gaussian input distributions

## Mixture models

as computational graphs


$$
\begin{aligned}
p(X=5) & =0.2 \cdot p_{1}\left(X_{1}=5\right) \\
& +0.8 \cdot p_{2}\left(X_{1}=5\right)
\end{aligned}
$$

$\Rightarrow$ inference $=$ feedforward evaluation

## Mixture models

as computational graphs


A simplified notation:
$\Rightarrow$ scopes attached to inputs $\Rightarrow$ edge directions omitted

## Factorizations

as computational graphs

$$
p\left(X_{1}, X_{2}, X_{3}\right)=p\left(X_{1}\right) \cdot p\left(X_{2}\right) \cdot p\left(X_{3}\right)
$$


$\Rightarrow$ e.g. modeling a multivariate Gaussian with diagonal covariance matrix...

## Factorizations

as computational graphs

$$
p\left(X_{1}, X_{2}, X_{3}\right)=p\left(X_{1}\right) \cdot p\left(X_{2}\right) \cdot p\left(X_{3}\right)
$$


$\Rightarrow$...with a product node over some univariate Gaussian distribution

## Factorizations

## as computational graphs

$$
p\left(x_{1}, x_{2}, x_{3}\right)=p\left(x_{1}\right) \cdot p\left(x_{2}\right) \cdot p\left(x_{3}\right)
$$


$\Rightarrow$ feedforward evaluation

## Factorizations

as computational graphs

$$
p\left(x_{1}, x_{2}, x_{3}\right)=p\left(x_{1}\right) \cdot p\left(x_{2}\right) \cdot p\left(x_{3}\right)
$$



$\Rightarrow$ feedforward evaluation

## A grammar for tractable models

## Recursive semantics of probabilistic circuits

## A grammar for tractable models

## Recursive semantics of probabilistic circuits



## A grammar for tractable models

## Recursive semantics of probabilistic circuits



## A grammar for tractable models

## Recursive semantics of probabilistic circuits



## A grammar for tractable models

Recursive semantics of probabilistic circuits


## Building PCs in Python with SPFlow

import spn.structure.leaves.parametric.Parametric as param from param import Categorical, Gaussian

$$
\begin{aligned}
& \mathrm{PC}=0.4 * \text { (Categorical }(\mathrm{p}=[0.2,0.8] \text {, scope=0) * } \\
& \text { (0.3 * (Gaussian (mean }=1.0 \text {, stdev=1.0, scope=1) * } \\
& \text { Categorical (p=[0.4, 0.6], scope=2)) } \\
& +0.7 \text { * (Gaussian (mean }=-1.0 \text {, stdev }=1.0 \text {, scope=1) * } \\
& \text { Categorical }(\mathrm{p}=[0.6,0.4], \mathrm{scope}=2))) \text { ) \} } \\
{+0.6 \text { * (Categorical (p=[0.2, 0.8], scope=0) * }} \\
{\text { Gaussian (mean }=0.0 \text {, stdev }=0.1 \text {, scope=1) * }} \\
{\text { Categorical (p=[0.4, 0.6], scope=2)) }}
\end{aligned}
$$

Molina et al., "SPFlow: An easy and extensible library for deep probabilistic learning using

## EVI queries $=$ feedforward evaluation

$$
p\left(X_{1}=-1.85, X_{2}=0.5, X_{3}=-1.3, X_{4}=0.2\right)
$$



## EVI queries $=$ feedforward evaluation

$$
p\left(X_{1}=-1.85, X_{2}=0.5, X_{3}=-1.3, X_{4}=0.2\right)
$$



## EVI queries $=$ feedforward evaluation

$$
p\left(X_{1}=-1.85, X_{2}=0.5, X_{3}=-1.3, X_{4}=0.2\right)=0.75
$$



## Just sum, products and distributions?


just arbitrarily compose them like a neural network!

## Just sum, products and distributions?



## Which structural constraints ensure tractability?

## Decomposability

A product node is decomposable if its children depend on disjoint sets of variables $\Rightarrow$ just like in factorization!

decomposable circuit

non-decomposable circuit

## Smoothness

aka completeness
A sum node is smooth if its children depend of the same variable sets
$\Rightarrow$ otherwise not accounting for some variables

smooth circuit

non-smooth circuit

## Smoothness + decomposability $=$ tractable MAR

Computing arbitrary integrations (or summations)

$$
\Rightarrow \quad \text { linear in circuit size! }
$$

E.g., suppose we want to compute $Z$ (the distribution's normalizing constant):

$$
\int \boldsymbol{p}(\mathbf{x}) d \mathbf{x}
$$

## Smoothness + decomposability $=$ tractable MAR

$$
\text { If } p(\mathbf{x})=\sum_{i} w_{i} p_{i}(\mathbf{x}), \text { (smoothness): }
$$

$$
\int p(\mathbf{x}) d \mathbf{x}=\int \sum_{i} w_{i} p_{i}(\mathbf{x}) d \mathbf{x}=
$$

$$
=\sum_{i} w_{i} \int p_{i}(\mathbf{x}) d \mathbf{x}
$$

$\Rightarrow$ integrals are "pushed down" to children


## Smoothness + decomposability $=$ tractable MAR

$$
\text { If } p(\mathbf{x}, \mathbf{y}, \mathbf{z})=p(\mathbf{x}) p(\mathbf{y}) p(\mathbf{z}) \text {, (decomposability): }
$$

$$
\begin{aligned}
& \iiint p(\mathbf{x}, \mathbf{y}, \mathbf{z}) d \mathbf{x} d \mathbf{y} d \mathbf{z}= \\
= & \iiint p(\mathbf{x}) p(\mathbf{y}) p(\mathbf{z}) d \mathbf{x} d \mathbf{y} d \mathbf{z}= \\
= & \int p(\mathbf{x}) d \mathbf{x} \int p(\mathbf{y}) d \mathbf{y} \int p(\mathbf{z}) d \mathbf{z}
\end{aligned}
$$

$\Rightarrow$ integrals decompose into easier ones


## Smoothness + decomposability $=$ tractable MAR

Forward pass evaluation for MAR $\Rightarrow$ linear in circuit size! E.g. to compute $p\left(x_{2}, x_{4}\right)$ :

- leafs over $X_{1}$ and $X_{3}$ output $Z_{i}=\int p\left(x_{i}\right) d x_{i}$
leafs over $X_{2}$ and $X_{4}$ output EVI feedforward evaluation (bottom-up)



## Smoothness + decomposability $=$ tractable MAR

## Forward pass evaluation for MAR

$\Rightarrow$ linear in circuit size!
E.g. to compute $p\left(x_{2}, x_{4}\right)$ :

- leafs over $X_{1}$ and $X_{3}$ output $Z_{i}=\int p\left(x_{i}\right) d x_{i}$
$\Rightarrow$ for normalized leaf distributions: 1.0
- leafs over $X_{2}$ and $X_{4}$ output EVI
- feedforward evaluation (bottom-up)



## Smoothness + decomposability $=$ tractable MAR

Forward pass evaluation for MAR
$\Rightarrow$ linear in circuit size!
E.g. to compute $p\left(x_{2}, x_{4}\right)$ :

- leafs over $X_{1}$ and $X_{3}$ output $Z_{i}=\int p\left(x_{i}\right) d x_{i}$
$\Rightarrow$ for normalized leaf distributions: 1.0
- leafs over $X_{2}$ and $X_{4}$ output EVI
$\square$ feedforward evaluation (bottom-up)



## Tractable MAR



Peharz et al., "Einsum Networks: Fast and Scalable Learning of Tractable Probabilistic Circuits", 2020

## Smoothness + decomposability $=$ tractable CON

Analogously, for arbitrary conditional queries:

$$
p(\mathbf{q} \mid \mathbf{e})=\frac{p(\mathbf{q}, \mathbf{e})}{p(\mathbf{e})}
$$

1. evaluate $p(\mathbf{q}, \mathbf{e}) \Rightarrow$ one feedforward pass
2. evaluate $p(\mathbf{e}) \Rightarrow$ another feedforward pass $\Rightarrow$...still linear in circuit size!


## Tractable CON

Original


Missing


Conditional sample

| 7 | 7 | 7 | 7 |  |
| :--- | :--- | :--- | :--- | :--- |
| 7 | 7 | 7 | 7 |  |
| 7 | 4 | 7 | 7 |  |
| 7 | 7 | 7 | 7 |  |
| 7 | 7 | 7 | 7 | 1 |

## Generative models are still hard to control

more reasoning!


Generate a sentence using "frisbee", "caught" and "dog", following the given order.

Generate a sentence using "frisbee", "caught" and "dog", following the given order.

After a perfect throw, the frisbee glided through the air, and the dog, with incredible agility, caught it mid-flight.

Generate a sentence using "frisbee", "caught" and "dog", following the given order.

After a perfect throw, the frisbee glided through the air, and the dog, with incredible agility, caught it mid-flight.
$\bigcirc$ That's not correct. Generate a sentence using "frisbee", "caught" and "dog". The keywords should appear in the order as specified.

Here's the correct sentence: The dog caught the frisbee in mid-air, showing off its amazing catching skills.

Generate a sentence using "frisbee", "caught" and "dog", following the given order.

After a perfect throw, the frisbee glided through the air, and the dog, with incredible agility, caught it mid-flight.

That's not correct. Generate a sentence using "frisbee", "caught" and "dog". The keywords should appear in the order as specified.

Here's the correct sentence: The dog caught the frisbee in mid-air, showing off its amazing catching skills.

A frisbee is caught by a dog.
A pair of frisbee players are caught in a dog fight.

## What do we have?

Prefix: "The weather is"
Constraint $\alpha$ : text contains "winter"
Model only does $p($ next-token $\mid$ prefix $)=$

| cold | 0.05 |
| :---: | :---: |
| warm | 0.10 |

## What do we need?

Prefix: "The weather is"
Constraint $\alpha$ : text contains "winter"

Generate from $p($ next-token $\mid$ prefix,$\alpha)=$| cold | 0.50 |
| :--- | :--- |
| warm | 0.01 |

$\propto \sum_{\text {text }} p($ next-token, text, prefix, $\alpha)$

## Marginalization!

## Computing $p\left(\alpha \mid x_{1: t+1}\right)$

For $\alpha$ in conjunctive normal form (CNF):

$$
\left(w_{1,1} \vee \ldots \vee w_{1, d 1}\right) \wedge \ldots \wedge\left(w_{m, 1} \vee \ldots \vee w_{m, d m}\right)
$$

where each $\mathrm{w}_{\mathrm{i}}$ is a keyword (i.e. a string of tokens), representing the constraint that $\mathrm{w}_{\mathrm{ij}}$ appears in the generated text.


## Computing $p\left(\alpha \mid x_{1: t+1}\right)$

For $\boldsymbol{\alpha}$ in conjunctive normal form (CNF):

$$
\left(w_{1,1} \vee \ldots \vee w_{1, d 1}\right) \wedge \ldots \wedge\left(w_{m, 1} \vee \ldots \vee w_{m, d m}\right)
$$

where each $w_{i \mathrm{i}}$ is a keyword (i.e. a string of tokens), representing the constraint that $w_{i j}$ appears in the generated text.
e.g., $\alpha=($ "swims" $\vee$ "like swimming") $\wedge($ "lake" $\vee$ "pool")

## Efficient algorithm:

For $m$ clauses and sequence length $n$, time-complexity for generation is $O\left(2^{|m|} n\right)$ when $p$ is a hidden Markov model (see general probabilistic circuit case later).

