Polynomial semantics of probabilistic circuits

Oliver Broadrick, UCLA

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Based on joint work with Honghua Zhang and Guy Van den Broeck

▶ Tractable probabilistic models

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- Probabilistic circuits
- ▶ Several circuit semantics in the literature ...
- ▶ are equivalent! (for binary random variables)
- ▶ And, don't all extend to non-binary variables

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AI research

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▶ Deep learning and formal methods: "neuro-symbolic AI"

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AI research

- ▶ Deep learning and formal methods: "neuro-symbolic AI"
- ▶ Applications: images, language, audio, medicine, science, economics, etc.

X_1	X_2	\Pr
0	0	.1
0	1	.2
1	0	.3
1	1	.4

► Expressive-efficient representation

X_1	X_2	\Pr
0	0	.1
0	1	.2
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- ► Expressive-efficient representation
- ▶ Tractable inference





X_2	\Pr	
0	.1	
1	.2	
0	.3	
1	.4	
	$egin{array}{c} X_2 \\ 0 \\ 1 \\ 0 \\ 1 \end{array}$	$\begin{array}{c c} X_2 & \Pr \\ \hline 0 & .1 \\ 1 & .2 \\ 0 & .3 \\ 1 & .4 \\ \end{array}$

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$$X = Y \sqcup Z$$
, then what is $\Pr[Y = y]$?

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In general:

$$\Pr[\boldsymbol{Y} = \boldsymbol{y}] = \sum_{\boldsymbol{z}} \Pr[\boldsymbol{Y} = \boldsymbol{y}, \boldsymbol{Z} = \boldsymbol{z}]$$

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For example:

$$Pr[X_1 = 1] = Pr[X_1 = 1, X_2 = 0] + Pr[X_1 = 1, X_2 = 1]$$

= 0.3 + 0.4
= 0.7

 $\begin{array}{c|cc} X_1 & X_2 & \Pr\\ \hline 0 & 0 & .1 \end{array}$

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For example:
$$Pr[X_1 = 1] = Pr[X_1 = 1, X_2 = 0] + Pr[X_1 = 1, X_2 = 0.3 + 0.4]$$
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Goal: Find a model of polysize that supports marginal inference in polytime, for as large a set of probability distributions as possible.

1

▶ Bayesian Networks (of bounded treewidth) (BNs)

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- Probabilistic Circuits!

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Marginal Inference:



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$$\Pr[X_1 = 1]?$$

Set $x_1 = 1, \, \bar{x}_1 = 0, \, x_2 = 1,$
 $\bar{x}_2 = 1$

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 $p(x_1, x_2, \dots, x_n, \bar{x}_1, \bar{x}_2, \dots, \bar{x}_n) = .1\bar{x}_1\bar{x}_2 + .2\bar{x}_1x_2 + .3x_1\bar{x}_2 + .4x_1x_2$



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Circuit Semantics

Polynomial	Notation	Inference	Citation
Network polynomial	$p(x_1,\ldots,x_n,\bar{x}_1,\ldots,\bar{x}_n)$	\checkmark	Darwiche [2003]
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Generating function	$g(x_1,\ldots,x_n)$	\checkmark	Zhang et al. [2021]

Circuit Semantics

[2009]

How do they relate?



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$$p(x_1, x_2) = .2x_1 + .1x_2 + .1$$

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Can we do marginal inference? Relation to network polynomial?

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Transformation from network to a likelihood: $p(x_1, x_2, \bar{x}_1, \bar{x}_2) = .1\bar{x}_1\bar{x}_2 + .2\bar{x}_1x_2 + .3x_1\bar{x}_2 + .4x_1x_2$

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Can we do marginal inference? Relation to network polynomial?

Transfe	ormation from network to a likelihood.
$p(x_1, x)$	$(x_1, x_2) = .1x_1x_2 + .2x_1x_2 + .3x_1x_2 + .4x_1x_2$
$p(x_1, x)$	$(2, 1 - x_1, 1 - x_2)$
= .1(1)	$-x_1(1-x_2) + .2(1-x_1)x_2 + .3x_1(1-x_2) + .4x_1x_2$
$= .2x_1$	$+.1x_2 + .1$

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Idea:
$$\left(\prod_{i=1}^{n} (x_i + \bar{x}_i)\right) p\left(\frac{x_1}{x_1 + \bar{x}_1}, \dots, \frac{x_n}{x_n + \bar{x}_n}\right)$$

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= $.1\bar{x}_1\bar{x}_2 + .2\bar{x}_1x_2 + .3x_1\bar{x}_2 + .4x_1x_2 = p(x_1, x_2, \bar{x}_1, \bar{x}_2)$

Wait a second...
$$\left(\prod_{i=1}^{n} (x_i + \bar{x}_i)\right) p\left(\frac{x_1}{x_1 + \bar{x}_1}, \dots, \frac{x_n}{x_n + \bar{x}_n}\right)$$

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Theorem 2 (Strassen). You can remove the divisions in polynomial time!

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Lemma 3. For a circuit computing f of degree d, we can obtain circuits computing $H_0[f], H_1[f], \ldots, H_d[f]$ the homogeneous parts of f, i.e. $H_i[f]$ has degree i and $f = \sum_i H_i[f]$.

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Idea: Move divisions to the root using a/b + c/d = (ad + bc)/bd and a/b * c/d = ac/bd.

