# Polynomial semantics of probabilistic circuits 

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Based on joint work with Honghua Zhang and Guy Van den Broeck

## Outline

- Tractable probabilistic models


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- Probabilistic circuits


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- Several circuit semantics in the literature ...


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- Tractable probabilistic models
- Probabilistic circuits
- Several circuit semantics in the literature ...
- are equivalent! (for binary random variables)
- And, don't all extend to non-binary variables


## Probabilistic Models

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How we think about the world: models with uncertainty

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AI research

- Deep learning and formal methods: "neuro-symbolic AI"
- Applications: images, language, audio, medicine, science, economics, etc.


## The problem

| $X_{1}$ | $X_{2}$ | $\operatorname{Pr}$ |
| :---: | :---: | :---: |
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| 0 | 1 | .2 |
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## The problem

- Expressive-efficient representation

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## Marginal inference

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\text { If } \boldsymbol{X}=\boldsymbol{Y} \sqcup \boldsymbol{Z} \text {, then what is } \operatorname{Pr}[\boldsymbol{Y}=\boldsymbol{y}] \text { ? }
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## Marginal inference

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For example:

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\operatorname{Pr}\left[X_{1}=1\right] & =\operatorname{Pr}\left[X_{1}=1, X_{2}=0\right]+\operatorname{Pr}\left[X_{1}=1, X_{2}=1\right] \\
& =0.3+0.4 \\
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Goal: Find a model of polysize that supports marginal inference in polytime, for as large a set of probability distributions as possible.

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If $X_{i}=b$ for $b \in\{0,1\}$, set $x_{i}=b$ and $\bar{x}_{i}=1-b$.

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If $X_{i}=b$ for $b \in\{0,1\}$, set
$x_{i}=b$ and $\bar{x}_{i}=1-b$.
If $X_{i}$ is not assigned, set $x_{1}=1$ and $\bar{x}_{1}=1$.

## Probabilistic Circuits



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## Probabilistic Circuits


"Network polynomial"

## Circuit Semantics

| Polynomial | Notation | Inference | Citation |
| :--- | :--- | :--- | :--- |
| Network polynomial | $p\left(x_{1}, \ldots, x_{n}, \bar{x}_{1}, \ldots, \bar{x}_{n}\right)$ | $\checkmark$ | Darwiche [2003] |

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| Fourier transform | $\hat{p}\left(x_{1}, \ldots, x_{n}\right)$ | $\checkmark$ | Yu et al. [2023] |

## How do they relate?

$\begin{array}{ll}p(x, \bar{x}) & \text { Network polynomial } \\ p(x) & \text { Likelihood polynomial }\end{array}$
$g(x) \quad$ Generating function
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## Likelihood polynomials

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Transformation from network to a likelihood: $p\left(x_{1}, x_{2}, \bar{x}_{1}, \bar{x}_{2}\right)=.1 \bar{x}_{1} \bar{x}_{2}+.2 \bar{x}_{1} x_{2}+.3 x_{1} \bar{x}_{2}+.4 x_{1} x_{2}$

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## Transformation from likelihood to network

Theorem 1. Let $\operatorname{Pr}$ be a probability distribution on $n$ binary random variables. Then a circuit of size s computing the likelihood polynomial for $\operatorname{Pr}$ can be transformed to a circuit of size $O\left(s n^{2}\right)$ computing the network polynomial for $\operatorname{Pr}$.

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## Divisions?

Wait a second... $\left(\prod_{i=1}^{n}\left(x_{i}+\bar{x}_{i}\right)\right) p\left(\frac{x_{1}}{x_{1}+\bar{x}_{1}}, \ldots, \frac{x_{n}}{x_{n}+\bar{x}_{n}}\right)$

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Lemma 3. For a circuit computing $f$ of degree d, we can obtain circuits computing $H_{0}[f], H_{1}[f], \ldots, H_{d}[f]$ the homogeneous parts of $f$, i.e. $H_{i}[f]$ has degree $i$ and $f=\sum_{i} H_{i}[f]$.

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Idea: Move divisions to the root using $a / b+c / d=(a d+b c) / b d$ and $a / b * c / d=a c / b d$.
Then for circuit $a / b$ computing polynomial $f=a / b$ of degree $d$, assume $b(0)=1$, and we have

$$
H_{i}[f]=H_{i}\left[a\left(1+(1-b)+(1-b)^{2}+\ldots+(1-b)^{d}\right] .\right.
$$

## Progress update

$p(x, \bar{x}) \quad$ Network polynomial
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## Generating functions:

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Circuits computing network polynomials (S, D, positive)


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## Generating functions

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g(x)=.1+.2 x_{2}+.3 x_{1}+.4 x_{1} x_{2}
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Can we do marginal inference?
For $X_{i}=1$, set $x_{i}=t$

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## Generating functions

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g(x)=.1+.2 x_{2}+.3 x_{1}+.4 x_{1} x_{2}
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Can we do marginal inference?
For $X_{i}=1$, set $x_{i}=t$
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$$
\text { For } X_{i}=?, \text { set } x_{i}=1
$$

$$
\begin{aligned}
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& p\left(x_{1}, x_{2}, 1,1\right)=.1+.2 x_{2}+.3 x_{1}+.4 x_{1} x_{2}=g\left(x_{1}, x_{2}\right)
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Transformation from generating to network

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Theorem 4. Let $\operatorname{Pr}$ be a probability distribution on $n$ binary random variables. Then a circuit of size s computing the probability generating function for $\operatorname{Pr}$ can be transformed to a circuit of size $O\left(s n^{2}\right)$ computing the network polynomial for $\operatorname{Pr}$.

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## Progress update

$p(x, \bar{x}) \quad$ Network polynomial
$p(x) \quad$ Likelihood polynomial
$g(x) \quad$ Generating function
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Proposition 1. For binary random variables, probability generating functions $g(x)$ and Fourier polynomials $\hat{p}(x)$ are the same function(!), on respective domains $\{-1,1\}^{n}$ and $\{0,1\}^{n}$, up to the bijection $\phi:\{0,1\} \rightarrow\{-1,1\}$ given by $\phi(b)=(-1)^{b}$ applied bitwise .

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What have we done?

- several distinct circuit-based models are equally succinct
- distinct inference algorithms in a common framework

Non-binary distributions?

## Non-binary distributions?

Let $\operatorname{Pr}: K^{n} \rightarrow \mathbb{R}$ be a probability mass function with $K=\{0,1,2, \ldots, k-1\}$. Then the probability generating polynomial of Pr is

$$
\begin{equation*}
g(x)=\sum_{\left(d_{1}, d_{2}, \ldots, d_{n}\right) \in K^{n}} \operatorname{Pr}\left(d_{1}, \ldots, d_{n}\right) x_{1}^{d_{1}} x_{2}^{d_{2}} \cdots x_{n}^{d_{n}} \tag{1}
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Theorem 5. For $k \geq 4$, computing likelihoods on a circuit for $g(x)$ is \#P-hard. Proof idea: Reduce from 0,1-permanent.

## Conclusion

What we did:

- Several distinct circuit-based models are equally succinct
- Distinct inference algorithms in a common framework
- Inference is hard in circuits computing generating functions for $k \geq 4$ categories


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What's next?

- Are there more succinct tractable representations? e.g., do we need multilinearity?
- Can we characterize all tractable marginal inference?
- How can theoretically more expressive models be learned/constructed in practice?