Trick: dynamic programming with clever preprocessing and local belief updates

## CommonGen: a Challenging Benchmark

Given 3-5 concepts (keywords), our goal is to generate a sentence using all keywords, which can appear in any order and any form of inflections. e.g.,

## Input: snow drive car

Reference 1: A car drives down a snow covered road.
Reference 2: Two cars drove through the snow.

$$
\left(w_{1,1} \vee \ldots \vee w_{1, \mathrm{~d} 1}\right) \wedge \ldots \wedge\left(w_{m, 1} \vee \ldots \vee w_{m, d m}\right)
$$

Each clause represents the inflections for one keyword.

GeLaTo Overview

Lexical Constraint $\alpha$ : sentence contains keyword "winter"


GeLaTo Overview

Lexical Constraint $\alpha$ : sentence contains keyword "winter"


## Step 2: Control $p_{g p t}$ via $p_{h m m}$

## Unsupervised

Language model is not
fine-tuned/prompted to satisfy constraints

By Bayes rule:

$$
p_{g p t}\left(x_{t+1} \mid x_{1: t}, \alpha\right) \propto p_{g p t}\left(\alpha \mid x_{1: t+1}\right) \cdot p_{g p t}\left(x_{t+1} \mid x_{1: t}\right)
$$

Assume $p_{h m m}\left(\alpha \mid x_{1: t+1}\right) \approx p_{g p t}\left(\alpha \mid x_{1: t+1}\right)$, we generate from:

$$
p\left(x_{t+1} \mid x_{1: t}, \alpha\right) \propto p_{h m m}\left(\alpha \mid x_{1: t+1}\right) \cdot p_{g p t}\left(x_{t+1} \mid x_{1: t}\right)
$$

| Method | Generation Quality |  |  |  |  |  |  |  | Constraint Satisfaction |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | ROU | E-L |  |  |  |  |  |  |  |  | Succe | Rate |
| Unsupervised | dev | test | dev | test | dev | test | dev | test | dev | test | dev | test |
| InsNet (Lu et al., 2022a) | - | - | 18.7 | - | - | - | - | - | 100.0 | - | 100.0 | - |
| NeuroLogic (Lu et al., 2021) | - | 41.9 | - | 24.7 | - | 14.4 | - | 27.5 | - | 96.7 | - | - |
| A*esque (Lu et al., 2022b) | - | 44.3 | - | 28.6 | - | 15.6 | - | 29.6 | - | 97.1 | - | - |
| NADO (Meng et al., 2022) | - | - | 26.2 | - | - | - | - | - | 96.1 | - | - | - |
| GeLaTo | 44.6 | 44.1 | 29.9 | 29.4 | 16.0 | 15.8 | 31.3 | 31.0 | 100.0 | 100.0 | 100.0 | 100.0 |

## Step 2: Control $p_{g p t}$ via $p_{h m m}$

## Supervised

Language model is fine-tuned to perform constrained generation (e.g. seq2seq)

Empirically $p_{H M M}\left(\alpha \mid x_{1: t+1}\right) \approx p_{g p t}\left(\alpha \mid x_{1: t+1}\right)$ does not hold well enough;
we view $p_{H M M}\left(x_{t+1} \mid x_{1: t}, \alpha\right)$ and $p_{g p t}\left(x_{t+1} \mid x_{1: t}\right)$ as classifiers trained for the same task with different biases; thus we generate from their weighted geometric mean:

$$
p\left(x_{t+1} \mid x_{1: t}, \alpha\right) \propto p_{h m m}\left(x_{t+1} \mid x_{1: t}, \alpha\right)^{w} \cdot p_{g p t}\left(x_{t+1} \mid x_{1: t}\right)^{1-w}
$$

| Method | Generation Quality |  |  |  |  |  |  |  | Constraint Satisfaction |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | ROUGE-L |  | BLEU-4 |  | CIDEr |  | SPICE |  | Coverage |  | Success Rate |  |
| Supervised | dev | test | dev | test | dev | test | dev | test | dev | test | dev | test |
| NeuroLogic (Lu et al., 2021) | - | 42.8 | - | 26.7 | - | 14.7 | - | 30.5 | - | 97.7 | - | $93.9{ }^{\dagger}$ |
| A*esque (Lu et al., 2022b) | - | 43.6 | - | 28.2 | - | 15.2 | - | 30.8 | - | 97.8 | - | $97.9^{\dagger}$ |
| NADO (Meng et al., 2022) | $44.4{ }^{\dagger}$ | - | 30.8 | - | $16.1{ }^{\dagger}$ | - | 32.0 ${ }^{\dagger}$ | - | 97.1 | - | $88.8{ }^{\dagger}$ | - |
| GeLaTo | 46.0 | 45.6 | 34.1 | 32.9 | 16.7 | 16.8 | 31.3 | 31.9 | 100.0 | 100.0 | 100.0 | 100.0 |

## Advantages of GeLaTo:

1. Constraint $\alpha$ is guaranteed to be satisfied: for any next-token $x_{t+1}$ that would make $\alpha$ unsatisfiable, $p\left(x_{t+1} \mid x_{1: t^{\prime}} \alpha\right)=0$ for both settings.
2. Training $p_{\text {hmm }}$ does not depend on $\alpha$, which is only imposed at inference (generation) time. Once $p_{\text {hmm }}$ is trained, we can impose whatever $\alpha$.
3. We can impose additional tractable constraints:

- The keywords are generated following a particular order.
- (Some) keywords must appear at a particular position.
- (Some) keywords must not appear in the generated sentence.

Conclusion: you can control an intractable generative model using a tractable probabilistic circuit.

## Smoothness + decomposalbility $=$ tractalble MAP

We can also decompose bottom-up a MAP query:

$$
\max _{\mathbf{q}} p(\mathbf{q} \mid \mathbf{e})
$$

## Smoothness + decomposability = ewetulenna

We cannot decompose bottom-up a MAP query:

$$
\max _{\mathbf{q}} p(\mathbf{q} \mid \mathbf{e})
$$

since for a sum node we are marginalizing out a latent variable

$$
\max _{\mathbf{q}} \sum_{i} w_{i} p_{i}(\mathbf{q}, \mathbf{e})=\max _{\mathbf{q}} \sum_{\mathbf{z}} p(\mathbf{q}, \mathbf{z}, \mathbf{e}) \neq \sum_{\mathbf{z}} \max _{\mathbf{q}} p(\mathbf{q}, \mathbf{z}, \mathbf{e})
$$

$\Rightarrow$ MAP for latent variable models is intractable (Conaty et al. 2017)

## Determinism

## aka selectivity

A sum node is deterministic if the output of only one children is non zero for any input $\Rightarrow$ e.g. if their distributions have disjoint support

deterministic circuit

non-deterministic circuit

## Determinism + decomposability $=$ tractable MAP

Computing maximization with arbitrary evidence $\mathbf{e}$ $\Rightarrow$ linear in circuit size!
E.g., suppose we want to compute:

$$
\max _{\mathbf{q}} p(\mathbf{q} \mid \mathbf{e})
$$



## Determinism + decomposability $=$ tractable MAP

$$
\text { If } p(\mathbf{q}, \mathbf{e})=\sum_{i} w_{i} \boldsymbol{p}_{i}(\mathbf{q}, \mathbf{e})=\max _{i} w_{i} \boldsymbol{p}_{i}(\mathbf{q}, \mathbf{e})
$$ (deterministic sum node):

$$
\begin{aligned}
\max _{\mathbf{q}} p(\mathbf{q}, \mathbf{e}) & =\max _{\mathbf{q}} \sum_{i} w_{i} p_{i}(\mathbf{q}, \mathbf{e}) \\
& =\max _{\mathbf{q}} \max _{i} w_{i} p_{i}(\mathbf{q}, \mathbf{e}) \\
& =\max _{i} \max _{\mathbf{q}} w_{i} \boldsymbol{p}_{i}(\mathbf{q}, \mathbf{e})
\end{aligned}
$$

$\Rightarrow$ one non-zero child term, thus sum is max


## Determinism + decomposability $=$ tractable MAP

If $p(\mathbf{q}, \mathbf{e})=p\left(\mathbf{q}_{\mathbf{x}}, \mathbf{e}_{\mathbf{x}}, \mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}}\right)=p\left(\mathbf{q}_{\mathbf{x}}, \mathbf{e}_{\mathbf{x}}\right) p\left(\mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}}\right)$ (decomposable product node):

$$
\begin{aligned}
& \max _{\mathbf{q}} p(\mathbf{q} \mid \mathbf{e})=\max _{\mathbf{q}} p(\mathbf{q}, \mathbf{e}) \\
& \quad=\max _{\mathbf{q}_{\mathbf{x}}, \mathbf{q}_{\mathbf{y}}} p\left(\mathbf{q}_{\mathbf{x}}, \mathbf{e}_{\mathbf{x}}, \mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}}\right) \\
& \quad=\max _{\mathbf{q}_{\mathbf{x}}} p\left(\mathbf{q}_{\mathbf{x}}, \mathbf{e}_{\mathbf{x}}\right) \cdot \max _{\mathbf{q}_{\mathbf{y}}} p\left(\mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}}\right) \\
& \quad \Rightarrow \text { solving optimization independently }
\end{aligned}
$$



## Determinism + decomposability $=$ tractable MAP

Evaluating the circuit twice: bottom-up and top-down $\Rightarrow$ still linear in circuit size!


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Evaluating the circuit twice: bottom-up and top-down $\Rightarrow$ still linear in circuit size!
E.g., for $\operatorname{argmax}_{x_{1}, x_{3}} p\left(x_{1}, x_{3} \mid x_{2}, x_{4}\right)$ :

1. turn sum into max nodes and distributions into max distributions
2. evaluate $p\left(x_{2}, x_{4}\right)$ bottom-up
3. retrieve max activations top-down
4. compute Mapstates for $X_{1}$ and $X_{3}$ at leaves


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## MAP inference: image segmentation

Input Image


Multiscale sum-product


Semantic segmentation is MAP over joint pixel and label space
Even approximate MAP for non-deterministic circuits (SPNs) delivers good performances.