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Idea: Move divisions to the root using a/b + c/d = (ad + bc)/bd and a/b * c/d = ac/bd. Then for circuit a/b computing polynomial f = a/b of degree d, assume b(0) = 1, and we have

$$H_i[f] = H_i[a(1 + (1 - b) + (1 - b)^2 + \dots + (1 - b)^d].$$

Progress update



Progress update



Progress update



Generating functions: Why?

Generating functions: Why?



Generating functions: Why?



$$g(x) = .1 + .2x_2 + .3x_1 + .4x_1x_2$$

Can we do marginal inference?

X_1	X_2	\Pr
0	0	.1
0	1	.2
1	0	.3
1	1	.4

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Can we do marginal inference? For $X_i = 1$, set $x_i = t$



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Can we do marginal inference? For $X_i = 1$, set $x_i = t$ For $X_i = 0$, set $x_i = 0$ For $X_i = ?$, set $x_i = 1$ $\Pr[X_i = 1]^2$

X_1	X_2	\Pr
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0	1	.2
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$$\begin{aligned} &\Pr[X_1 = 1]?\\ &g(t,1) = .1 + .2 + .3t + .4t = .3 + .7t \end{aligned}$$
Generating functions

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Can we do marginal inference? For $X_i = 1$, set $x_i = t$ For $X_i = 0$, set $x_i = 0$ For $X_i = ?$, set $x_i = 1$ $\Pr[X_1 = 1]?$ g(t, 1) = .1 + .2 + .3t + .4t = .3 + .7t

Relation to network polynomial?

X_1	X_2	Pr
$\frac{1}{0}$	0	.1
0	1	.2
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X_1	X_2	Pr
0	0	.1
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Theorem 4. Let Pr be a probability distribution on n binary random variables. Then a circuit of size s computing the probability generating function for Pr can be transformed to a circuit of size $O(sn^2)$ computing the network polynomial for Pr.

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Idea:
$$\left(\prod_{i=1}^{n} \bar{x}_i\right) g\left(\frac{x_1}{\bar{x}_1}, \frac{x_2}{\bar{x}_2}, \dots, \frac{x_n}{\bar{x}_n}\right)$$

Theorem 4. Let \Pr be a probability distribution on n binary random variables. Then a circuit of size s computing the probability generating function for \Pr can be transformed to a circuit of size $O(sn^2)$ computing the network polynomial for \Pr .

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$$(\bar{x}_1\bar{x}_2)\left(.1+.2\frac{x_2}{\bar{x}_2}+.3\frac{x_1}{\bar{x}_1}+.4\frac{x_1}{\bar{x}_1}\frac{x_2}{\bar{x}_2}\right)$$

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$$(\bar{x}_1 \bar{x}_2) \left(.1 + .2 \frac{x_2}{\bar{x}_2} + .3 \frac{x_1}{\bar{x}_1} + .4 \frac{x_1}{\bar{x}_1} \frac{x_2}{\bar{x}_2} \right)$$

= $.1 \bar{x}_1 \bar{x}_2 + .2 \bar{x}_1 x_2 + .3 x_1 \bar{x}_2 + .4 x_1 x_2$

Theorem 4. Let \Pr be a probability distribution on n binary random variables. Then a circuit of size s computing the probability generating function for \Pr can be transformed to a circuit of size $O(sn^2)$ computing the network polynomial for \Pr .

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$$\left(\prod_{i=1}^{n} \bar{x}_{i}\right) g\left(\frac{x_{1}}{\bar{x}_{1}}, \frac{x_{2}}{\bar{x}_{2}}, \dots, \frac{x_{n}}{\bar{x}_{n}}\right)$$

$$(\bar{x}_1 \bar{x}_2) \left(.1 + .2 \frac{x_2}{\bar{x}_2} + .3 \frac{x_1}{\bar{x}_1} + .4 \frac{x_1}{\bar{x}_1} \frac{x_2}{\bar{x}_2} \right) = .1 \bar{x}_1 \bar{x}_2 + .2 \bar{x}_1 x_2 + .3 x_1 \bar{x}_2 + .4 x_1 x_2 = p(x, \bar{x})$$

Theorem 4. Let \Pr be a probability distribution on n binary random variables. Then a circuit of size s computing the probability generating function for \Pr can be transformed to a circuit of size $O(sn^2)$ computing the network polynomial for \Pr .

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Example: Starting with $g(x) = .1 + .2x_2 + .3x_1 + .4x_1x_2$, we form

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Strassen to the rescue!

- $p(x, \bar{x})$ Network polynomial p(x)
 - Likelihood polynomial
- g(x)Generating function
- $\hat{p}(x)$ Fourier transform



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Proposition 1. For binary random variables, probability generating functions g(x) and Fourier polynomials $\hat{p}(x)$ are the same function(!), on respective domains $\{-1,1\}^n$ and $\{0,1\}^n$, up to the bijection $\phi: \{0,1\} \rightarrow \{-1,1\}$ given by $\phi(b) = (-1)^b$ applied bitwise.

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Theorem 5. For $k \ge 4$, computing likelihoods on a circuit for g(x) is #P-hard. Proof idea: Reduce from 0, 1-permanent.

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- ▶ Can we characterize *all* tractable marginal inference?
- ▶ How can theoretically more expressive models be learned/constructed in practice?