[^4]
## How expressive?

| Dataset | Sparse PC (ours) | HCLT | RatSPN | IDF | BitSwap | BB-ANS | McBits |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| MNIST | $\mathbf{1 . 1 4}$ | 1.20 | 1.67 | 1.90 | 1.27 | 1.39 | 1.98 |
| EMNIST(MNIST) | $\mathbf{1 . 5 2}$ | 1.77 | 2.56 | 2.07 | 1.88 | 2.04 | 2.19 |
| EMNIST(Letters) | $\mathbf{1 . 5 8}$ | 1.80 | 2.73 | 1.95 | 1.84 | 2.26 | 3.12 |
| EMNIST(Balanced) | $\mathbf{1 . 6 0}$ | 1.82 | 2.78 | 2.15 | 1.96 | 2.23 | 2.88 |
| EMNIST(ByClass) | $\mathbf{1 . 5 4}$ | 1.85 | 2.72 | 1.98 | 1.87 | 2.23 | 3.14 |
| FashionMNIST | $\mathbf{3 . 2 7}$ | 3.34 | 4.29 | 3.47 | 3.28 | 3.66 | 3.72 |

## competitive with Flows and VAEs!

## How scalable?

| Dataset | TPMs |  |  |  |  |  |  |  |  |  |  |  | DGMs |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | LVD (ours) | HCLT | EiNet | RAT-SPN |  | Glow | RealNVP | BIVA |  |  |  |  |  |  |  |
| ImageNet32 | $\mathbf{4 . 3 9}_{ \pm 0.01}$ | 4.82 | 5.63 | 6.90 |  | 4.09 | 4.28 | 3.96 |  |  |  |  |  |  |  |
| ImageNet64 | $\mathbf{4 . 1 2}_{ \pm 0.00}$ | 4.67 | 5.69 | 6.82 |  | 3.81 | 3.98 | - |  |  |  |  |  |  |  |
| CIFAR | $\mathbf{4 . 3 8}_{ \pm 0.02}$ | 4.61 | 5.81 | 6.95 |  | 3.35 | 3.49 | 3.08 |  |  |  |  |  |  |  |



## up to billions of parameters

## Questions answered today

1. What are probabilistic queries? Are current models tractable? (Guy)
2. What are probabilistic circuits and why are they tractable? (Guy)
3. What is the connection to logical circuit languages? (YooJung)
4. How do I compile my favorite model into a circuit? (YooJung)
5. How are circuit size and tractability related? (Yoolung)
6. What's the most impressive query we can efficiently compute? (YooJung)
7. Are all tractable distributions probabilistic circuits? (Guy)
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## Logical Circuits

## Tractability to other semi-rings

Tractable probabilistic inference exploits efficient summation for decomposable functions in the probability commutative semiring:

$$
(\mathbb{R},+, \times, 0,1)
$$

analogously efficient computations can be done in other semi-rings:
$\left(\mathbb{S}, \oplus, \otimes, 0_{\oplus}, 1_{\otimes}\right)$
$\Rightarrow$ Algebraic model counting (Kimmig et al. 2017), Semi-ring programming (Belle et al. 2016)
Historically, very well studied for boolean functions:

$$
(\mathbb{B}=\{0,1\}, \vee, \wedge, 0,1) \quad \Rightarrow \text { logical circuits! }
$$

## Logical circuits


s/d-D/NNFs
(Darwiche et al. 2002a)


O/BDDs
(Bryant 1986)


SDDs
(Darwiche 2011a)

Logical circuits are compact representations for boolean functions...

## Logical circuits

structural properties
...and like probabilitistic circuits, one can define structural properties: (structured) decomposability, smoothness, determinism allowing for tractable computations


## Logica/ circuits

a knowledge compilation map
...inducing a hierarchy of tractable logical circuit families


## Knowledge Compilation

encoding



## NNF Circuits

$$
\begin{aligned}
& P \vee L \\
& A \Rightarrow P \\
& K \Rightarrow(P \vee L) \\
& \hline
\end{aligned}
$$



## Decomposability (DNNF)

Darwiche, JACM 2001
SAT in linear time


## Determinism (d-DNNF)



## Decomposability + determinism $=$ tractable (W)MC

Model counting problem: given a Boolean formula $\Delta$, compute the number of satisfying assignments.

Weighted model counting (WMC):

$$
\operatorname{WMC}(\Delta, w)=\sum_{\mathbf{x} \models \Delta} \prod_{l \in \mathbf{x}} w(l)
$$

$\Rightarrow$ linear in circuit size!

## Decomposability + determinism $=$ tractable (W)MC

To compute $\mathrm{WMC}(\Delta, w)$ :

- Turn OR gates to sum nodes and AND gates to product nodes
- Renlace each literal l with its weight $\omega(l)$
- bottom-up evaluation



## Decomposability + determinism $=$ tractable $(W) M C$

To compute $\mathrm{WMC}(\Delta, w)$ :

- Turn OR gates to sum nodes and AND gates to product nodes
- Replace each literal $l$ with its weight $w(l)$
- bottom-up evaluation



## Probabilistic inference by WMC

connection to probabilistic circuits through WMC

1. Encode probabilistic model as WMC formula ( $\Delta, w)$
2. Compile $\Delta$ into a logical circuit (e.g. d-DNNF, OBDD, SDD, etc.)
3. Tractable MAR/CON by tractable WMC on circuit
4. Answer complex queries tractably by enforcing more structural properties!

## Probabilistic inference by WMC

connection to probabilistic circuits through WMC

Resulting compiled WMC circuit equivalent to probabilistic circuit
$\Rightarrow$ parameter variables $\rightarrow$ edge parameters


Compiled circuit of WMC encoding


Equivalent probabilistic circuit

## Questions answered today

1. What are probabilistic queries? Are current models tractable? (Guy)
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## From tree BN to circuits

via compilation


## From tree BN to circuits

via compilation
Bottom-up compilation: starting from leaves...


## From tree BN to circuits

via compilation
...compile a leaf CPT


## From tree BN to circuits

via compilation
...compile a leaf CPT


## From tree BN to circuits

## via compilation

...compile a leaf CPT...for all leaves...


## From tree BN to circuits

## via compilation

...and recurse over parents...


## From tree BN to circuits

## via compilation

...while reusing previously compiled nodes!...


## From tree BN to circuits

via compilation


## Hidden Markov Models

as computational graphs


## Compilation: probabilistic programming

```
```

x = flip( ( }\mp@subsup{|}{1}{\prime}\mathrm{ );

```
```

x = flip( ( }\mp@subsup{|}{1}{\prime}\mathrm{ );
if(x) {
if(x) {
y = flip( (
y = flip( (
} else {
} else {
y = x
y = x
}

```
```

    }
    ```
```



Chavira et al., "Compiling relational Bayesian networks for exact inference", 2006 Holtzen et al., "Symbolic Exact Inference for Discrete Probabilistic Programs", 2019 De Raedt et al.; Riguzzi; Fierens et al.; Vlasselaer et al., "ProbLog: A Probabilistic Prolog and Its Application in Link Discovery."; "A top down interpreter for LPAD and CP-logic"; "Inference and Learning in Probabilistic Logic Programs using Weighted Boolean Formulas"; "Anytime Inference in Probabilistic Logic Programs with Tp-compilation", 2007; 2007; 2015; 2015
Olteanu et al.; Van den Broeck et al., "Using OBDDs for efficient query evaluation on probabilistic databases"; Query Processing on Probabilistic Data: A Survey, 2008; 2017
Vlasselaer et al., "Exploiting Local and Repeated Structure in Dynamic Bayesian Networks", 2016

## Decision Diagrams

$\square$
FBDDs (Free binary
decision diagrams;
read-once)
OBDDs (Ordered BDDs)
SDDs (Sentential decision diagrams)

$\Rightarrow B D D$ as circuit

## Structured Decomposability

Pipatsrisawat \& Darwiche, AAAI 2008


## Structured Decomposability

Pipatsrisawat \& Darwiche, AAAI 2008


## Structured Decomposability

Pipatsrisawat \& Darwiche, AAAI 2008


## Partitioned Determinism (SDDs)

Darwiche, IJCAI 2011


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Darwiche, IJCAI 2011


## Decision Diagrams

$\square$
FBDDs (Free binary decision diagrams; read-once)
$\square$ OBDDs (Ordered BDDs)

- SDDs (Sentential decision diagrams)

$$
\Rightarrow \quad S D D \& O B D D \text { for }
$$

$$
(A \wedge B) \vee(C \wedge D)
$$


(a) vtree

(b) SDD

(c) OBDD

## Probability of logical events


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## Probability of logical events

$\mathrm{q}_{8}$ : What is the probability of having a traffic jam on my route to campus?

$$
\mathrm{q}_{8}(\mathbf{m})=p_{\mathbf{m}}\left(\bigvee_{i \in \text { route }} \operatorname{Jam}_{\operatorname{Str} i}\right)
$$

$$
\Rightarrow \text { marginals + logical events }
$$


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## Smoothness + structured decomp. = tractable PR

Computing $\boldsymbol{p}(\alpha)$ : the probability of arbitrary logical formula

Multilinear in circuit sizes if the logical circuit:

- is smooth, structured decomposable, deterministic
$\square$ shares the same vtree



## Smoothness + structured decomp. = tractable PR

If $p(\mathbf{x})=\sum_{i} w_{i} p_{i}(\mathbf{x}), \alpha=\bigvee_{j} \alpha_{j}$, (smooth $p$ )
(smooth + deterministic $\alpha$ ):

$$
p(\alpha)=\sum_{i} w_{i} p_{i}\left(\bigvee_{j} \alpha_{j}\right)=\sum_{i} w_{i} \sum_{j} p_{i}\left(\alpha_{j}\right)
$$

$\Rightarrow$ probabilities are "pushed down" to children


## Smoothness + structured decomp. $=$ tractable PR

If $p(\mathbf{x}, \mathbf{y})=p(\mathbf{x}) p(\mathbf{y}), \alpha=\beta \wedge \gamma$,
(structured decomposability):

$$
p(\alpha)=p(\beta \wedge \gamma) \cdot p(\beta \wedge \gamma)=p(\beta) \cdot p(\gamma)
$$

$\Rightarrow$ probabilities decompose into simpler ones


## Smoothness + structured decomp. = tractable PR

To compute $p(\alpha)$ :
compute the probability for each pair of probabilistic and logical circuit nodes for the same vtree node
$\Rightarrow \quad$ cache the values!

- feedforward evaluation (bottom-up)



## Smoothness + structured decomp. = tractable PR

To compute $p(\alpha)$ :
compute the probability for each pair of probabilistic and logical circuit nodes for the same vtree node
$\Rightarrow$ cache the values!
feedforward evaluation (bottom-up)


## structured decomposability $=$ tractable...

Symmetric and group queries (exactly-k, odd-number, etc.) (Bekker et al. 2015)
For the "right" vtree

- Marginal MAP (Oztok et al. 2016)Probability of logical circuit event in probabilistic circuit (Choi et al. 2015b)Multiply two probabilistic circuits (Shen et al. 2016)KL Divergence between probabilistic circuits (Liang et al. 2017)Same-decision probability (Oztok et al. 2016)
Expected same-decision probability (Choi et al. 2017)
- Expected classifier agreement (Choi et al. 2018)Expected predictions (Khosravi et al. 2019b)


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## Succinctness of circuits

Expressive efficiency
Tractability is defined with respect to the size of the model.
How do structural constraints affect expressive efficiency (succinctness) of probabilistic/logical circuits?


## Succinctness of circuits

Expressive efficiency

A family of circuits $\mathcal{M}_{1}$ is at least as succinct as $\mathcal{M}_{2}$
iff for every $\mathbf{m}_{2} \in \mathcal{M}_{2}$, there exists $\mathbf{m}_{1} \in \mathcal{M}_{1}$ that represents
the same function and $\left|m_{1}\right| \leq\left|\operatorname{poly}\left(m_{2}\right)\right|$.
$\Rightarrow \quad$ denoted $\mathcal{M}_{1} \leq \mathcal{M}_{2}$
$\Rightarrow$ strictly more succinct $\left(\mathcal{M}_{1}<\mathcal{M}_{2}\right)$

$$
\text { iff } \mathcal{M}_{1} \leq \mathcal{M}_{2} \text { and } \mathcal{M}_{1} \nsupseteq \mathcal{M}_{2}
$$

## Succinctness of circuits

Expressive efficiency
Strict succinctness ordering: DNNF < d-DNNF < FBDD < OBDD


## Succinctness of circuits

Expressive efficiency
Strict succinctness ordering: DNNF < d-DNNF < FBDD < OBDD
$\square$ d-DNNF $\not \leq$ DNNF unless the polynomial hierarchy collapses (Darwiche et al. 2002a).

- The Sauerhoff function has DNNF of size $O\left(n^{2}\right)$ but d-DNNF of size $2^{\Omega(n)}$ (Bova et a


## Succinctness of circuits

Expressive efficiency
Strict succinctness ordering: DNNF < d-DNNF < FBDD < OBDD
$\square$ d-DNNF $\not \leq$ DNNF unless the polynomial hierarchy collapses (Darwiche et al. 2002a).
The Sauerhoff function has DNNF of size $O\left(n^{2}\right)$ but d-DNNF of size $2^{\Omega(n)}$ (Bova et al. 2016).
$\Rightarrow$ Unconditional exponential separation for d-DNNF $\not \leq$ DNNF
$\Rightarrow$ Using a connection between circuits and communication complexity

## Succinctness of circuits

Expressive efficiency
SDD < OBDD: SDDs are strictly more succinct than OBDDsSDD $\leq$ OBDD: OBDDs are SDDs with right-linear vtrees
$\square$ SDD $\nsupseteq$ OBDD: The hidden weighted bit function has SDD of size $O\left(n^{3}\right)$ but OBDD of size $2^{\Omega}(n)$.


## Query compilation

## Möbius Über Alles



How precise is the characterization of tractable circuits by structural properties?

## Smoothness + decomposability $=$ tractable MAR

Recall: Smoothness and decomposability allow marginal inference by feedforward (sum-product) evaluation.


## Smoothness + decomposability $=$ tractable MAR

Recall: Smoothness and decomposability allow marginal inference by feedforward (sum-product) evaluation.
$\Rightarrow$ Are these properties necessary?


## Smoothness + decomposability $=$ tractable MAR

Recall: Smoothness and decomposability allow marginal inference by feedforward (sum-product) evaluation.
$\Rightarrow$ Are these properties necessary?
$\Rightarrow$ Yes! Otherwise, integrals do not decompose.


## Determinism + decomposability $=$ tractable MAP

Recall: Determinism and decomposability allow MAP inference by feedforward (max-product) evaluation.


## Determinism + decomposability $=$ tractable MAP

Recall: Determinism and decomposability allow MAP inference by feedforward (max-product) evaluation.
$\Rightarrow$ However, decomposability is not necessary!


## Determinism + decomposability $=$ tractable MAP

Recall: Determinism and decomposability allow MAP inference by feedforward (max-product) evaluation.
$\Rightarrow$ However, decomposability is not necessary!
$\Rightarrow$ A weaker condition, consistency, suffices.


## Consistency

A product node is consistent if any variable shared between its children appears in a single leaf node
$\Rightarrow$ decomposability implies consistency

consistent circuit

inconsistent circuit

## Determinism + consistency $=$ tractable MAP

## Determinism + consistency $=$ tractable MAP

If $\max _{\mathbf{q}_{\text {shared }}} p(\mathbf{q}, \mathbf{e})=$ $\max _{\mathbf{q}_{\text {shared }}} p\left(\mathbf{q}_{\mathbf{x}}, \mathbf{e}_{\mathbf{x}}\right) \cdot \max _{\mathbf{q}_{\text {shared }}} p\left(\mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}}\right)$ (consistent):

$$
\begin{aligned}
\max _{\mathbf{q}} p(\mathbf{q}, \mathbf{e}) & =\max _{\mathbf{q}_{\mathbf{x}}, \mathbf{q}_{\mathbf{y}}} p\left(\mathbf{q}_{\mathbf{x}}, \mathbf{e}_{\mathbf{x}}, \mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}}\right) \\
& =\max _{\mathbf{q}_{\mathbf{x}}} p\left(\mathbf{q}_{\mathbf{x}}, \mathbf{e}_{\mathbf{x}}\right) \cdot \max _{\mathbf{q}_{\mathbf{y}}} p\left(\mathbf{q}_{\mathbf{y}}, \mathbf{e}_{\mathbf{y}}\right) \\
& \Rightarrow \text { solving optimization independently }
\end{aligned}
$$



## Expressive efficiency of circuits

MAR
smooth \& Decomp.
det. \& cons.
MAP

Are smooth \& decomposable circuits as succinct as deterministic \& consistent ones, or vice versa?

## Expressive efficiency of circuits



- Smooth \& decomposable circuits strictly more succinct than deterministic \& decomposable ones

Smooth \& consistent circuits are equally succinct as smooth \& decomposable ones

## Expressive efficiency of circuits



## Expressive efficiency of circuits


$\longrightarrow$ : strictly more succinct
工 : equally succinct

## Expressive efficiency of circuits



工_ : equally succinct

## Expressive efficiency of circuits


$\longrightarrow$ : strictly more succinct
工 : equally succinct

## Expressive efficiency of circuits


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工 : equally succinct

## Expressive efficiency of circuits


$\longrightarrow$ : strictly more succinct
工 : equally succinct

## Expressive efficiency of circuits


$\longrightarrow$ : strictly more succinct
工 : equally succinct

## Expressive efficiency of circuits


$\square$ Neither smooth \& decomposable nor deterministic \& consistent circuits are more succinct than the other!
$\Rightarrow$ Choose tractable circuit family based on your query
$\square$ More theoretical questions remaining
$\Rightarrow$ "Complete the map"
$\longrightarrow$ : strictly more succinct
工 : equally succinct

## Expressive efficiency of circuits



Succinctness map for monotone circuits
$\Rightarrow$ (s)mooth, (d)eterministic, (D)ecomposable, (w)eak (D)ecomposable (i.e. consistent)

## Expressive efficiency of circuits



Succinctness map for monotone circuits


Succinctness map for positive circuits (non-negative output, but weights may be negative)
$\Rightarrow$ (s)mooth, (d)eterministic, (D)ecomposable, (w)eak (D)ecomposable (i.e. consistent)

## Questions answered today

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## Gocl

Given a class of queries can we systematically find a class of probabilistic circuits that is tractable for it?

## A language for queries

Integral expressions that can be formed by composing these operators

$$
+, \times, \text { pow, log, exp and } /
$$

$\Rightarrow$ many divergences and information-theoretic queries

## A language for queries

Integral expressions that can be formed by composing these operators

$\Rightarrow$ many divergences and information-theoretic queries

Represented as higher-order computational graphs—pipelines—operating over circuits! $\Rightarrow$ re-using intermediate transformations across queries

## $\mathbb{K L L D}(p \| q)=\int_{\text {val }(\mathbf{X})} p(\mathbf{x}) \times \log (p(\mathbf{x}) / q(\mathbf{x})) d \mathbf{X}$



$$
\mathbb{K} \mathbb{L D}(p \| q)=\int_{\text {val }(\mathbf{X})} p(\mathbf{x}) \times \log (p(\mathbf{x}) / q(\mathbf{x})) d \mathbf{X}
$$



## $\mathbb{K} \mathbb{L D}(p \| q)=\int_{\text {val }(\mathbf{X})} p(\mathbf{x}) \times \log (p(\mathbf{x}) / q(\mathbf{x})) d \mathbf{X}$



## $\mathbb{K L L D}(p \| q)=\int_{\text {val }(\mathbf{X})} p(\mathbf{x}) \times \log (p(\mathbf{x}) / q(\mathbf{x})) d \mathbf{X}$



## $\mathbb{X E N T}(p \| q)=\int p(\mathbf{x}) \times \log q(\mathbf{x}) d \mathbf{X}$



$$
\mathbb{E}_{\mathbf{x}^{m} \sim p\left(\mathbf{x}^{m} \mid \mathbf{x}^{o}\right)}\left[q^{\alpha}\left(\mathbf{x}^{m}, \mathbf{x}^{o}\right)\right]
$$



## Compatibility

Two circuits are compatible if they have the same hierarchical scope partitioning $\Rightarrow$ generalizes "structured decomposability with same vtree"

compatible circuits

non-compatible circuits

## Tractable operators


smooth, decomposable compatible

## Tractable operators


smooth, decomposable deterministic


Building an atlas of composable tractable atomic operations


To perform tractable integration we need $s$ to be smooth and decomposable...

hence we need $p$ and $r$ to be smooth, decomposable and compatible...

therefore $q$ must be smooth, decomposable and deterministic...

we can compute $\mathbb{X} \mathbb{E} \mathbb{N} \mathbb{T}$ tractably if $\boldsymbol{p}$ and $\boldsymbol{q}$ are smooth, decomposable, compatible and $\boldsymbol{q}$ is deterministic...

|  | Query | Tract. Conditions | Hardness |
| :---: | :---: | :---: | :---: |
| Cross Entropy | $-\int p(\boldsymbol{x}) \log q(\boldsymbol{x}) \mathrm{d} \mathbf{X}$ | Cmp, q Det | \#P-hard w/o Det |
| Shannon Entropy | $-\sum p(x) \log p(\boldsymbol{x})$ | Sm, Dec, Det | coNP-hard w/o Det |
| RÉNyi Entropy | $(1-\alpha)^{-1} \log \int p^{\alpha}(\boldsymbol{x}) d \mathbf{X}, \alpha \in \mathbb{N}$ | SD | \#P-hard w/o SD |
|  | $(1-\alpha)^{-1} \log \int p^{\alpha}(x) d \mathbf{X}, \alpha \in \mathbb{R}_{+}$ | Sm, Dec, Det | \#P-hard w/o Det |
| MUTUAL InFormation | $\int p(\boldsymbol{x}, \boldsymbol{y}) \log (p(\boldsymbol{x}, \boldsymbol{y}) /(p(\boldsymbol{x}) p(\boldsymbol{y}))$ ) | Sm, SD, Det* | coNP-hard w/o SD |
| Kullback-Leibler Div. | $\int p(\boldsymbol{x}) \log (p(\boldsymbol{x}) / q(\boldsymbol{x})) d \mathbf{X}$ | Cmp, Det | \#P-hard w/o Det |
| RÉNYI'S ALPHA DIV. | $(1-\alpha)^{-1} \log \int p^{\alpha}(\boldsymbol{x}) q^{1-\alpha}(\boldsymbol{x}) d \mathbf{X}, \alpha \in \mathbb{N}$ | Cmp, $q$ Det | \#P-hard w/o Det |
|  | $(1-\alpha)^{-1} \log \int p^{\alpha}(\boldsymbol{x}) q^{1-\alpha}(\boldsymbol{x}) d \mathbf{X}, \alpha \in \mathbb{R}$ | Cmp, Det | \#P-hard w/o Det |
| ITAKURA-SAITO DIV. | $\int[p(\boldsymbol{x}) / q(\boldsymbol{x})-\log (p(x) / q(\boldsymbol{x}))-1] d \mathbf{X}$ | Cmp, Det | \#P-hard w/o Det |
| CaUchy-Schwarz Div. | $-\log \frac{\int p(\boldsymbol{x}) q(\boldsymbol{x}) d \mathbf{X}}{\sqrt{\int p^{2}(\boldsymbol{x}) d \mathbf{X} q^{2}(\boldsymbol{x}) d \mathbf{X}}}$ | Cmp | \#P-hard w/o Cmp |
| SQUARED LOSS | $\int(p(\boldsymbol{x})-q(\boldsymbol{x}))^{2} d \mathbf{X}$ | Cmp | \#P-hard w/o Cmp |

## compositionally derive the tractability of many more queries

|  | Query | Tract. Conditions | Hardness |
| :---: | :---: | :---: | :---: |
| Cross Entropy | $-\int p(\boldsymbol{x}) \log q(\boldsymbol{x}) \mathrm{d} \mathbf{X}$ | Cmp, q Det | \#P-hard w/o Det |
| Shannon Entropy | $-\sum p(x) \log p(\boldsymbol{x})$ | Sm, Dec, Det | coNP-hard w/o Det |
| RÉNyi Entropy | $(1-\alpha)^{-1} \log \int p^{\alpha}(\boldsymbol{x}) d \mathbf{X}, \alpha \in \mathbb{N}$ | SD | \#P-hard w/o SD |
|  | $(1-\alpha)^{-1} \log \int p^{\alpha}(x) d \mathbf{X}, \alpha \in \mathbb{R}_{+}$ | Sm, Dec, Det | \#P-hard w/o Det |
| MUTUAL InFormation | $\int p(\boldsymbol{x}, \boldsymbol{y}) \log (p(\boldsymbol{x}, \boldsymbol{y}) /(p(\boldsymbol{x}) p(\boldsymbol{y}))$ ) | Sm, SD, Det* | coNP-hard w/o SD |
| Kullback-Leibler Div. | $\int p(\boldsymbol{x}) \log (p(\boldsymbol{x}) / q(\boldsymbol{x})) d \mathbf{X}$ | Cmp, Det | \#P-hard w/o Det |
| RÉNYI'S ALPHA DIV. | $(1-\alpha)^{-1} \log \int p^{\alpha}(\boldsymbol{x}) q^{1-\alpha}(\boldsymbol{x}) d \mathbf{X}, \alpha \in \mathbb{N}$ | Cmp, $q$ Det | \#P-hard w/o Det |
|  | $(1-\alpha)^{-1} \log \int p^{\alpha}(\boldsymbol{x}) q^{1-\alpha}(\boldsymbol{x}) d \mathbf{X}, \alpha \in \mathbb{R}$ | Cmp, Det | \#P-hard w/o Det |
| ITAKURA-SAITO DIV. | $\int[p(\boldsymbol{x}) / q(\boldsymbol{x})-\log (p(x) / q(\boldsymbol{x}))-1] d \mathbf{X}$ | Cmp, Det | \#P-hard w/o Det |
| CaUchy-Schwarz Div. | $-\log \frac{\int p(\boldsymbol{x}) q(\boldsymbol{x}) d \mathbf{X}}{\sqrt{\int p^{2}(\boldsymbol{x}) d \mathbf{X} q^{2}(\boldsymbol{x}) d \mathbf{X}}}$ | Cmp | \#P-hard w/o Cmp |
| SQUARED LOSS | $\int(p(\boldsymbol{x})-q(\boldsymbol{x}))^{2} d \mathbf{X}$ | Cmp | \#P-hard w/o Cmp |

## and prove hardness when some input properties are not satisfied

## Composable tractable sub-routines

```
function kld(p, q)
    r = quotient(p, q)
    s}=\operatorname{log}(r
    t = product(p,s)
    return integrate(t)
end
```

function $x \in n t(p, q)$
$r=\log (q)$
$s=$ product ( $p, r)$
return -integrate (s)
$s=\operatorname{product}(p, r)$
return -integrate (s)
end

```
function alphadiv(p, q, alpha=1.5)
    r = real_pow(p, alpha)
    s = real_pow(q, 1.0-alpha)
    t = product(r,s)
    return log(integrate(t)) / (1.0-alpha)
end
```

Efficient inference algorithms in a couple lines of Julia code! ${ }^{2}$

[^5]
## Next up...

1. Learning and reasoning with symbolic constraints
2. Expected predictions: handling missing values, fairness
3. Exact inference of causal effects
$\Rightarrow$ using tractable operators

smooth, decomposable compatible

## Symbolic constraints

"How can neural nets
reason and learn with
symbolic constraints
reliably and efficiently?"

## When?



Ground Truth

## e.g. predict shortest path in a map

## When?



Ground Truth
given X // e.g. a tile map

## structured output prediction (SOP) tasks

## When?



Ground Truth
given $\mathbf{x}$ // e.g. a tile map
find $\mathbf{y}^{*}=\operatorname{argmax}_{\mathbf{y}} p_{\theta}(\mathbf{y} \mid \mathbf{x}) / /$ e.g. a configurations of edges in a grid

## structured output prediction (SOP) tasks

## When?



Ground Truth
given $\mathbf{x}$ // e.g. a tile map
find $\mathbf{y}^{*}=\operatorname{argmax}_{\mathbf{y}} p_{\theta}(\mathbf{y} \mid \mathbf{x}) / /$ e.g. a configurations of edges in a grid s.t. $\mathbf{y}=\mathrm{K} / /$ e.g., that form a valid path

## structured output prediction (SOP) tasks

## When?



Ground Truth
given $\mathbf{x}$ // e.g. a tile map
find $\mathbf{y}^{*}=\operatorname{argmax}_{\mathbf{y}} p_{\theta}(\mathbf{y} \mid \mathbf{x}) / /$ e.g. a configurations of edges in a grid s.t. $\mathbf{y}=\mathrm{K} / /$ e.g., that form a valid path
// for a $12 \times 12$ grid, $2^{144}$ states but only $10^{10}$ valid ones!

## structured output prediction (SOP) tasks

## When?



$$
\begin{aligned}
& \text { given } \mathbf{x} / / \text { e.g. a feature map } \\
& \text { find } \mathbf{y}^{*}=\operatorname{argmax} \mathrm{max}_{\boldsymbol{y}}(\mathbf{y} \mid \mathbf{x}) / / \text { e.g. labels of classes } \\
& \qquad \text { s.t. } \mathbf{y} \vDash \mathrm{K} / / \text { e.g., constraints over superclasses }
\end{aligned}
$$

$\mathrm{K}:\left(Y_{\text {cat }} \Longrightarrow Y_{\text {animal }}\right) \wedge\left(Y_{\text {dog }} \Longrightarrow Y_{\text {animal }}\right)$

## hierarchical multi-Iabel classification

## When?



Ground Truth


ResNet-18
neural nets struggle to satisfy domain constraints!

## How?


take an unreliable neural network architecture...

......and replace the last layer with a semantic probabilistic Iayer

$q_{\boldsymbol{\Theta}}(\mathbf{y} \mid g(\mathbf{z}))$ is an expressive distribution over labels

$$
c_{\mathrm{K}}(\mathbf{x}, \mathbf{y}) \text { encodes the constraint } \mathbb{1}\{\mathbf{x}, \mathbf{y} \models \mathrm{K}\}
$$



$$
p(\mathbf{y} \mid \mathbf{x})=q_{\boldsymbol{\Theta}}(\mathbf{y} \mid g(\mathbf{z})) \cdot c_{\mathbf{K}}(\mathbf{x}, \mathbf{y}) / \mathcal{Z}(\mathbf{x})
$$

$$
\mathcal{Z}(\mathbf{x})=\sum_{\mathbf{y}} q_{\boldsymbol{\Theta}}(\mathbf{y} \mid \mathbf{x}) \cdot c_{\mathrm{K}}(\mathbf{x}, \mathbf{y})
$$


a conditional circuit $q(\mathbf{y} ; \boldsymbol{\Theta}=g(\mathbf{z}))$

and a logical circuit $c(y, x)$ encoding $K$

## Tractable products


exactly compute $\mathcal{Z}$ in time $O(|q||c|)$

## SPL recipe

$$
\begin{aligned}
& \mathrm{K}:\left(Y_{1}=1 \Longrightarrow Y_{3}=1\right) \\
& \wedge\left(Y_{2}=1 \Longrightarrow Y_{3}=1\right)
\end{aligned}
$$

1) Take any
logical constraint

## SPL recipe



1) Take any
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2) Compile it into
a constraint circuit

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## SPL recipe

$\mathrm{K}:\left(Y_{1}=1 \Longrightarrow Y_{3}=1\right)$
$\wedge\left(Y_{2}=1 \Longrightarrow Y_{3}=1\right)$


1) Take any
logical constraint
2) Compile it into
a constraint circuit

3) Multiply it by a circuit distribution

## 4) train end-to-end by sgd!

## Guaranteeing consistency


cost: 39.31

cost: 57.31

cost: $\infty$

cost: $\infty$
$\mathcal{L}_{\mathrm{SL}}$

cost: $\infty$

cost: $\infty$

SPL

cost: 45.09

cost: 58.09

## Expected predictions

Reasoning about the output of a classifier or regressor $\boldsymbol{f}$ given a distribution $\boldsymbol{p}$ over the input features

$$
\mathbb{E}_{p}[f]=\int_{\operatorname{val}(\mathbf{X})} p(\mathbf{x}) \times f(\mathbf{x}) d \mathbf{X}
$$



## Handling missing values at test time



Given a partial observation $\mathbf{x}^{o}$, what is the expected output from $f$ ?

$$
\underset{\mathbf{x}^{m} \sim p\left(\mathbf{x}^{m} \mid \mathbf{x}^{o}\right)}{\mathbb{E}}\left[f\left(\mathbf{x}^{m}, \mathbf{x}^{o}\right)\right]
$$

## Fairness analysis

```
using ProbabilisticCircuits
pc = load_prob_circuit(zoo_psdd_file("insurance.psdd"));
rc = load_logistic_circuit(zoo_lc_file("insurance.circuit"), 1);
```

q : Is the predictive model biased by gender?

```
groups = make_observations([["male"], ["female"]])
exps, _ = Expectation(pc, rc, groups);
println("Female : \\$ \$(exps[2])");
println("Male : \\$ \$(exps[1])");
println("Diff : \\$ \$(exps[2] - exps[1])");
Female : \$ 14170.125469335406
Male : \$ 13196.548926381849
Diff : \$ 973.5765429535568
```


## Causal Inference

Given subsets $\boldsymbol{A}, \boldsymbol{Y} \subseteq \boldsymbol{X}$, interested in causal effect $p(\boldsymbol{Y} \mid \operatorname{do}(\boldsymbol{A}))$.

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- Specify (qualitative) assumptions on the system using a causal diagram $G($ here $\boldsymbol{A}, \boldsymbol{Y}, \boldsymbol{Z}, \boldsymbol{K} \subseteq \boldsymbol{X})$ ) :

(a) Backdoor

(b) Napkin


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- Specify (qualitative) assumptions on the system using a causal diagram $G($ here $\boldsymbol{A}, \boldsymbol{Y}, \boldsymbol{Z}, \boldsymbol{K} \subseteq \boldsymbol{X})$ ) :

(a) Backdoor

(b) Napkin
- Given causal diagram $G$, can derive expressions for causal effect $p(\boldsymbol{Y} \mid \boldsymbol{A})$ using do-calculus (Pearl 1995).

$$
\sum_{\boldsymbol{Z}} p(\boldsymbol{Z}) p(\boldsymbol{Y} \mid \boldsymbol{A}, \boldsymbol{Z})
$$

$$
\frac{\sum_{K} p(\boldsymbol{A}, \boldsymbol{Y} \mid K, \boldsymbol{Z}) p(K)}{\sum_{K} p(\boldsymbol{A} \mid K, \boldsymbol{Z}) p(K)}
$$

(a) Backdoor
(b) Napkin

## Tractability of Exact Causal Inference

Consider the backdoor query, for fixed values of the treatment a and outcome $\boldsymbol{y}$ :

$$
p(\boldsymbol{y} \mid d o(\boldsymbol{a})):=\sum_{\boldsymbol{Z}} p(\boldsymbol{Z}) \times p(\boldsymbol{y} \mid \boldsymbol{a}, \boldsymbol{Z})
$$

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$$

## Theorem (Wang \& Kwiatkowska 2023)

If $p$ is given as a structured decomposable and deterministic circuit, then the backdoor query is \#P-hard to compute.

## Applying the Atlas of Tractable Operations

Break down do-calculus query into compositions of basic operations, such as marginalization, products, and powers:

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Break down do-calculus query into compositions of basic operations, such as marginalization, products, and powers:


Problem: Cannot guarantee that input to POW is deterministic, even if $p(\boldsymbol{X})$ is deterministic.

## Marginal Determinism

## Definition (Marginal Determinism, Choi et al. 2020)

Given a subset of variables $\boldsymbol{Q} \subseteq \boldsymbol{X}$, a PC is $\boldsymbol{Q}$-deterministic if the children of a sum node $T$ correspond to different values of $\boldsymbol{Q}$ (for sum nodes with sc $(T) \cap \boldsymbol{Q} \neq \emptyset)$.

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(a) $\boldsymbol{Q}=\{A, Z\}$-deterministic
(b) $\boldsymbol{Q}=\{A, Y\}$-deterministic

Motivation: If a circuit is marginally deterministic w.r.t $\boldsymbol{Q}$, then we can marginalize out $\boldsymbol{X} \backslash \boldsymbol{Q}$ and obtain a deterministic circuit!

## Tractable Causal Inference

If (the circuit encoding) $p(\boldsymbol{X})$ is $(\boldsymbol{A} \cup \boldsymbol{Z})$-deterministic, then the input to POW is guaranteed to be deterministic.

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(b) Pipeline for entire backdoor query

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(a) Pipeline for $\operatorname{COND}(\cdot, \boldsymbol{A} \cup \boldsymbol{Z})$

(b) Pipeline for entire backdoor query
$\Longrightarrow$ all operations are tractable according to Atlas
$\Longrightarrow$ can compute causal effect in $O\left(|p|^{3}\right)$ time
(can improve to $O\left(|p|^{2}\right)$ )

## Open Questions

- Are all causal queries derived by the do-calculus tractable in PTIME (for some non-trivial marginal determinism condition)?
- What is the optimal complexity for these queries?


## Questions answered today

1. What are probabilistic queries? Are current models tractable? (Guy)
2. What are probabilistic circuits and why are they tractable? (Guy)
3. What is the connection to logical circuit languages? (YooJung)
4. How do I compile my favorite model into a circuit? (YooJung)
5. How are circuit size and tractability related? (YooJung)
6. What's the most impressive query we can efficiently compute? (YooJung)
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## Questions answered today

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8. How to learn probabilistic circuits from data? (Guy)
more tractable queries


## tractability vs expressive efficiency

## Smooth $V$ decomposable $V$ deterministic <br> $\checkmark$ structured decomposable PCs?

|  | smooth | dec. | det. |
| ---: | ---: | ---: | ---: |
| str.dec. |  |  |  |
| Arithmetic Circuits (ACs) (Darwiche 2003) |  |  |  |
| Sum-Product Networks (SPNs) (Poon et al. 2011) |  |  |  |
| Cutset Networks (CNets) (Rahman et al. 2014) |  |  |  |
| Probabilistic Decision Graphs (Jaeger 2004) |  |  |  |
| (Affine) ADDs (Hoey et al. 1999; Sanner et al. 2005) |  |  |  |
| AndOrGraphs (Dechter et al. 2007) |  |  |  |
| PSDDs (Kisa et al. 2014) |  |  |  |

## Low-treewidh PGMs

| Tree, polytrees and | Therefore they support |
| :--- | :---: |
| Thin Junction trees | tractable |
| can be turned into | EVI |
| $\square$ decomposable | MAR/CON |
| $\square$ smooth | MAP |
| $\square$ deterministic |  |
| circuits |  |



## Arithmetic Circuits (ACs)

ACs (Darwiche 2003) are
$\square$ decomposable
$\square$ smooth
$\square$ deterministic

They support tractable
EVI

- MAR/CON
- MAP

parameters are attached to the leaves $\Rightarrow \quad$...but can be moved to the sum node edges (Rooshenas et al. 2014)


## Sum-Product Networks (SPNs)


$\Rightarrow$ deterministic SPNs are also called selective (Peharz et al. 2014)

## Cutset Networks (CNets)

## CNets

(Rahman et al. 2014) are
decomposable
smooth
$\square$ deterministic

They support tractable
EVI
MAR/CON
MAP


Rahman et al., "Cutset Networks: A Simple, Tractable, and Scalable Approach for Improving the Accuracy of Chow-Liu Trees", 2014
Di Mauro et al., "Learning Accurate Cutset Networks by Exploiting Decomposability", 2015

## Probabilistic Sentential Decision Diagrams

PSDDs (Kisa et al. 2014) are
$\square$ structured decomposable
$\square$ smooth
deterministic

They support tractable
EVI
MAR/CON

- MAP

Complex queries!


[^6]
## Probabilistic Decision Graphs

PDGs (Jaeger 2004) are
structured decomposable
smooth
$\square$ deterministic

They support tractable
EVI
MAR/CON
MAP
Complex queries!


Jaeger, "Probabilistic decision graphs-combining verification and AI techniques for probabilistic inference", 2004
Jaeger et al., "Learning probabilistic decision graphs", 2006

## AndOrGraphs

AndOrGarphs
(Dechter et al. 2007) are
structured
decomposable
smooth
deterministic

They support tractable
EVI
MAR/CON
MAP
Complex queries!


Probabilistic circuits seem awfully general.

## Are all tractable probabilistic models probabilistic circuits?

## Enter: Determinantal Point Processes (DPPs)

DPPs are models where probabilities are specified by (sub)determinants

$$
\begin{aligned}
& L=\left[\begin{array}{cccc}
1 & 0.9 & 0.8 & 0 \\
0.9 & 0.97 & 0.96 & 0 \\
0.8 & 0.96 & 1 & 0 \\
0 & 0 & 0 & 1
\end{array}\right] \\
& \operatorname{Pr}_{L}\left(X_{1}=1, X_{2}=0, X_{3}=1, X_{4}=0\right)=\frac{1}{\operatorname{det}(L+I)} \operatorname{det}\left(L_{\{1,2\}}\right)
\end{aligned}
$$

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0.8 & 0.96 & 1 & 0 \\
0 & 0 & 0 & 1
\end{array}\right]
$$

Tractable likelihoods and marginals
$\square$
Global Negative Dependence

Diversity in recommendation systems

$$
\operatorname{Pr}_{L}\left(X_{1}=1, X_{2}=0, X_{3}=1, X_{4}=0\right)=\frac{1}{\operatorname{det}(L+I)} \operatorname{det}\left(L_{\{1,2\}}\right)
$$

## Are all tractable probabilistic models probabilistic circuits?



## Relationship between PCs and DPPs



## We cannot tractably represent DPPs with subclasses of PCs



## We cannot tractably represent DPPs with subclasses of PCs



## We cannot tractably represent DPPs with subclasses of PCs



## We cannot tractably represent DPPs with subclasses of PCs



## PCs and Circuit Lower Bounds

Theorem (Martens and Medabalimi, 2014). Let $P_{n}$ be the uniform distribution over spanning trees on $K_{n}$. For $n \geq 20$, the size of any smooth and decomposable $P C$ that represents $P_{n}$ is at least $2^{n / 30240}$.

Based on arithmetic circuit lower bounds by Ran Raz and Amir Yehudayoff

## Decomposable PCs are Syntactically Multilinear Arithmetic Circuits:

Definition 7 (Multilinear Arithmetic Circuit) If every node of an arithmetic circuit $\Phi$ over $y$ computes a multilinear polynomial in $y, \Phi$ is said to be a (semantically) multilinear arithmetic circuit. And if for every product node in $\Phi$, the scopes of its child nodes are pair-wise disjoint, $\Phi$ is said to be a syntactically multilinear arithmetic circuit.

## DPPs have No Compact Decomposable PCs

Theorem (Snell, 1995). The uniform distribution over spanning trees on the complete graph $K_{n}$ is a DPP over $\binom{n}{2}$ edges.

Theorem (Martens and Medabalimi, 2014). Let $P_{n}$ be the uniform distribution over spanning trees on $K_{n}$. For $n \geq 20$, the size of any smooth and decomposable $P C$ that represents $P_{n}$ is at least $2^{n / 30240}$.


## Probabilistic Generating Circuits



A Tractable Unifying Framework for PCs and DPPs

## Probability Generating Functions

| $X_{1}$ | $X_{2}$ | $X_{3}$ | $\operatorname{Pr}_{\beta}$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0.02 |
| 0 | 0 | 1 | 0.08 |
| 0 | 1 | 0 | 0.12 |
| 0 | 1 | 1 | 0.48 |
| 1 | 0 | 0 | 0.02 |
| 1 | 0 | 1 | 0.08 |
| 1 | 1 | 0 | 0.04 |
| 1 | 1 | 1 | 0.16 |

$$
\begin{aligned}
g_{\beta}= & \underbrace{0.16 z_{1} z_{2} z_{3}}+0.04 z_{1} z_{2}+0.08 z_{1} z_{3}+0.02 z_{1} \\
& +0.48 z_{2} z_{3}+0.12 z_{2}+0.08 z_{3}+0.02 .
\end{aligned}
$$

## Probability Generating Functions

| $X_{1}$ | $X_{2}$ | $X_{3}$ | $\operatorname{Pr}_{\beta}$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0.02 |
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& +0.48 z_{2} z_{3}+0.12 z_{2}+0.08 z_{3}+0.02
\end{aligned}
$$

## Probabilistic Generating Circuits (PGCs)

$$
g_{\beta}=\left(0.1\left(z_{1}+1\right)\left(6 z_{2}+1\right)-0.4 z_{1} z_{2}\right)\left(0.8 z_{3}+0.2\right)
$$



1. Sum nodes $\bigoplus$ with weighted edges to children.
2. Product nodes $\bigotimes$ with unweighted edges to children.
3. Leaf nodes: z_i or constant.

## PCs as PGCs

(Smooth \& Decomposable) PCs represents probability mass functions:

$$
\begin{aligned}
m_{\beta}= & 0.16 X_{1} X_{2} X_{3}+0.04 X_{1} X_{2} \overline{X_{3}}+0.08 X_{1} \overline{X_{2}} X_{3}+0.02 X_{1} \overline{X_{2}} \overline{X_{3}} \\
& +0.48 \overline{X_{1}} X_{2} X_{3}+0.12 \overline{X_{1}} X_{2} \overline{X_{3}}+0.08 \overline{X_{1}} \overline{X_{2}} X_{3}+0.02 \overline{X_{1}} \overline{X_{2}} \overline{X_{3}}
\end{aligned}
$$

PGCs represent probability generating functions:

$$
\begin{aligned}
g_{\beta}= & 0.16 z_{1} z_{2} z_{3}+0.04 z_{1} z_{2}+0.08 z_{1} z_{3}+0.02 z_{1} \\
& +0.48 z_{2} z_{3}+0.12 z_{2}+0.08 z_{1} z_{3}+0.02
\end{aligned}
$$

Given a smooth \& decomposable PC, by setting $\overline{X_{i}}$ to 1 , and $X_{i}$ to $z_{i}$, we obtain a PGC that represents the PC.

## Tractable Likelihood (EVID) or Marginals (MAR)?



## PGCs Support Tractable Likelihoods/Marginals



## PGCs Support Tractable Likelihoods/Marginals

$$
z_{i}=\left\{\begin{array}{lr}
t . & X_{i}=1 \\
0, & X_{i}=0 \\
1, & \text { otherwise }
\end{array}\right.
$$



$$
\operatorname{Pr}\left(X_{1}=1, X_{2}=0, \ldots\right)=?
$$



- Monomials setting to true variables that must be false are 0-ed out
- Other monomials contribute to result.
- Only monomials that set all required variables to true have max degree.
- Sum those up


## PGCs Support Tractable Likelihoods/Marginals



## Example



## Example



## Example



| $X_{1}$ | $X_{2}$ | $X_{3}$ | $\operatorname{Pr}_{\beta}$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0.02 |
| 0 | 0 | 1 | 0.08 |
| 0 | 1 | 0 | 0.12 |
| 0 | 1 | 1 | 0.48 |
| 1 | 0 | 0 | 0.02 |
| 1 | 0 | 1 | 0.08 |
| 1 | 1 | 0 | 0.04 |
| 1 | 1 | 1 | 0.16 |

## Inference Time Complexity

Given a PGC of size $m$ (\#edges) over $n$ random variables.
Algorithm 1 (Zhang et al., ICML 2021):

| Bottom-up pass <br> $\mathrm{w} / \mathrm{z} \_\mathrm{i}=\mathrm{t}, 0$ or 1 |
| :--- |
| Product/sum of degree-n <br> polynomials at each node |
|  |
| or $O(m n \log \mathrm{n} \log \log \mathrm{n})$ |

## Inference Time Complexity

Given a PGC of size $m$ (\#edges) over $n$ random variables.
Algorithm 1 (Zhang et al., ICML 2021):


| Product/sum of degree-n |
| :--- |
| polynomials at each node |$\rightleftharpoons O\left(m n^{2}\right)$

or $O(m n \log n \log \log n)$
Algorithm 2 (Harviainen et al., UAI 2023):
observation: the output of a PGC is a degree-n polynomial w/ respect to $t$

$$
\left.\begin{array}{c}
\text { Bottom-up pass } \\
\mathrm{w} / \mathrm{t}=0,1, \ldots, \mathrm{n}
\end{array}\right) \quad O(m n)
$$

## Syntactic vs. Semantic Restrictions

+ PGCs are tractable when semantically multilinear
+ No need for PC decomposability/syntactic multilinearity or other properties...
- Checking Validity of PGCs is Hard

Theorem (Harviainen et al.). It is NP-hard to check if a PGC encodes a valid probability generating polynomial

## DPPs as PGCs

The generating polynomial for a DPP with kernel $L$ is given by:

$$
g_{L}=\frac{1}{\operatorname{det}(L+I)} \operatorname{det}\left(I+L \operatorname{diag}\left(z_{1}, \ldots, z_{n}\right)\right)
$$

We need it as a sum of products to obtain a Probabilistic Generating Circuit

## DPPs as PGCs

The generating polynomial for a DPP with kernel $L$ is given by:

$$
g_{L}=\frac{1}{\operatorname{det}(L+I)} \operatorname{det}\left(I+L \operatorname{diag}\left(z_{1}, \ldots, z_{n}\right)\right) .
$$

Constant

We need it as a sum of products to obtain a Probabilistic Generating Circuit

## DPPs as PGCs

The generating polynomial for a DPP with kernel $L$ is given by:


## Experiment Results: Amazon Baby Registries

|  | DPP | Strudel | EiNet | MT | SimplePGC |
| :---: | :---: | :---: | :---: | :---: | :---: |
| apparel | -9.88 | -9.51 | -9.24 | -9.31 | $-9.10{ }^{*+\circ}$ |
| bath | -8.55 | -8.38 | -8.49 | -8.53 | $-8.29{ }^{* \dagger}{ }^{\text {a }}$ |
| bedding | -8.65 | -8.50 | -8.55 | -8.59 | $-8.41^{* \dagger}{ }^{\text {a }}$ |
| carseats | -4.74 | -4.79 | -4.72 | -4.76 | $-4.64{ }^{* \dagger}$ |
| diaper | -10.61 | -9.90 | -9.86 | -9.93 | $-9.72{ }^{* \dagger}$ |
| feeding | -11.86 | -11.42 | -11.27 | -11.30 | $-11.17^{*+\circ}$ |
| furniture | -4.38 | -4.39 | -4.38 | -4.43 | $-4.34^{* \dagger}$ |
| gear | -9.14 | -9.15 | -9.18 | $-9.23$ | $-9.04{ }^{* \dagger}{ }^{\text {a }}$ |
| gifts | -3.51 | -3.39 | -3.42 | -3.48 | $-3.47^{\circ}$ |
| health | -7.40 | -7.37 | -7.47 | -7.49 | $-7.24{ }^{* \dagger}$ |
| media | -8.36 | -7.62 | -7.82 | $-7.93$ | $-7.69{ }^{\text {¢o }}$ |
| moms | -3.55 | -3.52 | -3.48 | $-3.54$ | $-3.53{ }^{\circ}$ |
| safety | -4.28 | -4.43 | -4.39 | -4.36 | $-4.28{ }^{*+}$ |
| strollers | -5.30 | -5.07 | -5.07 | -5.14 | $-5.00^{* \dagger}$ |
| toys | -8.05 | -7.61 | -7.84 | $-7.88$ | $-7.62{ }^{\dagger}$ |

## Beyond DPPs: Strongly Rayleigh Distributions

## DPPs are strongly Rayleigh distributions

Definition. A probability distribution over binary random variables $X_{1}, \ldots, X_{n}$ (or equivalently, subsets of $[n]:=\{1,2, \ldots, n\}$ ) is strongly Rayleigh if its probability generating polynomial $g$ is real-stable; that is, for $z_{i} \in \mathbb{C}$, if $\operatorname{Im}\left(z_{i}\right)>0$ for all $z_{i}$, then $g\left(z_{1}, \cdots, z_{n}\right) \neq 0$.

We can efficiently sample from strongly Rayleigh distributions by MCMC (with polynomial bound on mixing time)

## Efficient Sampling from SR Distributions

Theorem (Li et al., 2016). Let $\pi$ be a strongly Rayleigh distribution over $[n]$, we can efficiently sample from $\pi$ by sampling from its symmetric homogenization $\pi_{s h}$; for $S \subset[2 n]$, define

$$
\pi_{s h}(S):=\left\{\begin{array}{l}
\pi(S \cap[n])\binom{n}{S \cap[n]}^{-1}, \quad \text { if }|S|=n \\
0, \quad \text { otherwise }
\end{array}\right.
$$

in particular, $\pi_{s h}$ is also strongly Rayleigh and the mixing time of a Gibbsexchange sampler with initial set $S_{0}$ is bounded as

$$
\tau(\epsilon) \leq 2 n^{2}\left(\log \binom{n}{\left|S_{0}\right|}+\log \pi\left(S_{0}\right)^{-1}+\log \epsilon^{-1}\right)
$$

## Relationship between PGCs and SR Distributions



## Relationship between PGCs and SR Distributions



## Not All SR Distributions have Compact PGCs (Bläser 2023)

Let $K_{m, n}=(U \cup V, E)$ be a complete bipartite graph, the signed double function generating polynomial is defined as

$$
D F_{m, n}(e)=\sum_{F, H}(-1)^{|F|+|H|} \prod_{(i, j) \in F} e_{i, j} \prod_{\left(i^{\prime}, j^{\prime}\right) \in H} e_{i^{\prime}, j^{\prime}}
$$

where the sum is taken over all partial functions $U \rightarrow V$ and $V \rightarrow U$, respectively. Each pair of $(F, H)$ is a double function of $K_{m, n}$.


Figure 4. The thick edges are a matching of size two.


Figure 5. The thick edges form a total function $U \rightarrow V$, which is not injective.


Figure 6. The thick edges form a partial function from $V$ to $U$.


Figure 7. A double function.

## Not All SR Distributions have Compact PGCs (Bläser 2023)

$$
D F_{m, n}(e)=\sum_{F, H}(-1)^{|F|+|H|} \prod_{(i, j) \in F} e_{i, j} \prod_{\left(i^{\prime}, j^{\prime}\right) \in H} e_{i^{\prime}, j^{\prime}}
$$

$\square$ Generalize to bipartite multigraph $K_{m, n}^{(d)}$
d : each edge from U to V has d copies

$$
D F_{m, n}^{(d)}\left(e^{(d)}\right)=\sum_{F, H}(-1)^{|F|+|H|} \prod_{(i, j) \in F \backslash H} \sum_{\delta=1}^{d} e_{i, j}^{(\delta)} \prod_{\left(i^{\prime}, j^{\prime}\right) \in H \cap F} \sum_{1 \leq \delta^{\prime}<\gamma \leq d} e_{i, j^{\prime}, j^{\prime}}^{\left(\delta^{\prime}\right)}, e_{i^{(\lambda), j j^{\prime}}}^{\left(\lambda^{\prime}\right)} \prod_{\left(i^{\prime}, j^{\prime}\right) \in H \backslash F \mid} \sum_{\delta^{\prime \prime}=1}^{d} e_{i}^{\left(i^{\prime \prime}, j^{\prime \prime}\right)}
$$

$$
D F_{n, n}^{(n+2)} \text { is real-stable and its evaluation is \#P-hard. }
$$

$$
D F_{n, n}^{(n+2)} \text { does not define an SR distribution as it has negative coefficients }
$$

## Not All SR Distributions have Compact PGCs (Bläser 2023)

Definition. For a polynomial $f\left(z_{1}, \ldots, z_{n}\right)$ with $z_{i}$ of degree $k_{i}$, the inversion of $f$ is defined as $\prod_{i} z_{i}{ }^{k_{i}} f\left(-1 / z_{1}, \ldots,-1 / z_{i}, \ldots,-1 / z_{n}\right)$.

The inversion of a real stable polynomial is also real stable

Let $P_{n}$ be the inversion of $D F_{n, n}^{(n+2)}$, then $P_{n}$ is a mutilinear and real stable polynomial with all coefficients non-negative.

Theorem (Bläser, 2023). Assuming $P^{\# P} \nsubseteq P /$ Poly. Let $\hat{P}_{n}$ be the normalized $P_{n}$, then $\hat{P}_{n}$ cannot be represented as polynomial-size PGCs.

## Relationship between PGCs and SR Distributions



## Probabilistic generating circuits seem awfully general.

Are all tractable probabilistic models probabilistic generating circuits?


## Questions answered today

1. What are probabilistic queries? Are current models tractable? (Guy)
2. What are probabilistic circuits and why are they tractable? (Guy)
3. What is the connection to logical circuit languages? (YooJung)
4. How do I compile my favorite model into a circuit? (YooJung)
5. How are circuit size and tractability related? (YooJung)
6. What's the most impressive query we can efficiently compute? (YooJung)
7. Are all tractable distributions probabilistic circuits? (Guy)
8. How to learn probabilistic circuits from data? (Guy)

## Questions answered today

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Building Probabilistic Circuits

## Information Prior Knowledge domain assumptions constraints other models <br> Data <br> experimental data samples measurements <br> learning <br> Circuits <br> decomposability smoothness <br> determinism <br> compatibility <br> Structure <br>  <br> Parameters <br> $\boldsymbol{\theta}, \boldsymbol{w} \quad$ generative discriminative Bayesian <br> credal

## Origins: Compilation

## Compiling probabilistic graphical models

## Arithmetic circuits

(Darwiche 2002, 2003, 2009)

- Compile a given Bayesian network into an arithmetic circuit-a smooth, decomposable and deterministic PCs
$\square$
Either via logic encoding of Bayesian network + knowledge compilation
$\square$ Or record "execution trace" (sum and product operations) of traditional inference algorithms (junction tree, variable elimination)



## Compilation

Selected references
Logic circuits, interplay between structural properties and tractable reasoning (Darwiche et al. 2002a)
Converting probabilistic graphical models via knowledge compilation
(Darwiche 2002)

## Logic circuit compilers

(Darwiche 2004; Muise et al. 2012; Bova et al. 2015; Lagniez et al. 2017; Oztok et al. 2018)
Neuro-symbolic models using logic circuits
(Ahmed et al. 2022a,b)

Parameter Learning

## Gradient descent (of course)

PCs are computational graphs
$\square$ Hence we can just learn them as any other neural net using SGD
$\square$ Use re-parameterization if parameters should satisfy constraints:
soft-max for sum-weights (non-negative, sum-to-one)
soft-plus for variances
low-rank plus diagonal for covariance matrices

- Allows for conditional distributions


## Conditional PCs

(Shao et al. 2019)

chain rule of probabilities


## Maximum Iikelihood (frequentist)

PCs can be interpreted as hierarchical latent variable models, where each sum node corresponds to a discrete latent variable (Peharz et al. 2016). This allows to perform classical maximum-likelihood estimation.


## Closed-form maximum Ikellhood

When the circuit is deterministic, there is even an closed-form ML solution, which is incredible fast:

```
julia> using ProbabilisticCircuits;
julia> data, structure = load(...);
julia> num_examples(data)
17412
julia> num_edges(structure)
270448
julia> @btime estimate_parameters(structure, data);
    63.585 ms (1182350 allocations: 65.97 MiB)
```

Custom SIMD and CUDA kernels to parallelize over layers and training examples.

## Expectation-Maximization

When the PC is not deterministic, we can still apply expectation-maximization (Peharz et al. 2016). EM can piggy-back on autodfiff:

```
train_x, valid_x, test_x = get_mnist_images([7])
graph = Graph.poon_domingos_structure(shape=(28,28), delta=[7])
args = EinsumNetwork.Args(num_var=train_x.shape [1], num_dims=1,
    num_classes=1, num_sums=28,
    num_input_distributions=28,
    exponential_family=EinsumNetwork.BinomialArray,
    exponential_family_args={'N':255},
    online_em_frequency=1, online_em_stepsize=0.05)
```

PC = EinsumNetwork.EinsumNetwork(graph, args)
PC.initialize()
PC.to('cuda')

## Expectation-Maximization

```
for epoch_count in range(10):
    train_ll, valid_ll, test_ll = compute_loglikelihood()
    start_t = time.time()
    for idx in get_batches(train_x, 100):
        outputs = PC.forward(train_x[idx, :])
        log_likelihood = EinsumNetwork.log_likelihoods(outputs).sum()
        log_likelihood.backward()
        PC.em_process_batch()
    print_performance(epoch_count, train_ll, valid_ll, test_ll, time.time() - start_t)
```


## Expectation-Maximization

```
# train sample: 5175
# parameters: 1573486
```

[epoch 0] train LL -140936.80
[epoch 1] train LL -15916.14
[epoch 2] train LL -10865.67
[epoch 3] train LL -10388.53
[epoch 4] train LL -10264.11
[epoch 5] train LL -10212.66
[epoch 6] train LL -10192.21
[epoch 7] train LL -10153.97
[epoch 8] train LL -10112.95
[epoch 9] train LL -10093.31

| valid LL | -140955.72 | test LL | -141033.80 |
| :--- | :--- | :--- | :--- |
| valid LL | -15693.25 | test LL | -15976.43 |
| valid LL | -10616.72 | test LL | -10943.56 |
| valid LL | -10158.84 | test LL | -10475.49 |
| valid LL | -10041.66 | test LL | -10352.59 |
| valid LL | -10001.09 | test LL | -10319.35 |
| valid LL | -9965.98 | test LL | -10314.84 |
| valid LL | -9920.09 | test LL | -10261.41 |
| valid LL | -9882.48 | test LL | -10236.34 |
| valid LL | -9862.15 | test LL | -10200.94 |

... elapsed time 3.621 sec
... elapsed time 3.438 sec
... elapsed time 3.436 sec
... elapsed time 3.473 sec
... elapsed time 3.497 sec
... elapsed time 3.584 sec
... elapsed time 3.508 sec
... elapsed time 3.446 sec
... elapsed time 3.579 sec
... elapsed time 3.483 sec

Peharz et al., "Einsum Networks: Fast and Scalable Learning of Tractable Probabilistic Circuits", 2020

## Structure Learning

## Region graphs

Laying out the PC structure on a high level
$\square$ Region graphs (RGs) describe decompositional structure
$\square$ RGs are bipartite, directed graphs containing regions $(\mathcal{R})$ and partitions $(\mathcal{P})$

$\square$Input and output nodes of the RG are regions
$\square$ Regions have a scope (receptive field), denoted as $s c(\mathcal{R}) \subseteq \mathbf{X}$
$\square$ For every partition $\mathcal{P}$ it holds that

$$
\begin{array}{ll}
s c\left(\mathcal{R}_{\text {out }}\right)=\bigcup_{\mathcal{\mathcal { R } _ { \text { in } } \in \text { inputs } ( \mathcal { P } )}} s c\left(\mathcal{R}_{\text {in }}\right) \\
s c\left(\mathcal{R}^{\prime}\right) \cap s c\left(\mathcal{R}^{\prime \prime}\right)=\emptyset, & \mathcal{R}^{\prime} \neq \mathcal{R}^{\prime \prime} \in \operatorname{inputs}(\mathcal{P})
\end{array}
$$

## Example region graph



## From region graphs to PCs

-••


## From region graphs to PCs

Equip each input region with non-linear units
©


## From region graphs to PCs

Equip each internal region with sum nodes
©®


## From region graphs to PCs

©e


## From region graphs to PCs

Equip partitions with products, combining units
©e

$\mathcal{P}$


## From region graphs to PCs

Equip partitions with products, combining units
-••

$\mathcal{P}$


## From region graphs to PCs

Connect products to sum units above


## From region graphs to PCs

- Equip each input region (leaf) $\mathcal{R}$ with $K$ units $\phi_{1}, \ldots, \phi_{K}$, which are non-linear functions over $s c(\mathcal{R})$. Usually, $\phi_{1}, \ldots, \phi_{K}$ are probability densities. $K$ can be different for each input region.
- Equip each other region with $K$ sum units. $K$ can be different for each internal region. Often, for the root region $K=1$, when PC is used as density estimator.
$\square$ Equip each partition $\mathcal{P}$ with as many products as there are combinations of units in the input regions to $\mathcal{P}$, selecting one unit from each region. Formally, if $\mathcal{P}$ has input regions $\mathcal{R}_{1}, \mathcal{R}_{2} \ldots, \mathcal{R}_{I}$, insert one product $\prod_{i=1}^{I} u_{i}$ for each $\left(u_{1}, u_{2}, \ldots, u_{I}\right) \in \mathcal{R}_{1} \times \mathcal{R}_{2} \times \cdots \times \mathcal{R}_{I}$.
- Connect each $\prod_{i=1}^{I} u_{i}$ in $\mathcal{P}$ to all sum units in the output regions of $\mathcal{P}$.


## From region graphs to PCs

$\square$ Resulting PC has alternating sum and product units (not a strong constraint)
$\square$ We can easily scale the PC (overparameterize, increase expressivity) by equipping regions with more units
$\square$ RGs can be seen as a vectorized version of PCs - each region and partition can be seen as as a module
$\square$ Resulting PC will be smooth and decomposable, i.e., we can integrate, marginalize, and take conditionals
$\square$ After the PC has been constructed, we might discard the RG

## Scaling up image models

Latent Variable Distillation

| Dataset | TPMs |  |  |  |  |  |  | DGMs |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | LVD (ours) | HCLT | EiNet | RAT-SPN |  | Glow | RealNVP | BIVA |  |  |
| ImageNet32 | $\mathbf{4 . 3 8}$ | 4.82 | 5.63 | 6.90 |  | 4.09 | 4.28 | 3.96 |  |  |
| ImageNet64 | $\mathbf{4 . 1 2}$ | 4.67 | 5.69 | 6.82 |  | 3.81 | 3.98 | - |  |  |
| CIFAR | $\mathbf{4 . 3 7}$ | 4.61 | 5.81 | 6.95 |  | 3.35 | 3.49 | 3.08 |  |  |



## How to construct and learn RGs?

## Random regions graphs

The "no-learning" option
Generating a random region graph, by recursively splitting $\mathbf{X}$ into two random parts:


## Image-tallored circuit structure

"Recursive image slicing"
Images yield a natural region graph by using axis-aligned splits:
$\square$ Start with the full image (=output region)

- Define partitions by applying horizontal and vertical splits
$\square$ Recurse on the newly generated sub-images (internal regions)
- Structure somewhat reminiscent to convolutions
$\square$ Generates RGs which are "true DAGs," i.e. regions get re-used






## Data-driven structure Iearning

"Recursive data slicing"

Expand regions with clustering


## Data-driven structure Iearning

"Recursive data slicing"

Number of clusters = number of partitions


## Data-driven structure learning

"Recursive data slicing"

Try to find independent groups of variables (e.g. independence tests)


## Data-driven structure learning

"Recursive data slicing"

Success $\rightarrow \boldsymbol{p a r t i t i o n}$ into new regions


## Data-driven structure learning

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Try to find independent groups of variables (e.g. independence tests)


## Data-driven structure learning

"Recursive data slicing"

Success $\rightarrow \boldsymbol{p a r t i t i o n}$ into new regions


## Data-driven structure Iearning

"Recursive data slicing"

Single variable


## Data-driven structure Iearning

"Recursive data slicing"

Single variable $\rightarrow$ input region


## Data-driven structure learning

"Recursive data slicing"

Expand regions with clustering


## Data-driven structure learning

"Recursive data slicing"

Number of clusters = number of partitions

And so on...


## Data-driven structure learning

"Recursive data slicing"
$\square$ Stopping conditions: minimal number of features, samples, depth, ...
$\square$ Clustering ratios also deliver (initial) parametersSmooth \& Decomposable Circuits
$\square$ Tractable integration


## LearnSPN

Selected references

- ID-SPN (Rooshenas et al. 2014)
- LearnSPN-b/T/B (Vergari et al. 2015)
- For heterogeneous data (Molina et al. 2018)
- Using k-means (Butz et al. 2018) or SVD splits (Adel et al. 2015)
$\square$ Learning DAGs (Dennis et al. 2015; Jaini et al. 2018)
■ Approximating independence tests (Di Mauro et al. 2018)


## Cutset networks

Besides clustering, decision tree learning can be used as PC learner. Cutset networks, decision trees over simple probabilistic models (Chow-Liu trees) (Rahman et al. 2014):


Cutset networks can easily be converted into smooth, decomposable and deterministic PCs.

## Decision trees as PCs

Also vanilla decision tree learners can be used to learn PCs, by augmenting the leaves with distributions over inputs (Correia et al. 2020). Allows to treat missing features and outlier detection.


## Information Prior Knowledge domain assumptions constraints other models <br> Data <br> experimental data samples measurements <br> learning <br> Circuits <br> decomposability smoothness <br> determinism <br> compatibility <br> Structure <br>  <br> Parameters <br> $\boldsymbol{\theta}, \boldsymbol{w}$ generative discriminative Bayesian <br> credal

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